

# Matrix Equations and Model Reduction

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# **Outline**

- Introduction
- Mathematical Basics
- Model Reduction by Projection
- 4 Interpolatory Model Reduction
- Balanced Truncation
- 6 Solving Large-Scale Matrix Equations
- Final Remarks

# **Outline**

- Introduction
  - Model Reduction for Dynamical Systems
  - Application Areas
  - Motivating Examples
- Mathematical Basics
- Model Reduction by Projection
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 Mathematical Basics
 MOR by Projection
 RatInt
 Balanced Truncation
 Matrix Equations

#### Introduction

Model Reduction — Abstract Definition

#### Problem

Given a physical problem with dynamics described by the states  $x \in \mathbb{R}^n$ , where n is the dimension of the state space.

Because of redundancies, complexity, etc., we want to describe the dynamics of the system using a reduced number of states.

This is the task of model reduction (also: dimension reduction, order reduction).

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#### Introduction

Model Reduction for Dynamical Systems

# Dynamical Systems

$$\Sigma : \left\{ \begin{array}{lcl} \dot{x}(t) & = & f(t, x(t), u(t)), & x(t_0) = x_0, \\ y(t) & = & g(t, x(t), u(t)) \end{array} \right.$$

with

- states  $x(t) \in \mathbb{R}^n$ ,
- inputs  $u(t) \in \mathbb{R}^m$ ,
- outputs  $y(t) \in \mathbb{R}^q$ .



### Original System

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- states  $\hat{x}(t) \in \mathbb{R}^r$ ,  $r \ll n$
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#### Reduced-Order Model (ROM)

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#### Goal

 $\|y - \hat{y}\| < \text{tolerance} \cdot \|u\|$  for all admissible input signals.

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#### Goal:

 $\|y - \hat{y}\| < \text{tolerance} \cdot \|u\|$  for all admissible input signals.

Secondary goal: reconstruct approximation of x from  $\hat{x}$ .

Parameter-Dependent Dynamical Systems

### Dynamical Systems

$$\Sigma(p): \begin{cases} E(p)\dot{x}(t;p) &= f(t,x(t;p),u(t),p), & x(t_0) = x_0, \\ y(t;p) &= g(t,x(t;p),u(t),p) \end{cases}$$
 (a)

#### with

- (generalized) states  $x(t; p) \in \mathbb{R}^n$  ( $E \in \mathbb{R}^{n \times n}$ ),
- inputs  $u(t) \in \mathbb{R}^m$ ,
- outputs  $y(t; p) \in \mathbb{R}^q$ , (b) is called output equation,
- $p \in \Omega \subset \mathbb{R}^d$  is a parameter vector,  $\Omega$  is bounded.

#### **Applications:**

- Repeated simulation for varying material or geometry parameters, boundary conditions,
- Control, optimization and design.

Requirement: keep parameters as symbolic quantities in ROM.

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**Linear Systems** 

# Linear, Time-Invariant (LTI) Systems

$$\begin{array}{lcl} E\dot{x} & = & f(t,x,u) & = & Ax+Bu, \quad E,A\in\mathbb{R}^{n\times n}, \qquad B\in\mathbb{R}^{n\times m}, \\ y & = & g(t,x,u) & = & Cx+Du, \quad C\in\mathbb{R}^{q\times n}, & D\in\mathbb{R}^{q\times m}. \end{array}$$

Linear Systems

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#### Linear, Time-Invariant Parametric Systems

$$E(p)\dot{x}(t;p) = A(p)x(t;p) + B(p)u(t),$$
  
$$y(t;p) = C(p)x(t;p) + D(p)u(t),$$

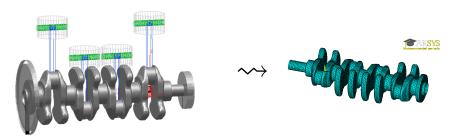
where  $A(p), E(p) \in \mathbb{R}^{n \times n}, B(p) \in \mathbb{R}^{n \times m}, C(p) \in \mathbb{R}^{q \times n}, D(p) \in \mathbb{R}^{q \times m}$ .

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# **Application Areas**

Structural Mechanics / Finite Element Modeling

since  $\sim$ 1960ies



- Resolving complex 3D geometries ⇒ millions of degrees of freedom.
- Analysis of elastic deformations requires many simulation runs for varying external forces.

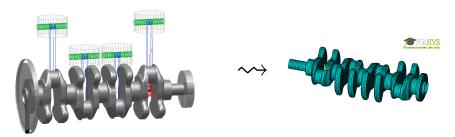
Standard MOR techniques in structural mechanics: modal truncation, combined with Guyan reduction (static condensation)  $\leadsto$  Craig-Bampton method.

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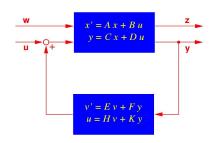
since  $\sim 1980$ ies

#### Feedback Controllers

A feedback controller (dynamic compensator) is a linear system of order N, where

- input = output of plant,
- output = input of plant.

Modern (LQG- $/\mathcal{H}_2$ - $/\mathcal{H}_\infty$ -) control design: N > n.



Practical controllers require small N ( $N\sim 10$ , say) due to

- real-time constraints,
- increasing fragility for larger N.
- $\implies$  reduce order of plant (n) and/or controller (N).

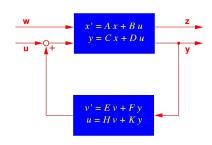
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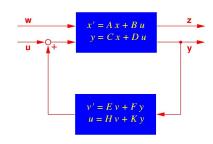
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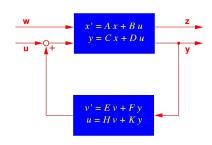
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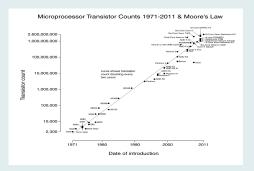
# **Application Areas**

Micro Electronics/Circuit Simulation

since  $\sim \! 1990 ies$ 

# Progressive miniaturization

- Verification of VLSI/ULSI chip design requires high number of simulations for different input signals.
- Moore's Law (1965/75) states that the number of on-chip transistors doubles each 24 months.



Source: http://en.wikipedia.org/wiki/File:Transistor\_Count\_and\_Moore'sLaw\_-\_2011.svg

Micro Electronics/Circuit Simulation

11/96

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#### **Application Areas** Micro Electronics/Circuit Simulation

## Progressive miniaturization

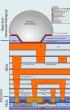
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- network topology (Kirchhoff's laws) and characteristic element/semiconductor equations.
- Increase in packing density and multilayer technology requires modeling of interconncet to ensure that thermic/electro-magnetic effects do not disturb signal transmission.

| Intel 4004 (1971)           | Intel Core 2 Extreme (quad-core) (2007) |
|-----------------------------|---|
| 1 layer, $10\mu$ technology | 9 layers, 45 <i>nm</i> technology       |
| 2,300 transistors           | > 8, 200, 000 transistors               |
| 64 kHz clock speed          | > 3 GHz clock speed.                    |

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Source: http://en.wikipedia.org/wiki/Image:Silicon\_chip\_3d.png.

#### Progressive miniaturization

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- Here: mostly MOR for linear systems, they occur in micro electronics through modified nodal analysis (MNA) for RLC networks. e.g., when
  - decoupling large linear subcircuits,
  - modeling transmission lines,
  - modeling pin packages in VLSI chips,
  - modeling circuit elements described by Maxwell's equation using partial element equivalent circuits (PEEC).

Micro Electronics/Circuit Simulation

# Progressive miniaturization

- Verification of VLSI/ULSI chip design requires high number of simulations for different input signals.
- network topology (Kirchhoff's laws) and characteristic element/semiconductor equations.

 ∼→ Clear need for model reduction techniques in order to facilitate or even. enable circuit simulation for current and future VLSI design.

Micro Electronics/Circuit Simulation

# Progressive miniaturization

- Verification of VLSI/ULSI chip design requires high number of simulations for different input signals.

→ Clear need for model reduction techniques in order to facilitate or even enable circuit simulation for current and future VLSI design.

Standard MOR techniques in circuit simulation:

Krylov subspace / Padé approximation / rational interpolation methods.

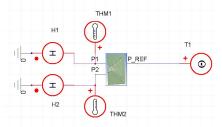
# **Application Areas**

Many other disciplines in computational sciences and engineering like

- computational fluid dynamics (CFD),
- computational electromagnetics,
- chemical process engineering,
- design of MEMS/NEMS (micro/nano-electrical-mechanical systems),
- computational acoustics,
- . . .

Electro-Thermic Simulation of Integrated Circuit (IC)

SIMPLORER<sup>®</sup> test circuit with 2 transistors.



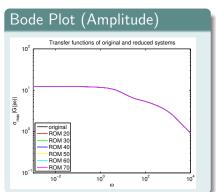
- Conservative thermic sub-system in SIMPLORER: voltage → temperature, current → heat flow.
- Original model: n = 270.593,  $m = q = 2 \Rightarrow$  Computing time (on Intel Xeon dualcore 3GHz, 1 Thread):
  - Main computational cost for set-up data  $\approx 22min$ .
  - Computation of reduced models from set-up data: 44–49sec. (r = 20-70).
  - Bode plot (MATLAB on Intel Core i7, 2,67GHz, 12GB):
     7.5h for original system, < 1min for reduced system.</li>
  - Speed-up factor: 18 including / ≥ 450 excluding reduced model generation!

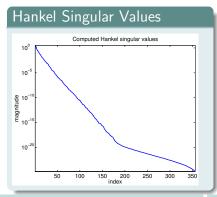
## **Motivating Examples**

#### **Electro-Thermic Simulation of Integrated Circuit (IC)**

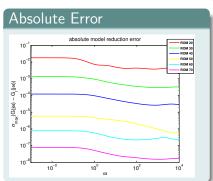
[Source: Evgenii Rudnyi, CADFEM GmbH]

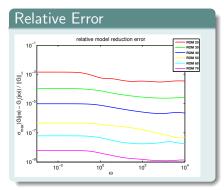
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#### Motivating Examples

#### A Nonlinear Model from Computational Neurosciences: the FitzHugh-Nagumo System

• Simple model for neuron (de-)activation [Chaturantabut/Sorensen 2009]

$$\epsilon v_t(x,t) = \epsilon^2 v_{xx}(x,t) + f(v(x,t)) - w(x,t) + g,$$
  

$$w_t(x,t) = hv(x,t) - \gamma w(x,t) + g.$$

with f(v) = v(v - 0.1)(1 - v) and initial and boundary conditions

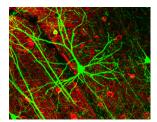
$$v(x,0) = 0,$$
  $w(x,0) = 0,$   $x \in [0,1]$ 

$$(x,0)\equiv 0, \qquad x\in [0,1]$$

$$v_{x}(0,t) = -i_{0}(t),$$
  $v_{x}(1,t) = 0,$   $t \geq 0,$ 

where 
$$\epsilon = 0.015$$
,  $h = 0.5$ ,  $\gamma = 2$ ,  $g = 0.05$ ,  $i_0(t) = 50000t^3 \exp(-15t)$ .





Source: http://en.wikipedia.org/wiki/Neuron

# **Motivating Examples**

#### A Nonlinear Model from Computational Neurosciences: the FitzHugh-Nagumo System

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$$v(x,0) = 0,$$
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where 
$$\epsilon = 0.015$$
,  $h = 0.5$ ,  $\gamma = 2$ ,  $g = 0.05$ ,  $i_0(t) = 50000t^3 \exp(-15t)$ .

- Parameter g handled as an additional input.
- Original state dimension  $n = 2 \cdot 400$ , QBDAE dimension  $N = 3 \cdot 400$ , reduced QBDAE dimension r = 26, chosen expansion point  $\sigma = 1$ .

troduction Mathematical Basics MOR by Projection RatInt Balanced Truncation Matrix Equations

# **Motivating Examples**

A Nonlinear Model from Computational Neurosciences: the FitzHugh-Nagumo System

duction Mathematical Basics MOR by Projection RatInt Balanced Truncation Matrix Equations

#### **Motivating Examples**

Parametric MOR: Applications in Microsystems/MEMS Design

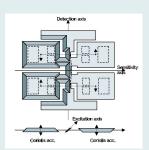
# Microgyroscope (butterfly gyro)



- Voltage applied to electrodes induces vibration of wings, resulting rotation due to Coriolis force yields sensor data.
- FE model of second order:  $N = 17.361 \rightsquigarrow n = 34.722, m = 1, q = 12.$
- Sensor for position control based on acceleration and rotation.

Source: The Oberwolfach Benchmark Collection http://www.imtek.de/simulation/benchmark

Application: inertial navigation.



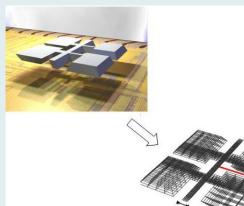
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### **Motivating Examples**

Parametric MOR: Applications in Microsystems/MEMS Design

## Microgyroscope (butterfly gyro)

Parametric FE model:  $M(d)\ddot{x}(t) + D(\Phi, d, \alpha, \beta)\dot{x}(t) + T(d)x(t) = Bu(t)$ .



### **Motivating Examples**

Parametric MOR: Applications in Microsystems/MEMS Design

## Microgyroscope (butterfly gyro)

Parametric FF model:

$$M(d)\ddot{x}(t) + D(\Phi, d, \alpha, \beta)\dot{x}(t) + T(d)x(t) = Bu(t),$$

wohei

$$M(d) = M_1 + dM_2,$$
  
 $D(\Phi, d, \alpha, \beta) = \Phi(D_1 + dD_2) + \alpha M(d) + \beta T(d),$   
 $T(d) = T_1 + \frac{1}{d}T_2 + dT_3,$ 



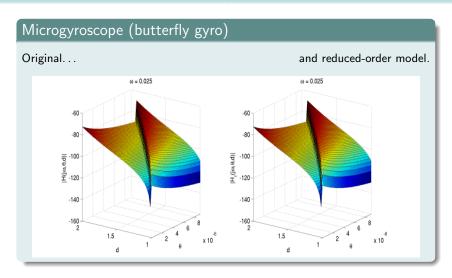
with

- width of bearing: d,
- angular velocity: Φ,
- Rayleigh damping parameters:  $\alpha, \beta$ .

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### **Motivating Examples**

Parametric MOR: Applications in Microsystems/MEMS Design



Mathematical Basics MOR by Projection RatInt Balanced Truncation Matrix Equations

## **Outline**

- 1 Introduction
- Mathematical Basics
  - Numerical Linear Algebra
  - Systems and Control Theory
  - Qualitative and Quantitative Study of the Approximation Error
- Model Reduction by Projection
- 4 Interpolatory Model Reduction
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## **Numerical Linear Algebra**

Image Compression by Truncated SVD

- A digital image with  $n_x \times n_y$  pixels can be represented as matrix  $X \in \mathbb{R}^{n_x \times n_y}$ , where  $x_{ij}$  contains color information of pixel (i,j).
- Memory (in single precision):  $4 \cdot n_x \cdot n_y$  bytes.

## $\mathsf{Theorem}\;(\mathsf{Schmidt} ext{-Mirsky/Eckart-Young})$

Best rank-r approximation to  $X \in \mathbb{R}^{n_x \times n_y}$  w.r.t. spectral norm:

$$\widehat{X} = \sum\nolimits_{j=1}^r \sigma_j u_j v_j^T,$$

where  $X = U\Sigma V^T$  is the singular value decomposition (SVD) of X. The approximation error is  $||X - \widehat{X}||_2 = \sigma_{r+1}$ .

### Idea for dimension reduction

Instead of X save  $u_1, \ldots, u_r, \sigma_1 v_1, \ldots, \sigma_r v_r$ .  $\rightarrow$  memory =  $4r \times (n_x + n_y)$  bytes.

## **Numerical Linear Algebra**

Image Compression by Truncated SVD

- A digital image with  $n_x \times n_y$  pixels can be represented as matrix  $X \in \mathbb{R}^{n_x \times n_y}$ , where  $x_{ij}$  contains color information of pixel (i, j).
- Memory (in single precision):  $4 \cdot n_x \cdot n_y$  bytes.

## Theorem (Schmidt-Mirsky/Eckart-Young)

Best rank-r approximation to  $X \in \mathbb{R}^{n_x \times n_y}$  w.r.t. spectral norm:

$$\widehat{X} = \sum_{j=1}^{r} \sigma_{j} u_{j} v_{j}^{T},$$

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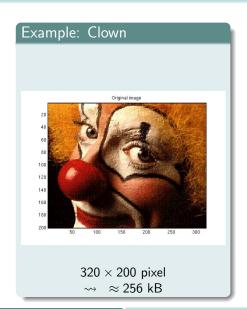
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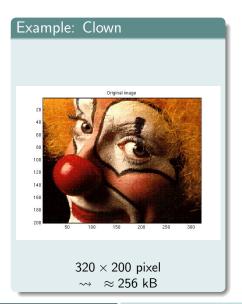
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## **Example: Image Compression by Truncated SVD**

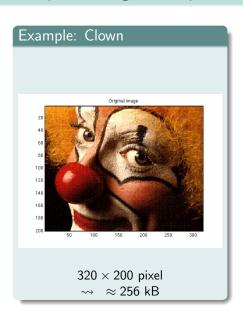


## **Example: Image Compression by Truncated SVD**



• rank r = 50,  $\approx 104$  kB





• rank r = 50,  $\approx 104$  kB



• rank r = 20,  $\approx 42$  kB

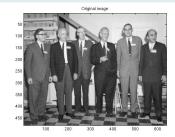


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## **Dimension Reduction via SVD**

## Example: Gatlinburg

Organizing committee Gatlinburg/Householder Meeting 1964: James H. Wilkinson, Wallace Givens, George Forsythe, Alston Householder, Peter Henrici, Fritz L. Bauer.

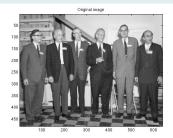


 $640 \times 480$  pixel,  $\approx 1229$  kB

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 $640 \times 480$  pixel,  $\approx 1229$  kB

### • rank r = 100, $\approx 448$ kB

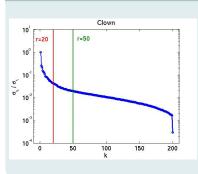


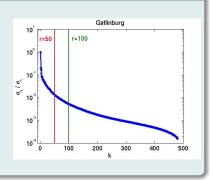
### • rank r = 50, $\approx 224$ kB



Image data compression via SVD works, if the singular values decay (exponentially).







The Laplace transform

### Definition

The Laplace transform of a time domain function  $f \in L_{1,loc}$  with  $dom(f) = \mathbb{R}_0^+$  is

$$\mathcal{L}:f(t)\mapsto f(s):=\mathcal{L}\{f(t)\}(s):=\int_0^\infty \mathrm{e}^{-st}f(t)\,dt,\quad s\in\mathbb{C}.$$

F is a function in the (Laplace or) frequency domain.

**Note:** for frequency domain evaluations ("frequency response analysis"), one takes re s=0 and im  $s\geq 0$ . Then  $\omega:=\operatorname{im} s$  takes the role of a frequency (in [rad/s], i.e.,  $\omega=2\pi v$  with v measured in [Hz]).

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$$\mathcal{L}\{\dot{f}(t)\}(s) = sF(s).$$

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Note: for ease of notation, in the following we will use lower-case letters for both, a function and its Laplace transform!

Mathematical Basics MOR by Projection RatInt Balanced Truncation Matrix Equations

### **Systems and Control Theory**

The Model Reduction Problem as Approximation Problem in Frequency Domain

## Linear Systems in Frequency Domain

Application of Laplace transform  $(x(t)\mapsto x(s),\,\dot{x}(t)\mapsto sx(s))$  to linear system

$$E\dot{x}(t) = Ax(t) + Bu(t), \quad y(t) = Cx(t) + Du(t)$$

with x(0) = 0 yields:

$$sEx(s) = Ax(s) + Bu(s), \quad y(s) = Cx(s) + Du(s),$$

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⇒ I/O-relation in frequency domain:

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**Goal:** Fast evaluation of mapping  $u \rightarrow y$ .

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Example.

The Model Reduction Problem as Approximation Problem in Frequency Domain

## Formulating model reduction in frequency domain

Approximate the dynamical system

$$\begin{array}{lll} E\dot{x} &=& Ax+Bu, & E,A\in\mathbb{R}^{n\times n},\ B\in\mathbb{R}^{n\times m},\\ y &=& Cx+Du, & C\in\mathbb{R}^{q\times n},\ D\in\mathbb{R}^{q\times m}, \end{array}$$

by reduced-order system

$$\begin{array}{lll} \hat{E}\dot{\hat{x}} & = & \hat{A}\hat{x} + \hat{B}u, & \hat{E}, \hat{A} \in \mathbb{R}^{r \times r}, \ \hat{B} \in \mathbb{R}^{r \times m}, \\ \hat{y} & = & \hat{C}\hat{x} + \hat{D}u, & \hat{C} \in \mathbb{R}^{q \times r}, \ \hat{D} \in \mathbb{R}^{q \times m} \end{array}$$

of order  $r \ll n$ , such that

$$||y - \hat{y}|| = ||Gu - \hat{G}u|| \le ||G - \hat{G}|| \cdot ||u|| < \text{tolerance} \cdot ||u||.$$

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 $\implies$  Approximation problem:  $\min_{\text{order } (\hat{G}) < r} ||G - \hat{G}||$ .

# Systems and Control Theory Properties of linear systems

### Definition

A linear system

$$E\dot{x}(t) = Ax(t) + Bu(t), \quad y(t) = Cx(t) + Du(t)$$

is stable if its transfer function G(s) has all its poles in the left half plane and it is asymptotically (or Lyapunov or exponentially) stable if all poles are in the open left half plane  $\mathbb{C}^- := \{z \in \mathbb{C} \mid \Re(z) < 0\}$ .

#### Lemma

Sufficient for asymptotic stability is that A is asymptotically stable (or Hurwitz), i.e., the spectrum of  $A - \lambda E$ , denoted by  $\Lambda(A, E)$ , satisfies  $\Lambda(A, E) \subset \mathbb{C}^-$ .

Note that by abuse of notation, often *stable system* is used for asymptotically stable systems.

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Mathematical Basics MOR by Projection RatInt Balanced Truncation Matrix Equations

## Systems and Control Theory

Properties of linear systems

Further properties to be discussed:

- Controllability/reachability
- Observability
- Stabilizability
- Detectability

See handout "Mathematical Basics".

Realizations of Linear Systems (with  $E = I_n$  for simplicity)

### Definition

For a linear (time-invariant) system

$$\Sigma: \begin{cases} \dot{x}(t) = Ax(t) + Bu(t), & \text{with transfer function} \\ y(t) = Cx(t) + Du(t), & G(s) = C(sI - A)^{-1}B + D, \end{cases}$$

the quadruple  $(A, B, C, D) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m} \times \mathbb{R}^{q \times n} \times \mathbb{R}^{q \times m}$  is called a realization of  $\Sigma$ .

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### Realizations are not unique!

Transfer function is invariant under state-space transformations,

$$\mathcal{T}: \left\{ \begin{array}{ccc} x & \rightarrow & Tx, \\ (A,B,C,D) & \rightarrow & (TAT^{-1},TB,CT^{-1},D), \end{array} \right.$$

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### Realizations are not unique!

Transfer function is invariant under addition of uncontrollable/unobservable states:

$$\frac{d}{dt} \begin{bmatrix} x \\ x_1 \end{bmatrix} = \begin{bmatrix} A & 0 \\ 0 & A_1 \end{bmatrix} \begin{bmatrix} x \\ x_1 \end{bmatrix} + \begin{bmatrix} B \\ B_1 \end{bmatrix} u(t), \quad y(t) = \begin{bmatrix} C & 0 \end{bmatrix} \begin{bmatrix} x \\ x_1 \end{bmatrix} + Du(t),$$

$$\frac{d}{dt} \begin{bmatrix} x \\ x_2 \end{bmatrix} = \begin{bmatrix} A & 0 \\ 0 & A_2 \end{bmatrix} \begin{bmatrix} x \\ x_2 \end{bmatrix} + \begin{bmatrix} B \\ 0 \end{bmatrix} u(t), \quad y(t) = \begin{bmatrix} C & C_2 \end{bmatrix} \begin{bmatrix} x \\ x_2 \end{bmatrix} + Du(t),$$

for arbitrary  $A_i \in \mathbb{R}^{n_j \times n_j}$ , j = 1, 2,  $B_1 \in \mathbb{R}^{n_1 \times m}$ ,  $C_2 \in \mathbb{R}^{q \times n_2}$  and any  $n_1, n_2 \in \mathbb{N}$ .

Realizations of Linear Systems (with  $E = I_n$  for simplicity)

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## Realizations are not unique!

Hence.

$$(A, B, C, D), \qquad \left( \begin{bmatrix} A & 0 \\ 0 & A_1 \end{bmatrix}, \begin{bmatrix} B \\ B_1 \end{bmatrix}, \begin{bmatrix} C & 0 \end{bmatrix}, D \right),$$

$$(TAT^{-1}, TB, CT^{-1}, D), \qquad \left( \begin{bmatrix} A & 0 \\ 0 & A_2 \end{bmatrix}, \begin{bmatrix} B \\ 0 \end{bmatrix}, \begin{bmatrix} C & C_2 \end{bmatrix}, D \right),$$

are all realizations of  $\Sigma$ !

Realizations of Linear Systems (with  $E = I_n$  for simplicity)

### **Definition**

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### Definition

The McMillan degree of  $\Sigma$  is the unique minimal number  $\hat{n} \geq 0$  of states necessary to describe the input-output behavior completely.

A minimal realization is a realization  $(\hat{A}, \hat{B}, \hat{C}, \hat{D})$  of  $\Sigma$  with order  $\hat{n}$ .

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### **Theorem**

A realization (A, B, C, D) of a linear system is minimal  $\iff$  (A, B) is controllable and (A, C) is observable.

**Balanced Realizations** 

### Definition

A realization (A, B, C, D) of a linear system  $\Sigma$  is balanced if its infinite controllability/observability Gramians P/Q satisfy

$$P = Q = \operatorname{diag} \{\sigma_1, \dots, \sigma_n\}$$
 (w.l.o.g.  $\sigma_j \ge \sigma_{j+1}, \ j = 1, \dots, n-1$ ).

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When does a balanced realization exist?

# Systems and Control Theory Balanced Realizations

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When does a balanced realization exist? Assume A to be Hurwitz, i.e.  $\Lambda(A) \subset \mathbb{C}^-$ . Then:

### **Theorem**

Given a stable minimal linear system  $\Sigma : (A, B, C, D)$ , a balanced realization is obtained by the state-space transformation with

$$T_b := \Sigma^{-\frac{1}{2}} V^T R,$$

where  $P = S^T S$ ,  $Q = R^T R$  (e.g., Cholesky decompositions) and  $SR^T = U\Sigma V^T$  is the SVD of  $SR^T$ .

Proof. Exercise!

#### **Balanced Realizations**

### Definition

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 $\sigma_1, \ldots, \sigma_n$  are the Hankel singular values of  $\Sigma$ .

**Note:**  $\sigma_1, \ldots, \sigma_n \geq 0$  as  $P, Q \geq 0$  by definition, and  $\sigma_1, \ldots, \sigma_n > 0$  in case of minimality!

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### **Theorem**

The infinite controllability/observability Gramians P/Q satisfy the Lyapunov equations

$$AP + PA^{T} + BB^{T} = 0$$
,  $A^{T}Q + QA + C^{T}C = 0$ .

# Systems and Control Theory Balanced Realizations

# Theorem

The infinite controllability/observability Gramians P/Q satisfy the Lyapunov equations

$$AP + PA^T + BB^T = 0$$
,  $A^TQ + QA + C^TC = 0$ .

Proof. (For controllability Gramian only, observability case is analogous!)

$$AP + PA^{T} + BB^{T} = A \int_{0}^{\infty} e^{At}BB^{T}e^{A^{T}t}dt + \int_{0}^{\infty} e^{At}BB^{T}e^{A^{T}t}dt A^{T} + BB^{T}$$

$$= \int_{0}^{\infty} \underbrace{Ae^{At}BB^{T}e^{A^{T}t} + e^{At}BB^{T}e^{A^{T}t}A^{T}}_{=\frac{d}{dt}e^{At}BB^{T}e^{A^{T}t}} dt + BB^{T}$$

$$= \underbrace{\lim_{t \to \infty} e^{At}BB^{T}e^{A^{T}t}}_{=0} - \underbrace{e^{A\cdot 0}BB^{T}e^{A^{T}\cdot 0}}_{=I_{n}} + BB^{T}$$

$$= 0.$$

### Definition

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### Theorem

The Hankel singular values (HSVs) of a stable minimal linear system are system invariants, i.e. they are unaltered by state-space transformations!

# Systems and Control Theory Balanced Realizations

### Theorem

The Hankel singular values (HSVs) of a stable minimal linear system are system invariants, i.e. they are unaltered by state-space transformations!

**Proof.** In balanced coordinates, the HSVs are  $\Lambda(PQ)^{\frac{1}{2}}$ . Now let

$$(\hat{A}, \hat{B}, \hat{C}, D) = (TAT^{-1}, TB, CT^{-1}, D)$$

be any transformed realization with associated controllability Lyapunov equation

$$0 = \hat{A}\hat{P} + \hat{P}\hat{A}^{T} + \hat{B}\hat{B}^{T} = TAT^{-1}\hat{P} + \hat{P}T^{-T}A^{T}T^{T} + TBB^{T}T^{T}.$$

This is equivalent to

$$0 = A(T^{-1}\hat{P}T^{-T}) + (T^{-1}\hat{P}T^{-T})A^{T} + BB^{T}.$$

The uniqueness of the solution of the Lyapunov equation implies that  $\hat{P} = TPT^T$  and, analogously,  $\hat{Q} = T^{-T}QT^{-1}$ . Therefore,

$$\hat{P}\hat{Q} = TPQT^{-1},$$

showing that  $\Lambda(\hat{P}\hat{Q}) = \Lambda(PQ) = \{\sigma_1^2, \dots, \sigma_n^2\}.$ 

# **Systems and Control Theory**

#### **Balanced Realizations**

### Definition

A realization (A, B, C, D) of a stable linear system  $\Sigma$  is balanced if its infinite controllability/observability Gramians P/Q satisfy

$$P = Q = \operatorname{diag} \{\sigma_1, \dots, \sigma_n\}$$
 (w.l.o.g.  $\sigma_j \ge \sigma_{j+1}, \ j = 1, \dots, n-1$ ).

 $\sigma_1, \ldots, \sigma_n$  are the Hankel singular values of  $\Sigma$ .

**Note:**  $\sigma_1, \ldots, \sigma_n \geq 0$  as  $P, Q \geq 0$  by definition, and  $\sigma_1, \ldots, \sigma_n > 0$  in case of minimality!

### Remark

For non-minimal systems, the Gramians can also be transformed into diagonal matrices with the leading  $\hat{n} \times \hat{n}$  submatrices equal to  $\operatorname{diag}(\sigma_1, \dots, \sigma_{\hat{n}})$ , and

$$\hat{P}\hat{Q} = \operatorname{diag}(\sigma_1^2, \ldots, \sigma_{\hat{n}}^2, 0, \ldots, 0).$$

see [Laub/Heath/Paige/Ward 1987, Tombs/Postlethwaite 1987].

Consider transfer function

$$G(s) = C(sI - A)^{-1}B + D$$

and input functions  $u \in \mathcal{L}_2^m \cong L_2^m(-\infty,\infty)$ , with the  $L_2$ -norm

$$||u||_2^2 := \frac{1}{2\pi} \int_{-\infty}^{\infty} u(j\omega)^H u(j\omega) d\omega.$$

Assume A (asympotically) stable:  $\Lambda(A) \subset \mathbb{C}^- := \{z \in \mathbb{C} : \operatorname{re} z < 0\}.$ 

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$$\int_{-\infty}^{\infty} y(\jmath\omega)^{H} y(\jmath\omega) d\omega = \int_{-\infty}^{\infty} u(\jmath\omega)^{H} G(\jmath\omega)^{H} G(\jmath\omega) u(\jmath\omega) d\omega$$

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$$\implies y \in \mathcal{L}_2^q \cong \mathcal{L}_2^q(-\infty,\infty).$$

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$$\|G\|_{\infty} := \sup_{\|u\|_2 \neq 0} \frac{\|Gu\|_2}{\|u\|_2}$$

is well defined. It can be shown that

$$\|G\|_{\infty} = \sup_{\omega \in \mathbb{R}} \|G(\jmath \omega)\| = \sup_{\omega \in \mathbb{R}} \sigma_{max} (G(\jmath \omega)).$$

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Sketch of proof:

$$||G(\jmath\omega)u(\jmath\omega)|| \le ||G(\jmath\omega)||||u(\jmath\omega)|| \Rightarrow "\le ".$$
  
Construct  $u$  with  $||Gu||_2 = \sup_{\omega \in \mathbb{R}} ||G(\jmath\omega)|||u||_2.$ 

Consider transfer function

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## Hardy space $\mathcal{H}_{\infty}$

Function space of matrix-/scalar-valued functions that are analytic and bounded in  $\mathbb{C}^+$ .

The  $\mathcal{H}_{\infty}$ -norm is

$$\|F\|_{\infty} := \sup_{\mathsf{re}\,\mathsf{s}>0} \sigma_{\mathsf{max}}\left(F(\mathsf{s})\right) = \sup_{\omega\in\mathbb{R}} \sigma_{\mathsf{max}}\left(F(\jmath\omega)\right).$$

Stable transfer functions are in the Hardy spaces

- ullet  $\mathcal{H}_{\infty}$  in the SISO case (single-input, single-output, m=q=1);
- $\mathcal{H}_{\infty}^{q \times m}$  in the MIMO case (multi-input, multi-output, m > 1, q > 1).

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Paley-Wiener Theorem (Parseval's equation/Plancherel Theorem)

$$L_2(-\infty,\infty)\cong \mathcal{L}_2, \quad L_2(0,\infty)\cong \mathcal{H}_2$$

Consequently, 2-norms in time and frequency domains coincide!

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### $\mathcal{H}_{\infty}$ approximation error

Reduced-order model  $\Rightarrow$  transfer function  $\hat{G}(s) = \hat{C}(sl_r - \hat{A})^{-1}\hat{B} + \hat{D}$ .

$$||y - \hat{y}||_2 = ||Gu - \hat{G}u||_2 \le ||G - \hat{G}||_{\infty} ||u||_2.$$

 $\implies$  compute reduced-order model such that  $\|G - \hat{G}\|_{\infty} < tol!$ 

Note: error bound holds in time- and frequency domain due to Paley-Wiener!

Consider stable transfer function

$$G(s) = C(sI - A)^{-1}B$$
, i.e.  $D = 0$ .

## Hardy space $\mathcal{H}_2$

Function space of matrix-/scalar-valued functions that are analytic  $\mathbb{C}^+$  and bounded w.r.t. the  $\mathcal{H}_2$ -norm

$$||F||_2 := \frac{1}{2\pi} \left( \sup_{\mathsf{re}\,\sigma > 0} \int_{-\infty}^{\infty} ||F(\sigma + \jmath\omega)||_F^2 d\omega \right)^{\frac{1}{2}}$$
$$= \frac{1}{2\pi} \left( \int_{-\infty}^{\infty} ||F(\jmath\omega)||_F^2 d\omega \right)^{\frac{1}{2}}.$$

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# $\mathcal{H}_2$ approximation error for impulse response $(u(t)=u_0\delta(t))$

Reduced-order model  $\Rightarrow$  transfer function  $\hat{G}(s) = \hat{C}(sI_r - \hat{A})^{-1}\hat{B}$ .

$$||y - \hat{y}||_2 = ||Gu_0\delta - \hat{G}u_0\delta||_2 \le ||G - \hat{G}||_2||u_0||.$$

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# Theorem (Practical Computation of the $\mathcal{H}_2$ -norm)

$$||F||_2^2 = \operatorname{tr}\left(B^T Q B\right) = \operatorname{tr}\left(C P C^T\right),$$

where P,Q are the controllability and observability Gramians of the corresponding LTI system.

### Qualitative and Quantitative Study of the Approximation Error Approximation Problems

### Output errors in time-domain

$$||y - \hat{y}||_2 \le ||G - \hat{G}||_{\infty} ||u||_2 \Longrightarrow ||G - \hat{G}||_{\infty} < \text{tol}$$
$$||y - \hat{y}||_{\infty} \le ||G - \hat{G}||_2 ||u||_2 \Longrightarrow ||G - \hat{G}||_2 < \text{tol}$$

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$$\begin{aligned} \|y - \hat{y}\|_2 & \leq \|G - \hat{G}\|_{\infty} \|u\|_2 & \Longrightarrow \|G - \hat{G}\|_{\infty} < \text{tol} \\ \|y - \hat{y}\|_{\infty} & \leq \|G - \hat{G}\|_2 \|u\|_2 & \Longrightarrow \|G - \hat{G}\|_2 < \text{tol} \end{aligned}$$

| $\mathcal{H}_{\infty}$ -norm | best approximation problem for given reduced order $r$ in general open; balanced truncation yields suboptimal solution with computable $\mathcal{H}_{\infty}$ -norm bound. |
|------------------------------|--|
| $\mathcal{H}_2$ -norm        | necessary conditions for best approximation known; (local) optimizer computable with iterative rational Krylov algorithm (IRKA)  |
| Hankel-norm                  | optimal Hankel norm approximation (AAK theory).  |
| $\ G\ _H := \sigma_{max}$    |  |

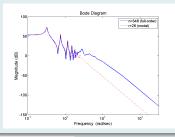
# Qualitative and Quantitative Study of the Approximation Error Computable error measures

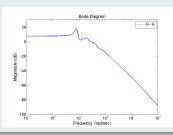
Evaluating system norms is computationally very (sometimes too) expensive.

### Other measures

- absolute errors  $\|G(\jmath\omega_j) \hat{G}(\jmath\omega_j)\|_2$ ,  $\|G(\jmath\omega_j) \hat{G}(\jmath\omega_j)\|_\infty$   $(j=1,\ldots,N_\omega)$ ;
- relative errors  $\frac{\|G(\jmath\omega_j) \hat{G}(\jmath\omega_j)\|_2}{\|G(\jmath\omega_i)\|_2}$ ,  $\frac{\|G(\jmath\omega_j) \hat{G}(\jmath\omega_j)\|_{\infty}}{\|G(\jmath\omega_i)\|_{\infty}}$ ;
- "eyeball norm", i.e. look at frequency response/Bode (magnitude) plot: for SISO system, log-log plot frequency vs.  $|G(\jmath\omega)|$  (or  $|G(\jmath\omega)-\hat{G}(\jmath\omega)|$ ) in decibels, 1 dB  $\simeq 20\log_{10}(\text{value})$ .

For MIMO systems,  $q \times m$  array of plots  $G_{ij}$ .





# **Outline**

- Introduction
- 2 Mathematical Basics
- Model Reduction by Projection
  - Projection and Interpolation
  - Modal Truncation
- 4 Interpolatory Model Reduction
- Balanced Truncation
- 6 Solving Large-Scale Matrix Equations
- Final Remarks

- Automatic generation of compact models.
- Satisfy desired error tolerance for all admissible input signals, i.e.,

$$||y - \hat{y}|| < \text{tolerance} \cdot ||u|| \qquad \forall u \in L_2(\mathbb{R}, \mathbb{R}^m).$$

- ⇒ Need computable error bound/estimate!
- Preserve physical properties:

$$\int_{-\infty}^{t} u(\tau)^{T} y(\tau) d\tau \ge 0 \quad \forall t \in \mathbb{R}, \quad \forall u \in L_{2}(\mathbb{R}, \mathbb{R}^{m})$$

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# Goals

- Automatic generation of compact models.
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**Projection Basics** 

# Definition 3.1 (Projector)

A projector is a matrix  $P \in \mathbb{R}^{n \times n}$  with  $P^2 = P$ . Let  $\mathcal{V} = \operatorname{range}(P)$ , then P is projector onto  $\mathcal{V}$ . On the other hand, if  $\{v_1, \ldots, v_r\}$  is a basis of  $\mathcal{V}$  and  $V = [v_1, \ldots, v_r]$ , then  $P = V(V^TV)^{-1}V^T$  is a projector onto  $\mathcal{V}$ .

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- If  $P = P^T$ , then P is an orthogonal projector (aka: Galerkin projection), otherwise an oblique projector (aka: Petrov-Galerkin projection).
- P is the identity operator on  $\mathcal{V}$ , i.e.,  $Pv = v \ \forall v \in \mathcal{V}$ .
- $\bullet$  I-P is the complementary projector onto ker P.
- If  $\mathcal V$  is an A-invariant subspace corresponding to a subset of A's spectrum, then we call P a spectral projector.
- Let  $W \subset \mathbb{R}^n$  be another r-dimensional subspace and  $W = [w_1, \ldots, w_r]$  be a basis matrix for W, then  $P = V(W^T V)^{-1} W^T$  is an oblique projector onto V along W.

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#### Methods:

- Modal Truncation
- Rational Interpolation (Padé-Approximation and (rational) Krylov Subspace Methods)
- Balanced Truncation
- many more...

Joint feature of these methods: computation of reduced-order model (ROM) by projection!

# Projection and Interpolation

Joint feature of these methods: computation of reduced-order model (ROM) by projection!

Assume trajectory x(t;u) is contained in low-dimensional subspace  $\mathcal V$ . Thus, use Galerkin or Petrov-Galerkin-type projection of state-space onto  $\mathcal V$  along complementary subspace  $\mathcal W\colon x\approx VW^Tx=:\tilde x$ , where

range 
$$(V) = V$$
, range  $(W) = W$ ,  $W^T V = I_r$ .

Then, with  $\hat{x} = W^T x$ , we obtain  $x \approx V \hat{x}$  so that

$$||x - \tilde{x}|| = ||x - V\hat{x}||,$$

and the reduced-order model is

$$\hat{A} := W^T A V$$
,  $\hat{B} := W^T B$ ,  $\hat{C} := C V$ ,  $(\hat{D} := D)$ .

### Projection and Interpolation

Joint feature of these methods: computation of reduced-order model (ROM) by projection! Assume trajectory x(t;u) is contained in low-dimensional subspace  $\mathcal V$ . Thus, use Galerkin or Petrov-Galerkin-type projection of state-space onto  $\mathcal V$  along complementary subspace  $\mathcal W\colon x\approx VW^Tx=:\tilde x$ , and the reduced-order model is  $\hat x=W^Tx$ 

$$\hat{A} := W^T A V, \quad \hat{B} := W^T B, \quad \hat{C} := C V, \quad (\hat{D} := D).$$

Important observation:

• The state equation residual satisfies  $\dot{\tilde{x}} - A\tilde{x} - Bu \perp \mathcal{W}$ , since

$$W^{T} \left( \dot{\tilde{x}} - A\tilde{x} - Bu \right) = W^{T} \left( VW^{T} \dot{x} - AVW^{T} x - Bu \right)$$

Joint feature of these methods: computation of reduced-order model (ROM) by projection! Assume trajectory x(t;u) is contained in low-dimensional subspace  $\mathcal V$ . Thus, use Galerkin or Petrov-Galerkin-type projection of state-space onto  $\mathcal V$  along complementary subspace  $\mathcal W\colon x\approx VW^Tx=:\tilde x$ , and the reduced-order model is  $\hat x=W^Tx$ 

$$\hat{A} := W^{T}AV, \quad \hat{B} := W^{T}B, \quad \hat{C} := CV, \quad (\hat{D} := D).$$

Important observation:

• The state equation residual satisfies  $\dot{\tilde{x}} - A\tilde{x} - Bu \perp \mathcal{W}$ , since

$$W^{T} \left( \dot{\tilde{x}} - A\tilde{x} - Bu \right) = W^{T} \left( VW^{T} \dot{x} - AVW^{T} x - Bu \right)$$
$$= \underbrace{W^{T} \dot{x}}_{\dot{\hat{x}}} - \underbrace{W^{T} AV}_{=\hat{A}} \underbrace{W^{T} x}_{=\hat{x}} - \underbrace{W^{T} B}_{=\hat{B}} u$$

#### Model Reduction by Projection **Projection and Interpolation**

Joint feature of these methods:

computation of reduced-order model (ROM) by projection! Assume trajectory x(t; u) is contained in low-dimensional subspace  $\mathcal{V}$ . Thus, use Galerkin or Petrov-Galerkin-type projection of state-space onto  $\mathcal V$  along complementary subspace  $W: x \approx VW^Tx =: \tilde{x}$ , and the reduced-order model is  $\hat{x} = W^T x$ 

$$\hat{A} := W^T A V, \quad \hat{B} := W^T B, \quad \hat{C} := C V, \quad (\hat{D} := D).$$

Important observation:

• The state equation residual satisfies  $\dot{\tilde{x}} - A\tilde{x} - Bu \perp \mathcal{W}$ , since

$$W^{T} \left( \dot{\tilde{x}} - A\tilde{x} - Bu \right) = W^{T} \left( VW^{T} \dot{x} - AVW^{T} x - Bu \right)$$

$$= \underbrace{W^{T} \dot{x}}_{\dot{\hat{x}}} - \underbrace{W^{T} AV}_{=\hat{A}} \underbrace{W^{T} x}_{=\hat{x}} - \underbrace{W^{T} B}_{=\hat{B}} u$$

$$= \dot{\hat{x}} - \hat{A}\hat{x} - \hat{B}u = 0.$$

Projection and Interpolation

### Projection → Rational Interpolation

Given the ROM

$$\hat{A} = W^T A V$$
,  $\hat{B} = W^T B$ ,  $\hat{C} = C V$ ,  $(\hat{D} = D)$ ,

the error transfer function can be written as

$$G(s) - \hat{G}(s) = (C(sI_n - A)^{-1}B + D) - (\hat{C}(sI_r - \hat{A})^{-1}\hat{B} + \hat{D})$$

Projection and Interpolation

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$$= C\left( (sI_n - A)^{-1} - V(sI_r - \hat{A})^{-1}W^T \right) B$$

Projection and Interpolation

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$$= C\left((sI_n - A)^{-1} - V(sI_r - \hat{A})^{-1}W^T\right)B$$

$$= C\left(I_n - \underbrace{V(sI_r - \hat{A})^{-1}W^T(sI_n - A)}_{=:P(s)}\right)(sI_n - A)^{-1}B.$$

Projection and Interpolation

### Projection → Rational Interpolation

Given the ROM

$$\hat{A} = W^T A V$$
,  $\hat{B} = W^T B$ ,  $\hat{C} = C V$ ,  $(\hat{D} = D)$ ,

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=  $C(I_n - \underbrace{V(sI_r - \hat{A})^{-1}W^T(sI_n - A)}_{=:P(s)})(sI_n - A)^{-1}B.$ 

If  $s_* \in \mathbb{C} \setminus (\Lambda(A) \cup \Lambda(\hat{A}))$ , then  $P(s_*)$  is a projector onto  $\mathcal{V}$ :

range  $(P(s_*)) \subset \text{range}(V)$ , all matrices have full rank  $\Rightarrow$  "=",

$$P(s_*)^2 = V(s_*I_r - \hat{A})^{-1}W^T(s_*I_n - A)V(s_*I_r - \hat{A})^{-1}W^T(s_*I_n - A)$$

#### Projection and Interpolation

### Projection → Rational Interpolation

Given the ROM

$$\hat{A} = W^T A V$$
,  $\hat{B} = W^T B$ ,  $\hat{C} = C V$ ,  $(\hat{D} = D)$ ,

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$$G(s) - \hat{G}(s) = \left(C(sI_n - A)^{-1}B + D\right) - \left(\hat{C}(sI_r - \hat{A})^{-1}\hat{B} + \hat{D}\right)$$

$$= C\left(I_n - \underbrace{V(sI_r - \hat{A})^{-1}W^T(sI_n - A)}_{=:P(s)}\right)(sI_n - A)^{-1}B.$$

If  $s_* \in \mathbb{C} \setminus (\Lambda(A) \cup \Lambda(\hat{A}))$ , then  $P(s_*)$  is a projector onto  $\mathcal{V}$ :

 $\operatorname{range}(P(s_*)) \subset \operatorname{range}(V)$ , all matrices have full rank  $\Rightarrow$  "=",

$$P(s_*)^2 = V(s_*I_r - \hat{A})^{-1}W^T(s_*I_n - A)V(s_*I_r - \hat{A})^{-1}W^T(s_*I_n - A)$$

$$= V(s_*I_r - \hat{A})^{-1}\underbrace{(s_*I_r - \hat{A})(s_*I_r - \hat{A})^{-1}}_{-I_*}W^T(s_*I_n - A) = P(s_*).$$

#### Model Reduction by Projection Projection and Interpolation

### Projection → Rational Interpolation

Given the ROM

$$\hat{A} = W^T A V$$
,  $\hat{B} = W^T B$ ,  $\hat{C} = C V$ ,  $(\hat{D} = D)$ ,

the error transfer function can be written as

$$G(s) - \hat{G}(s) = \left(C(sI_n - A)^{-1}B + D\right) - \left(\hat{C}(sI_r - \hat{A})^{-1}\hat{B} + \hat{D}\right)$$
  
=  $C(I_n - \underbrace{V(sI_r - \hat{A})^{-1}W^T(sI_n - A)}_{=:P(s)})(sI_n - A)^{-1}B.$ 

If  $s_* \in \mathbb{C} \setminus (\Lambda(A) \cup \Lambda(\hat{A}))$ , then  $P(s_*)$  is a projector onto  $\mathcal{V} \Longrightarrow$ 

if 
$$(s_*I_n - A)^{-1}B \in \mathcal{V}$$
, then  $(I_n - P(s_*))(s_*I_n - A)^{-1}B = 0$ .

hence

$$G(s_*) - \hat{G}(s_*) = 0 \implies G(s_*) = \hat{G}(s_*)$$
, i.e.,  $\hat{G}$  interpolates  $G$  in  $s_*$ !

Projection and Interpolation

### Projection → Rational Interpolation

Given the ROM

$$\hat{A} = W^T A V$$
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the error transfer function can be written as

$$G(s) - \hat{G}(s) = \left(C(sI_n - A)^{-1}B + D\right) - \left(\hat{C}(sI_r - \hat{A})^{-1}\hat{B} + \hat{D}\right)$$

$$= C\left(I_n - \underbrace{V(sI_r - \hat{A})^{-1}W^T(sI_n - A)}_{=:P(s)}\right)(sI_n - A)^{-1}B.$$

Analogously, = 
$$C(sI_n - A)^{-1}(I_n - \underbrace{(sI_n - A)V(sI_r - \hat{A})^{-1}W^T})B$$
.

If  $s_* \in \mathbb{C} \setminus (\Lambda(A) \cup \Lambda(\hat{A}))$ , then  $Q(s)^H$  is a projector onto  $\mathcal{W} \Longrightarrow$ 

if 
$$(s_*I_n - A)^{-*}C^T \in \mathcal{W}$$
, then  $C(s_*I_n - A)^{-1}(I_n - Q(s_*)) = 0$ ,

hence

$$G(s_*) - \hat{G}(s_*) = 0 \implies G(s_*) = \hat{G}(s_*)$$
, i.e.,  $\hat{G}$  interpolates  $G$  in  $s_*!$ 

### Model Reduction by Projection Projection and Interpolation

# Theorem 3.3

#### [Grimme '97, Villemagne/Skelton '87]

Given the ROM

$$\hat{A} = W^T A V$$
,  $\hat{B} = W^T B$ ,  $\hat{C} = C V$ ,  $(\hat{D} = D)$ ,

and  $s_* \in \mathbb{C} \setminus (\Lambda(A) \cup \Lambda(\hat{A}))$ , if either

• 
$$(s_* I_n - A)^{-1} B \in \text{range}(V)$$
, or

• 
$$(s_*I_n - A)^{-*}C^T \in \text{range}(W)$$
,

then the interpolation condition

$$G(s_*) = \hat{G}(s_*).$$

in s\* holds.

Note: extension to Hermite interpolation conditions later!

#### Basic method:

Assume A is diagonalizable,  $T^{-1}AT = D_A$ , project state-space onto A-invariant subspace  $\mathcal{V} = \operatorname{span}(t_1, \ldots, t_r)$ ,  $t_k = \operatorname{eigenvectors}$  corresp. to "dominant" modes / eigenvalues of A. Then with

$$V = T(:,1:r) = [t_1,\ldots,t_r], \quad \tilde{W}^H = T^{-1}(1:r,:), \quad W = \tilde{W}(V^H\tilde{W})^{-1},$$

reduced-order model is

$$\hat{A} := W^H A V = \operatorname{diag} \{\lambda_1, \dots, \lambda_r\}, \quad \hat{B} := W^H B, \quad \hat{C} = C V$$

Also computable by truncation:

$$T^{-1}AT = \begin{bmatrix} \hat{A} \\ A_2 \end{bmatrix}, \quad T^{-1}B = \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix}, \quad CT = [\hat{C}, C_2], \quad \hat{D} = D.$$

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Assume A is diagonalizable,  $T^{-1}AT = D_A$ , project state-space onto A-invariant subspace  $\mathcal{V} = \operatorname{span}(t_1, \ldots, t_r)$ ,  $t_k =$  eigenvectors corresp. to "dominant" modes / eigenvalues of A. Then with

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### Properties:

Simple computation for large-scale systems, using, e.g., Krylov subspace methods (Lanczos, Arnoldi), Jacobi-Davidson method.

#### Basic method:

$$T^{-1}AT = \begin{bmatrix} \hat{A} \\ A_2 \end{bmatrix}, \quad T^{-1}B = \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix}, \quad CT = [\hat{C}, C_2], \quad \hat{D} = D.$$

### Properties:

Error bound:

$$\|G - \hat{G}\|_{\infty} \le \|C_2\| \|B_2\| \frac{1}{\min_{\lambda \in \Lambda(A_2)} |\operatorname{Re}(\lambda)|}.$$

Proof.

$$G(s) = C(sI - A)^{-1}B + D = CTT^{-1}(sI - A)^{-1}TT^{-1}B + D$$

$$= CT(sI - T^{-1}AT)^{-1}T^{-1}B + D$$

$$= [\hat{C}, C_2] \begin{bmatrix} (sI_r - \hat{A})^{-1} \\ (sI_{n-r} - A_2)^{-1} \end{bmatrix} \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix} + D$$

$$= \hat{G}(s) + C_2(sI_{n-r} - A_2)^{-1}B_2,$$

#### Basic method:

$$T^{-1}AT = \begin{bmatrix} \hat{A} \\ A_2 \end{bmatrix}, \quad T^{-1}B = \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix}, \quad CT = [\hat{C}, C_2], \quad \hat{D} = D.$$

### Properties:

Error bound:

$$\|G-\hat{G}\|_{\infty} \leq \|C_2\|\|B_2\|\frac{1}{\min_{\lambda \in \Lambda(A_2)}|\operatorname{Re}(\lambda)|}.$$

Proof:

$$G(s) = \hat{G}(s) + C_2(sI_{n-r} - A_2)^{-1}B_2$$

observing that  $||G - \hat{G}||_{\infty} = \sup_{\omega \in \mathbb{R}} \sigma_{\max}(C_2(\gamma \omega I_{n-r} - A_2)^{-1}B_2)$ , and

$$C_2(\jmath\omega I_{n-r}-A_2)^{-1}B_2=C_2\mathrm{diag}\left(\frac{1}{\jmath\omega-\lambda_{r+1}},\ldots,\frac{1}{\jmath\omega-\lambda_n}\right)B_2.$$

#### Basic method:

Assume A is diagonalizable,  $T^{-1}AT = D_A$ , project state-space onto A-invariant subspace  $\mathcal{V} = \operatorname{span}(t_1, \ldots, t_r)$ ,  $t_k =$  eigenvectors corresp. to "dominant" modes / eigenvalues of A. Then reduced-order model is

$$\hat{A} := W^H A V = \operatorname{diag} \{\lambda_1, \dots, \lambda_r\}, \quad \hat{B} := W^H B, \quad \hat{C} = C V$$

Also computable by truncation:

$$T^{-1}AT = \begin{bmatrix} \hat{A} \\ A_2 \end{bmatrix}, \quad T^{-1}B = \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix}, \quad CT = [\hat{C}, C_2], \quad \hat{D} = D.$$

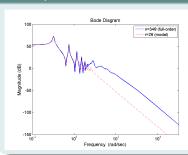
### Difficulties:

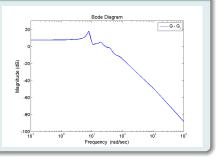
- Eigenvalues contain only limited system information.
- Dominance measures are difficult to compute.
   ([Litz '79] use Jordan canoncial form; otherwise merely heuristic criteria,
   e.g., [Varga '95]. Recent improvement: dominant pole algorithm.)
- Error bound not computable for really large-scale problems.

Example

**BEAM**, SISO system from SLICOT Benchmark Collection for Model Reduction, n = 348, m = q = 1, reduced using 13 dominant complex conjugate eigenpairs, error bound yields  $\|G - \hat{G}\|_{\infty} \le 1.21 \cdot 10^3$ 





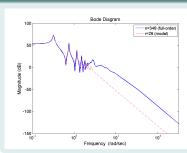


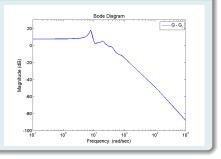
MATLAB® demo

Example

**BEAM**, SISO system from SLICOT Benchmark Collection for Model Reduction, n = 348, m = q = 1, reduced using 13 dominant complex conjugate eigenpairs, error bound yields  $||G - \hat{G}||_{\infty} \le 1.21 \cdot 10^3$ 







MATLAB® demo.

Extensions

#### Base enrichment

Static modes are defined by setting  $\dot{x}=0$  and assuming unit loads, i.e.,  $u(t)\equiv e_{j},\,j=1,\ldots,m$ :

$$0 = Ax(t) + Be_j \implies x(t) \equiv -A^{-1}b_j.$$

Projection subspace V is then augmented by  $A^{-1}[b_1,\ldots,b_m]=A^{-1}B$ .

Interpolation-projection framework  $\implies G(0) = \hat{G}(0)!$ 

If two sided projection is used, complimentary subspace can be augmented by  $A^{-T}C^T \Longrightarrow G'(0) = \hat{G}'(0)!$  (If  $m \neq q$ , add random vectors or delete some of the columns in  $A^{-T}C^T$ ).

Extensions

### Guyan reduction (static condensation)

Partition states in masters  $x_1 \in \mathbb{R}^r$  and slaves  $x_2 \in \mathbb{R}^{n-r}$  (FEM terminology) Assume stationarity, i.e.,  $\dot{x} = 0$  and solve for  $x_2$  in

$$0 = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} B_1 \\ B_2 \end{bmatrix} u$$

$$\Rightarrow x_2 = -A_{22}^{-1}A_{21}x_1 - A_{22}^{-1}B_2u.$$

Inserting this into the first part of the dynamic system

$$\dot{x}_1 = A_{11}x_1 + A_{12}x_2 + B_1u, \quad y = C_1x_1 + C_2x_2$$

then yields the reduced-order model

$$\dot{x}_1 = (A_{11} - A_{12}A_{22}^{-1}A_{21})x_1 + (B_1 - A_{12}A_{22}^{-1}B_2)u 
y = (C_1 - C_2A_{22}^{-1}A_{21})x_1 - C_2A_{22}^{-1}B_2u.$$

**Dominant Poles** 

#### Pole-Residue Form of Transfer Function

Consider partial fraction expansion of transfer function with D=0:

$$G(s) = \sum_{k=1}^{n} \frac{R_k}{s - \lambda_k}$$

with the residues  $R_k := (Cx_k)(y_k^H B) \in \mathbb{C}^{q \times m}$ .

**Dominant Poles** 

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with the residues  $R_k := (Cx_k)(y_k^H B) \in \mathbb{C}^{q \times m}$ .

**Note:** this follows using the spectral decomposition  $A = XDX^{-1}$ , with

 $X = [x_1, \dots, x_n]$  the right and  $X^{-1} =: Y = [y_1, \dots, y_n]^H$  the left eigenvector matrices:

$$G(s) = C(sI - XDX^{-1})^{-1}B = CX(sI - \operatorname{diag}\{\lambda_1, \dots, \lambda_n\})^{-1}YB$$

$$= [Cx_1, \dots, Cx_n] \begin{bmatrix} \frac{1}{s - \lambda_1} & & \\ & \ddots & \\ & & \frac{1}{s - \lambda_n} \end{bmatrix} \begin{bmatrix} y_1^H B \\ \vdots \\ y_n^H B \end{bmatrix}$$

$$= \sum_{k=1}^n \frac{(Cx_k)(y_k^H B)}{s - \lambda_k}.$$

**Dominant Poles** 

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with the residues  $R_k := (Cx_k)(y_k^H B) \in \mathbb{C}^{q \times m}$ .

**Note:**  $R_k = (Cx_k)(y_k^H B)$  are the residues of G in the sense of the residue theorem of complex analysis:

$$\operatorname{res} (G, \lambda_{\ell}) = \lim_{s \to \lambda_{\ell}} (s - \lambda_{\ell}) G(s) = \sum_{k=1}^{n} \underbrace{\lim_{s \to \lambda_{\ell}} \frac{s - \lambda_{\ell}}{s - \lambda_{k}}}_{= \begin{cases} 0 \text{ for } k \neq \ell \\ 1 \text{ for } k = \ell \end{cases}}_{= \begin{cases} R_{\ell} = R_{\ell}.$$

**Dominant Poles** 

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with the residues  $R_k := (Cx_k)(y_k^H B) \in \mathbb{C}^{q \times m}$ .

As projection basis use spaces spanned by right/left eigenvectors corresponding to dominant poles, i.e.,  $(\lambda_i, x_i, y_i)$  with largest

$$||R_k||/|\operatorname{re}(\lambda_k)|$$
.

**Dominant Poles** 

#### Pole-Residue Form of Transfer Function

Consider partial fraction expansion of transfer function with D=0:

$$G(s) = \sum_{k=1}^{n} \frac{R_k}{s - \lambda_k}$$

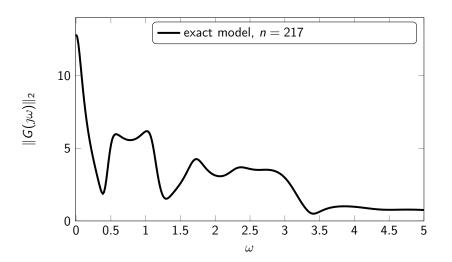
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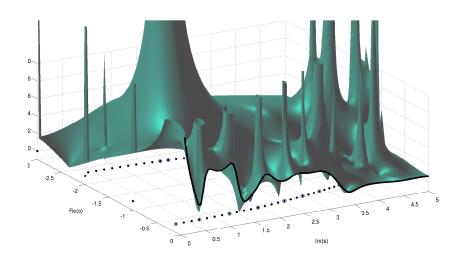
$$||R_k||/|\operatorname{re}(\lambda_k)|$$
.

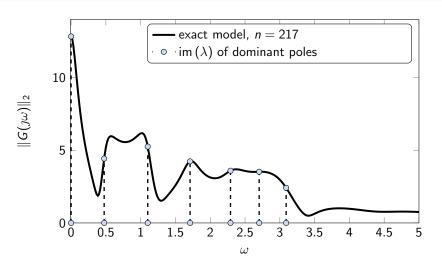
#### Remark

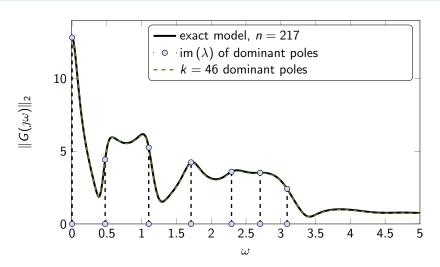
The dominant modes have most important influence on the input-output behavior of the system and are responsible for the "peaks" in the frequency response.

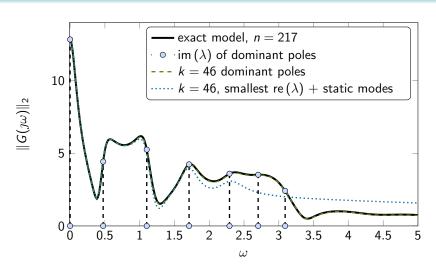


**Dominant Poles**Random SISO Example  $(B, C^T \in \mathbb{R}^n)$ 





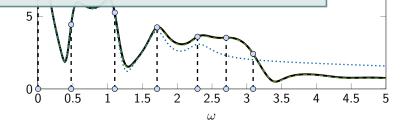




Random SISO Example ( $B, C^T \in \mathbb{R}^n$ )

Algorithms for computing dominant poles and eigenvectors:

- Subspace Accelerated Dominante Pole Algorithm (SADPA),
- Rayleigh-Quotient-Iteration (RQI),
- Jacobi-Davidson-Method.



modes

Basics MOR by Projection RatInt Balanced Truncation Matrix Equations

### **Outline**

- Introduction
- 2 Mathematical Basics
- Model Reduction by Projection
- 4 Interpolatory Model Reduction
  - Padé Approximation
  - A Change of Perspective: Rational Interpolation
  - H2-Optimal Model Reduction
- Balanced Truncation
- 6 Solving Large-Scale Matrix Equations
- Final Remarks

### Padé Approximation

### Idea:

• Consider (even for possibly singular E if  $\lambda E - A$  regular):

$$E\dot{x} = Ax + Bu, \quad y = Cx$$

with transfer function  $G(s) = C(sE - A)^{-1}B$ .

### Padé Approximation

### Idea:

• Consider (even for possibly singular E if  $\lambda E - A$  regular):

$$E\dot{x} = Ax + Bu, \quad y = Cx$$

with transfer function  $G(s) = C(sE - A)^{-1}B$ .

• For  $s_0 \notin \Lambda(A, E)$ :

$$G(s) = C((s_0E - A) + (s - s_0)E)^{-1}B$$

#### Idea:

• Consider (even for possibly singular E if  $\lambda E - A$  regular):

$$E\dot{x} = Ax + Bu, \quad y = Cx$$

with transfer function  $G(s) = C(sE - A)^{-1}B$ .

• For  $s_0 \notin \Lambda(A, E)$ :

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- For  $s_0 = 0$ :  $m_k := -C(A^{-1}E)^k A^{-1}B \implies \text{moments}$ .  $(m_k = -CA^{-(k+1)}B \text{ for } E = I_n)$
- For  $s_0 = \infty$  and  $E = I_n$ :  $m_0 = 0$ ,  $m_k := CA^{k-1}B$  for  $k > 1 \rightsquigarrow$ Markov parameters.

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• As reduced-order model use rth Padé approximant  $\hat{G}$  to G:

$$G(s) = \hat{G}(s) + \mathcal{O}((s-s_0)^{2r}),$$

i.e.,  $m_k = \hat{m}_k$  for k = 0, ..., 2r - 1

 $\rightsquigarrow$  moment matching if  $s_0 < \infty$ ,

 $\rightsquigarrow$  partial realization if  $s_0 = \infty$ .

#### Asymptotic Waveform Evaluation (AWE) [PILLAGE/ROHRER 1990]

Consider SISO case (m = q = 1) and  $s_0 = 0$  for simplicity. With

$$\hat{G}(s) = \frac{\alpha_{r-1}s^{r-1} + \alpha_{r-2}s^{r-2} + \ldots + \alpha_{1}s + \alpha_{0}}{\beta_{r}s^{r} + \beta_{r-1}s^{r-1} + \ldots + \beta_{1}s + 1},$$

the solution of the Padé approximation problem is obtained via solving

$$M\begin{bmatrix} \beta_r \\ \vdots \\ \beta_1 \end{bmatrix} = \begin{bmatrix} -m_r \\ \vdots \\ -m_{2r-1} \end{bmatrix},$$

with the Hankel matrix 
$$M=\left[\begin{array}{ccccc} m_0 & m_1 & m_2 & \dots & m_{r-1} \\ m_1 & m_2 & & \ddots & m_r \\ \\ m_2 & & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ m_{r-1} & m_r & \dots & m_{2r-2} \end{array}\right].$$

Then, with  $\beta_0 := 1$ :  $\alpha_j = \sum_{k=0}^j m_k \beta_{j-k}$ .

The Padé-Lanczos Connection [Gallivan/Grimme/Van Dooren 1994, Freund/Feldmann 1994]

### Theorem [Grimme '97, VILLEMAGNE/SKELTON '87]

Let  $s_* \not\in \Lambda(A, E)$  and

$$\tilde{A} := (s_*E - A)^{-1}E, \qquad \tilde{B} := (s_*E - A)^{-1}B, 
\tilde{A}^* := (s_*E - A)^{-T}E^T, \qquad \tilde{C} := (s_*E - A)^{-T}C^T.$$

If the reduced-order model is obtained by oblique projection onto  $\mathcal{V}\subset\mathbb{R}^n$  along  $\mathcal{W}\subset\mathbb{R}^n$ , and

$$\begin{split} \operatorname{span}\left\{\tilde{B},\tilde{A}\tilde{B},\dots,\tilde{A}^{K-1}\tilde{B}\right\} &\subset & \mathcal{V}, \\ \operatorname{span}\left\{\tilde{C},\tilde{A}^*\tilde{C},\dots,(\tilde{A}^*)^{K-1}\tilde{C}\right\} &\subset & \mathcal{W}, \end{split}$$

then 
$$G(s_*) = \hat{G}(s_*)$$
,  $\frac{d^k}{ds^k}G(s_*) = \frac{d^k}{ds^k}\hat{G}(s_*)$  for  $k = 1, \dots, \ell - 1$ , where 
$$\ell \begin{cases} = 2K & \text{if } m = q = 1; \\ \geq \lfloor \frac{K}{m} \rfloor + \lfloor \frac{K}{d} \rfloor & \text{if } m \neq 1 \text{ or } q \neq 1. \end{cases}$$

The Padé-Lanczos Connection [Gallivan/Grimme/Van Dooren 1994, Freund/Feldmann 1994]

### Padé-via-Lanczos Method (PVL)

• Padé approximation/moment matching yield:

$$m_k = \frac{1}{k!} G^{(k)}(s_0) = \frac{1}{k!} \hat{G}^{(k)}(s_0) = \hat{m}_k, \quad k = 0, \dots, 2K - 1,$$

i.e., Hermite interpolation in  $s_0$ .

 Recall interpolation via projection result ⇒ moments need not be computed explicitly; moment matching is equivalent to projecting state-space onto

$$\mathcal{V} = \operatorname{span}(\tilde{B}, \tilde{A}\tilde{B}, \dots, \tilde{A}^{K-1}\tilde{B}) =: \mathcal{K}_K(\tilde{A}, \tilde{B})$$
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• Computation via unsymmetric Lanczos method.

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**Remark:** Arnoldi (PRIMA) yields only  $G(s) = \hat{G}(s) + \mathcal{O}((s - s_0)^r)$ .

The Padé-Lanczos Connection [Gallivan/Grimme/Van Dooren 1994, Freund/Feldmann 1994]

### Padé-via-Lanczos Method (PVL)

- Computable error estimates/bounds for  $||y \hat{y}||_2$  often very pessimistic or expensive to evaluate.
- Mostly heuristic criteria for choice of expansion points.
   Optimal choice for second-order systems with proportional/Rayleigh damping (Beattie/Gugergin '05).
- Good approximation quality only locally.
- Preservation of physical properties only in special cases (e.g. PRIMA/Arnoldi: V<sup>T</sup>AV is stable if A is negative definite or dissipative exercises); usually requires post processing which (partially) destroys moment matching properties.

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A Change of Perspective: Rational Interpolation

### Computation of reduced-order model by projection

Given an LTI system  $\dot{x} = Ax + Bu$ , y = Cx with transfer function  $G(s) = C(sI_n - A)^{-1}B$ , a reduced-order model is obtained using projection approach with  $V, W \in \mathbb{R}^{n \times r}$  and  $W^T V = I_r$  by computing

$$\hat{A} = W^T A V, \quad \hat{B} = W^T B, \quad \hat{C} = C V.$$

Petrov-Galerkin-type (two-sided) projection:  $W \neq V$ ,

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### Rational Interpolation/Moment-Matching

Choose V, W such that

$$G(s_i) = \hat{G}(s_i), \quad j = 1, \ldots, k,$$

and

$$\frac{d^i}{ds^i}G(s_j) = \frac{d^i}{ds^i}\hat{G}(s_j), \quad i = 1, \dots, K_j, \quad j = 1, \dots, k.$$

A Change of Perspective: Rational Interpolation

### Theorem (simplified) [GRIMME '97, VILLEMAGNE/SKELTON '87]

lf

$$\operatorname{span} \left\{ (s_1 I_n - A)^{-1} B, \dots, (s_k I_n - A)^{-1} B \right\} \subset \operatorname{Ran}(V),$$
 
$$\operatorname{span} \left\{ (s_1 I_n - A)^{-T} C^T, \dots, (s_k I_n - A)^{-T} C^T \right\} \subset \operatorname{Ran}(W),$$

then

$$G(s_j) = \hat{G}(s_j), \quad \frac{d}{ds}G(s_j) = \frac{d}{ds}\hat{G}(s_j), \quad \text{for } j = 1, \dots, k.$$

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#### Remarks:

using Galerkin/one-sided projection yields  $G(s_i) = \hat{G}(s_i)$ , but in general

$$\frac{d}{ds}G(s_j)\neq \frac{d}{ds}\hat{G}(s_j).$$

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#### Remarks:

k = 1, standard Krylov subspace(s) of dimension  $K \leadsto$  moment-matching methods/Padé approximation,

$$\frac{d^i}{ds^i}G(s_1)=\frac{d^i}{ds^i}\hat{G}(s_1), \quad i=0,\ldots,K-1(+K).$$

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#### Remarks:

computation of V, W from rational Krylov subspaces, e.g.,

- dual rational Arnoldi/Lanczos [GRIMME '97],
- Iterative Rational Krylov-Algo. [Antoulas/Beattle/Gugercin '07].

### Best $\mathcal{H}_2$ -norm approximation problem

 $\operatorname{\mathsf{arg\,min}}_{\hat{G} \in \mathcal{H}_2 \ \operatorname{\mathsf{of\,order}} \ \leq r} \| \mathcal{G} - \hat{\mathcal{G}} \|_2.$ Find

## $\mathcal{H}_2$ -Optimal Model Reduction

### Best $\mathcal{H}_2$ -norm approximation problem

Find 
$$\arg \min_{\hat{G} \in \mathcal{H}_2 \text{ of order } \leq r} ||G - \hat{G}||_2$$
.

 $\rightsquigarrow$  First-order necessary  $\mathcal{H}_2$ -optimality conditions:

For SISO systems

$$G(-\mu_i) = \hat{G}(-\mu_i),$$
  

$$G'(-\mu_i) = \hat{G}'(-\mu_i),$$

where  $\mu_i$  are the poles of the reduced transfer function  $\hat{G}$ .

### $\mathcal{H}_2$ -Optimal Model Reduction

### Best $\mathcal{H}_2$ -norm approximation problem

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For MIMO systems

$$G(-\mu_i)\tilde{B}_i = \hat{G}(-\mu_i)\tilde{B}_i, \qquad \text{for } i = 1, \dots, r,$$

$$\tilde{C}_i^T G(-\mu_i) = \tilde{C}_i^T \hat{G}(-\mu_i), \qquad \text{for } i = 1, \dots, r,$$

$$\tilde{C}_i^T G'(-\mu_i)\tilde{B}_i = \tilde{C}_i^T \hat{G}'(-\mu_i)\tilde{B}_i, \qquad \text{for } i = 1, \dots, r,$$

where  $T^{-1}\hat{A}T = \mathrm{diag}\left\{\mu_1,\ldots,\mu_r\right\} = \mathrm{spectral}$  decomposition and

$$\tilde{B} = \hat{B}^T T^{-T}, \quad \tilde{C} = \hat{C}T.$$

→ tangential interpolation conditions.

Interpolation of the Transfer Function by Projection

Construct reduced transfer function by Petrov-Galerkin projection  $\mathcal{P} = VW^T$ , i.e.

$$\hat{G}(s) = CV (sI - W^T AV)^{-1} W^T B,$$

where V and W are given as the rational Krylov subspaces

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for  $i = 1, \ldots, r$  as desired.

 $\rightsquigarrow$  iterative algorithms (IRKA/MIRIAm) that yield  $\mathcal{H}_2$ -optimal models.

[Gugercin et al. '06], [Bunse-Gerstner et al. '07], [Van Dooren et al. '08]

Interpolation of the Transfer Function by Projection

Construct reduced transfer function by Petrov-Galerkin projection  $\mathcal{P} = VW^T$ , i.e.

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# **H<sub>2</sub>-Optimal Model Reduction**

The Basic IRKA Algorithm

#### **Algorithm 1** IRKA (MIMO version/MIRIAm)

**Input:** A stable, B, C,  $\hat{A}$  stable,  $\hat{B}$ ,  $\hat{C}$ ,  $\delta > 0$ .

Output: Aopt, Bopt, Copt

1: while 
$$(\max_{j=1,...,r}\left\{rac{|\mu_j-\mu_j^{
m old}|}{|\mu_j|}
ight\}>\delta)$$
 do

2: diag  $\{\mu_1, \dots, \mu_r\} := T^{-1}\hat{A}T$  = spectral decomposition,  $\tilde{B} = \hat{B}^H T^{-T}$ ,  $\tilde{C} = \hat{C}T$ .

3: 
$$V = \left[ (-\mu_1 I - A)^{-1} B \tilde{B}_1, \dots, (-\mu_r I - A)^{-1} B \tilde{B}_r \right]$$

4: 
$$W = \left[ (-\mu_1 I - A^T)^{-1} C^T \tilde{C}_1, \dots, (-\mu_r I - A^T)^{-1} C^T \tilde{C}_r \right]$$

5: 
$$V = \operatorname{orth}(V), W = \operatorname{orth}(W), W = W(V^H W)^{-1}$$

6: 
$$\hat{A} = W^H A V$$
,  $\hat{B} = W^H B$ ,  $\hat{C} = C V$ 

7: end while

8: 
$$A^{opt} = \hat{A}$$
,  $B^{opt} = \hat{B}$ ,  $C^{opt} = \hat{C}$ 

### **Outline**

- Introduction
- 2 Mathematical Basics
- Model Reduction by Projection
- 4 Interpolatory Model Reduction
- 5 Balanced Truncation
  - The basic method
  - Theoretical Background
  - Singular Perturbation Approximation
  - Balancing-Related Methods
- 6 Solving Large-Scale Matrix Equations
- Final Remarks

### Basic principle:

• Recall: a stable system  $\Sigma$ , realized by (A, B, C, D), is called balanced, if the Gramians, i.e., solutions P, Q of the Lyapunov equations

$$AP + PA^{T} + BB^{T} = 0, \qquad A^{T}Q + QA + C^{T}C = 0,$$

satisfy: 
$$P = Q = \operatorname{diag}(\sigma_1, \dots, \sigma_n)$$
 with  $\sigma_1 \ge \sigma_2 \ge \dots \ge \sigma_n > 0$ .

•  $\Lambda(PQ)^{\frac{1}{2}} = \{\sigma_1, \dots, \sigma_n\}$  are the Hankel singular values (HSVs) of  $\Sigma$ .

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$$\mathcal{T}: (A, B, C, D) \mapsto (TAT^{-1}, TB, CT^{-1}, D)$$

$$= \left( \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}, \begin{bmatrix} B_1 \\ B_2 \end{bmatrix}, \begin{bmatrix} C_1 & C_2 \end{bmatrix}, D \right)$$

• Truncation  $\rightsquigarrow$   $(\hat{A}, \hat{B}, \hat{C}, \hat{D}) := (A_{11}, B_1, C_1, D).$ 

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The HSVs  $\Lambda(PQ)^{\frac{1}{2}}=\{\sigma_1,\ldots,\sigma_n\}$  are system invariants: they are preserved under

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in transformed coordinates, the Gramians satisfy

$$(TAT^{-1})(TPT^{T}) + (TPT^{T})(TAT^{-1})^{T} + (TB)(TB)^{T} = 0,$$

$$(TAT^{-1})^{T}(T^{-T}QT^{-1}) + (T^{-T}QT^{-1})(TAT^{-1}) + (CT^{-1})^{T}(CT^{-1}) = 0$$

$$\Rightarrow (TPT^{T})(T^{-T}QT^{-1}) = TPQT^{-1},$$

hence  $\Lambda(PQ) = \Lambda((TPT^T)(T^{-T}QT^{-1})).$ 

#### Implementation: SR Method

- Compute (Cholesky) factors of the Gramians,  $P = S^T S$ ,  $Q = R^T R$ .
- $\odot$  ROM is  $(W^TAV, W^TB, CV, D)$ , where

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$$= \Sigma_{1}^{-\frac{1}{2}}[I_{r}, 0] \begin{bmatrix} \Sigma_{1} \\ \Sigma_{2} \end{bmatrix} \begin{bmatrix} I_{r} \\ 0 \end{bmatrix} \Sigma_{1}^{-\frac{1}{2}}$$

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 $\implies VW^T$  is an oblique projector, hence balanced truncation is a Petrov-Galerkin projection method.

## Properties:

- Reduced-order model is stable with HSVs  $\sigma_1, \ldots, \sigma_r$ .
- Adaptive choice of *r* via computable error bound:

$$||y - \hat{y}||_2 \le \left(2\sum_{k=r+1}^n \sigma_k\right) ||u||_2.$$

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## Linear, Time-Invariant (LTI) Systems

$$\begin{array}{lll} \dot{x} & = & Ax + Bu, & A \in \mathbb{R}^{n \times n}, & B \in \mathbb{R}^{n \times m}, \\ y & = & Cx + Du, & C \in \mathbb{R}^{q \times n}, & D \in \mathbb{R}^{q \times m}. \end{array}$$

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Assumptions (for now):  $t_0 = 0$ ,  $x_0 = x(0) = 0$ , D = 0.

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#### State-Space Description for I/O-Relation

Variation-of-constants ⇒

$$\mathcal{S}: u\mapsto y, \quad y(t)=\int_{-\infty}^t C\mathrm{e}^{A(t- au)}Bu( au)\,d au \quad ext{for all } t\in\mathbb{R}.$$

Theoretical Background

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- Recall:  $A \in \mathbb{R}^{n \times m}$  is a linear operator,  $A : \mathbb{R}^m \to \mathbb{R}^n$ !
- Basic Idea: use SVD approximation as for matrix A!
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Instead of

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Theoretical Background

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- $\rightsquigarrow$  Hankel singular values  $\{\sigma_j\}_{j=1}^{\infty}: \sigma_1 \geq \sigma_2 \geq \ldots \geq 0.$
- $\implies$  SVD-type approximation of  $\mathcal{H}$  possible!

Theoretical Background

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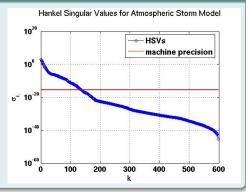
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#### Theoretical Background

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Theoretical Background

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#### Theorem

Let P, Q be the controllability and observability Gramians of an LTI system  $\Sigma$ . Then the Hankel singular values  $\Lambda(PQ)^{\frac{1}{2}} = \{\sigma_1, \dots, \sigma_n\}$  are the singular values of the Hankel operator associated to  $\Sigma$ .

The Hankel Singular Values are Singular Values!

#### Theorem

Let P, Q be the controllability and observability Gramians of an LTI system  $\Sigma$ . Then the Hankel singular values  $\Lambda(PQ)^{\frac{1}{2}} = \{\sigma_1, \dots, \sigma_n\}$  are the singular values of the Hankel operator associated to  $\Sigma$ .

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# Balanced Truncation

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## $\mathsf{Theorem}$

Let the reduced-order system  $\hat{\Sigma}: (\hat{A}, \hat{B}, \hat{C}, \hat{D})$  with  $r < \hat{n}$  be computed by balanced truncation. Then the reduced-order model  $\hat{\Sigma}$  is balanced, stable, minimal, and its HSVs are  $\sigma_1, \ldots, \sigma_r$ .

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Let the reduced-order system  $\hat{\Sigma}: (\hat{A}, \hat{B}, \hat{C}, \hat{D})$  with  $r < \hat{n}$  be computed by balanced truncation. Then the reduced-order model  $\hat{\Sigma}$  is balanced, stable, minimal, and its HSVs are  $\sigma_1, \ldots, \sigma_r$ .

Proof: Note that in balanced coordinates, the Gramians are diagonal and equal to

$$\operatorname{diag}(\Sigma_1, \Sigma_2) = \operatorname{diag}(\sigma_1, \dots, \sigma_r, \sigma_{r+1}, \dots, \sigma_n).$$

Hence, the Gramian satisfies

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} \Sigma_1 \\ \Sigma_2 \end{bmatrix} + \begin{bmatrix} \Sigma_1 \\ \Sigma_2 \end{bmatrix} \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}^T + \begin{bmatrix} B_1 \\ B_2 \end{bmatrix} \begin{bmatrix} B_1 \\ B_2 \end{bmatrix}^T = 0,$$

whence we obtain the "controllability Lyapunov equation" of the reduced-order system,

$$A_{11}\Sigma_1 + \Sigma_1 A_{11}^T + B_1 B_1^T = 0.$$

The result follows from  $\hat{A} = A_{11}, \hat{B} = B_1, \Sigma_1 > 0$ , the solution theory of Lyapunov equations and the analogous considerations for the observability Gramian. (Minimality is a simple consequence of  $\hat{P} = \Sigma_1 = \hat{Q} > 0$ .)

#### Assume the system

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} B_1 \\ B_2 \end{bmatrix} u, \quad y = \begin{bmatrix} C_1, C_2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + Du$$

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Balanced truncation would set  $x_2 = 0$  and use  $(A_{11}, B_1, C_1, D)$  as reduced-order model, thereby the information present in the remaining model is ignored!

# Singular Perturbation Approximation (aka Balanced Residualization)

Assume the system

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Balanced truncation would set  $x_2 = 0$  and use  $(A_{11}, B_1, C_1, D)$  as reduced-order model, thereby the information present in the remaining model is ignored!

Particularly, if  $G(0) = \hat{G}(0)$  ("zero steady-state error") is required, one can apply the same condensation technique as in Guyan reduction: instead of  $x_2 = 0$ , set  $\dot{x}_2 = 0$ . This yields the reduced-order model

$$\dot{x}_1 = (A_{11} - A_{12}A_{22}^{-1}A_{21})x_1 + (B_1 - A_{12}A_{22}^{-1}B_2)u, 
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with

- the same properties as the reduced-order model w.r.t. stability, minimality, error bound, but  $\hat{D} \neq D$ ;
- zero steady-state error,  $G(0) = \hat{G}(0)$  as desired.

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- the same properties as the reduced-order model w.r.t. stability, minimality, error bound, but  $\hat{D} \neq D$ ;
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#### Note:

- $A_{22}$  invertible as in balanced coordinates,  $A_{22}\Sigma_2 + \Sigma_2 A_{22}^T + B_2 B_2^T = 0$  and  $(A_{22}, B_2)$  controllable,  $\Sigma_2 > 0 \Rightarrow A_{22}$  stable.
- If the original system is not balanced, first compute a minimal realization by applying balanced truncation with  $r = \hat{n}$ .

Given positive semidefinite matrices  $P = S^T S$ ,  $Q = R^T R$ , compute balancing state-space transformation so that

$$P = Q = \operatorname{diag}(\sigma_1, \ldots, \sigma_n) = \Sigma, \quad \sigma_1 \ge \ldots \ge \sigma_n > 0,$$

and truncate corresponding realization at size r with  $\sigma_r > \sigma_{r+1}$ .

# **Balancing-Related Methods**

# Basic Principle

Given positive semidefinite matrices  $P = S^T S$ ,  $Q = R^T R$ , compute balancing state-space transformation so that

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## Classical Balanced Truncation (BT)

[Mullis/Roberts '76, Moore '81]

- P = controllability Gramian of system given by (A, B, C, D).
- Q = observability Gramian of system given by (A, B, C, D).
- P, Q solve dual Lyapunov equations

$$AP + PA^{T} + BB^{T} = 0, \qquad A^{T}Q + QA + C^{T}C = 0.$$

# **Balancing-Related Methods**

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# LQG Balanced Truncation (LQGBT) [Jonckheere/Silverman '83]

- P/Q = controllability/observability Gramian of closed-loop system based on LQG compensator.
- P, Q solve dual algebraic Riccati equations (AREs)

$$0 = AP + PA^{T} - PC^{T}CP + B^{T}B,$$

$$0 = A^T Q + QA - QBB^T Q + C^T C.$$

Given positive semidefinite matrices  $P = S^T S$ ,  $Q = R^T R$ , compute balancing state-space transformation so that

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# Balanced Stochastic Truncation (BST) [Desai/Pal '84, Green '88]

- P = controllability Gramian of system given by (A, B, C, D), i.e., solution of Lyapunov equation  $AP + PA^T + BB^T = 0$ .
- Q = observability Gramian of right spectral factor of power spectrum of system given by (A, B, C, D), i.e., solution of ARE

$$\hat{A}^T Q + Q \hat{A} + Q B_W (DD^T)^{-1} B_W^T Q + C^T (DD^T)^{-1} C = 0,$$

where  $\hat{A} := A - B_W(DD^T)^{-1}C$ ,  $B_W := BD^T + PC^T$ .

Given positive semidefinite matrices  $P = S^T S$ ,  $Q = R^T R$ , compute balancing state-space transformation so that

$$P = Q = \operatorname{diag}(\sigma_1, \ldots, \sigma_n) = \Sigma, \quad \sigma_1 \ge \ldots \ge \sigma_n > 0,$$

and truncate corresponding realization at size r with  $\sigma_r > \sigma_{r+1}$ .

# Positive-Real Balanced Truncation (PRBT)

- Based on positive-real equations, related to positive real (Kalman-Yakubovich-Popov-Anderson) lemma.
- P, Q solve dual AREs

$$0 = \bar{A}P + P\bar{A}^{T} + PC^{T}\bar{R}^{-1}CP + B\bar{R}^{-1}B^{T},$$

$$0 = \bar{A}^{T}Q + Q\bar{A} + QB\bar{R}^{-1}B^{T}Q + C^{T}\bar{R}^{-1}C,$$

where  $\bar{R} = D + D^T$ .  $\bar{A} = A - B\bar{R}^{-1}C$ .

Given positive semidefinite matrices  $P = S^T S$ ,  $Q = R^T R$ , compute balancing state-space transformation so that

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and truncate corresponding realization at size r with  $\sigma_r > \sigma_{r+1}$ .

# Other Balancing-Based Methods

- Bounded-real balanced truncation (BRBT) based on bounded real lemma [Opdenacker/Jonckheere '88];
- $H_{\infty}$  balanced truncation (HinfBT) closed-loop balancing based on  $H_{\infty}$  compensator [Mustafa/Glover '91].

Both approaches require solution of dual AREs.

Frequency-weighted versions of the above approaches.

# Balancing-Related Methods

#### **Properties**

- Guaranteed preservation of physical properties like
  - stability (all).
  - passivity (PRBT),
  - minimum phase (BST).
- Computable error bounds, e.g.,

$$\begin{split} \text{BT:} \quad & \|G - G_r\|_{\infty} \quad \leq 2 \; \sum_{j=r+1}^n \sigma_j^{BT}, \\ \text{LQGBT:} \quad & \|G - G_r\|_{\infty} \quad \leq \; 2 \sum_{j=r+1}^n \frac{\sigma_j^{LQG}}{\sqrt{1 + (\sigma_j^{LQG})^2}} \\ \text{BST:} \quad & \|G - G_r\|_{\infty} \quad \leq \left( \prod_{j=r+1}^n \frac{1 + \sigma_j^{BST}}{1 - \sigma_j^{BST}} - 1 \right) \|G\|_{\infty}, \end{split}$$

- Can be combined with singular perturbation approximation for steady-state performance.
- Computations can be modularized.

- Solving Large-Scale Matrix Equations
  - Linear Matrix Equations
  - Numerical Methods for Solving Lyapunov Equations
  - Solving Large-Scale Algebraic Riccati Equations
  - Software

Large-Scale Algebraic Lyapunov and Riccati Equations

Algebraic Riccati equation (ARE) for  $A, G = G^T, W = W^T \in \mathbb{R}^{n \times n}$  given and  $X \in \mathbb{R}^{n \times n}$  unknown:

$$0 = \mathcal{R}(X) := A^T X + XA - XGX + W.$$

 $G = 0 \Longrightarrow$ Lyapunov equation

$$0 = \mathcal{L}(X) := A^T X + XA + W.$$

- $n = 10^3 10^6 \ (\Longrightarrow 10^6 10^{12} \ \text{unknowns!})$
- A has sparse representation  $(A = -M^{-1}S)$  for FEM,
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Large-Scale Algebraic Lyapunov and Riccati Equations

Algebraic Riccati equation (ARE) for  $A, G = G^T, W = W^T \in \mathbb{R}^{n \times n}$  given and  $X \in \mathbb{R}^{n \times n}$  unknown:

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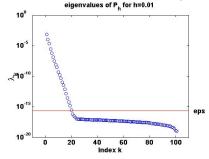
Low-Rank Approximation

Consider spectrum of ARE solution (analogous for Lyapunov equations).

# Example:

- Linear 1D heat equation with point control,
- $\Omega = [0, 1],$
- FEM discretization using linear B-splines,
- $h = 1/100 \implies n = 101$ .

Idea: 
$$X = X^T > 0 \implies$$



$$X = ZZ^{T} = \sum_{k=1}^{n} \lambda_{k} z_{k} z_{k}^{T} \approx Z^{(r)} (Z^{(r)})^{T} = \sum_{k=1}^{r} \lambda_{k} z_{k} z_{k}^{T}.$$

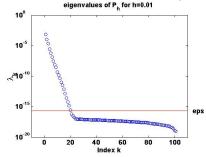
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Linear Matrix Equations

## **Equations without symmetry**

Sylvester equation

discrete Sylvester equation

$$AX + XB = W$$

$$AXB - X = W$$

with data  $A \in \mathbb{R}^{n \times n}$ ,  $B \in \mathbb{R}^{m \times m}$ ,  $W \in \mathbb{R}^{n \times m}$  and unknown  $X \in \mathbb{R}^{n \times m}$ .

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Solvability

Using the Kronecker (tensor) product, AX + XB = W is equivalent to

$$((I_m \otimes A) + (B^T \otimes I_n)) \operatorname{vec}(X) = \operatorname{vec}(W).$$

Hence,

Sylvester equation has a unique solution

$$M := (I_m \otimes A) + (B^T \otimes I_n)$$
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$$0 \not\in \Lambda(M) = \Lambda((I_m \otimes A) + (B^T \otimes I_n)) = \{\lambda_j + \mu_k, \mid \lambda_j \in \Lambda(A), \mu_k \in \Lambda(B)\}.$$

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# Corollary

A, B Hurwitz  $\Longrightarrow$  Sylvester equation has unique solution.

# Linear Matrix Equations

#### Complexity Issues

Solving the Sylvester equation

$$AX + XB = W$$

via the equivalent linear system of equations

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### requires

- LU factorization of  $nm \times nm$  matrix; for  $n \approx m$ , complexity is  $\frac{2}{3}n^6$ ;
- storing  $n \cdot m$  unknowns: for  $n \approx m$  we have  $n^4$  data for X.

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 $n = m = 1,000 \Rightarrow$  Gaussian elimination on an Intel core i7 (Westmere, 6) cores, 3.46 GHz  $\leftrightarrow$  83.2 GFLOP peak) would take > 94 DAYS and 7.3 TB of memory!

# Numerical Methods for Solving Lyapunov Equations Traditional Methods

Bartels-Stewart method for Sylvester and Lyapunov equation (1yap); Hessenberg-Schur method for Sylvester equations (1yap); Hammarling's method for Lyapunov equations  $AX + XA^T + GG^T = 0$  with A Hurwitz (1yapchol).

All based on the fact that if  $A, B^T$  are in Schur form, then

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is block-upper triangular. Hence, solve Mx = b by back-substitution.

- Clever implementation of back-substitution process requires nm(n+m) flops.
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  Hessenberg-Schur method).
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### The Sign Function Method

### Definition

For  $Z \in \mathbb{R}^{n \times n}$  with  $\Lambda(Z) \cap i\mathbb{R} = \emptyset$  and Jordan canonical form

$$Z = S \begin{bmatrix} J^+ & 0 \\ 0 & J^- \end{bmatrix} S^{-1}$$

the matrix sign function is

$$\operatorname{sign}(Z) := S \left[ \begin{array}{cc} I_k & 0 \\ 0 & -I_{n-k} \end{array} \right] S^{-1}.$$

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### Lemma

Let  $T \in \mathbb{R}^{n \times n}$  be nonsingular and Z as before, then

$$\operatorname{sign}(TZT^{-1}) = T\operatorname{sign}(Z)T^{-1}$$

# The Sign Function Method

# Computation of sign (Z)

 $\operatorname{sign}(Z)$  is root of  $I_n \Longrightarrow$  use Newton's method to compute it:

$$Z_0 \leftarrow Z, \qquad Z_{j+1} \leftarrow \frac{1}{2} \left( c_j Z_j + \frac{1}{c_j} Z_j^{-1} \right), \qquad j = 1, 2, \dots$$

$$\implies \operatorname{sign}(Z) = \lim_{j \to \infty} Z_j.$$

 $c_i > 0$  is scaling parameter for convergence acceleration and rounding error minimization, e.g.

$$c_j = \sqrt{\frac{\|Z_j^{-1}\|_F}{\|Z_j\|_F}},$$

based on "equilibrating" the norms of the two summands [Higham '86].

# Solving Lyapunov Equations with the Matrix Sign Function Method

#### Key observation:

If  $X \in \mathbb{R}^{n \times n}$  is a solution of  $AX + XA^T + W = 0$ , then

$$\underbrace{\begin{bmatrix} I_n & -X \\ 0 & I_n \end{bmatrix}}_{=T^{-1}} \underbrace{\begin{bmatrix} A & W \\ 0 & -A^T \end{bmatrix}}_{=:H} \underbrace{\begin{bmatrix} I_n & X \\ 0 & I_n \end{bmatrix}}_{=:T} = \begin{bmatrix} A & 0 \\ 0 & -A^T \end{bmatrix}.$$

Hence, if A is Hurwitz (i.e., asymptotically stable), then

$$\begin{aligned} \operatorname{sign}\left(H\right) &= & \operatorname{sign}\left(T\begin{bmatrix}A & 0 \\ 0 & -A^T\end{bmatrix}T^{-1}\right) = T\operatorname{sign}\left(\begin{bmatrix}A & 0 \\ 0 & -A^T\end{bmatrix}\right)T^{-1} \\ &= \begin{bmatrix}-I_n & 2X \\ 0 & I_n\end{bmatrix}. \end{aligned}$$

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$$= \begin{bmatrix} -I_n & 2X \\ 0 & I_n \end{bmatrix}.$$

Apply sign function iteration 
$$Z \leftarrow \frac{1}{2}(Z + Z^{-1})$$
 to  $H = \begin{bmatrix} A & W \\ 0 & -A^T \end{bmatrix}$ :

$$H + H^{-1} = \begin{bmatrix} A & W \\ 0 & -A^T \end{bmatrix} + \begin{bmatrix} A^{-1} & A^{-1}WA^{-T} \\ 0 & -A^{-T} \end{bmatrix}$$

⇒ Sign function iteration for Lyapunov equation:

$$A_0 \leftarrow A, \quad A_{j+1} \leftarrow \frac{1}{2} \left( A_j + A_j^{-1} \right),$$
  
 $W_0 \leftarrow G, \quad W_{j+1} \leftarrow \frac{1}{2} \left( W_j + A_j^{-1} W_j A_j^{-T} \right),$   $j = 0, 1, 2, \dots$ 

Define 
$$A_{\infty} := \lim_{j \to \infty} A_j$$
,  $W_{\infty} := \lim_{j \to \infty} W_j$ .

#### **Theorem**

If A is Hurwitz, then

$$A_{\infty} = -I_n$$
 and  $X = \frac{1}{2}W_{\infty}$ .

Recall sign function iteration for  $AX + XA^T + W = 0$ :

$$A_0 \leftarrow A, \quad A_{j+1} \leftarrow \frac{1}{2} (A_j + A_j^{-1}), W_0 \leftarrow G, \quad W_{j+1} \leftarrow \frac{1}{2} (W_j + A_i^{-1} W_j A_i^{-T}),$$
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Now consider the second iteration for  $W = BB^T$ , starting with  $W_0 = BB^T =: B_0B_0^T$ :

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Hence, obtain factored iteration

$$B_{j+1} \leftarrow \frac{1}{\sqrt{2}} \begin{bmatrix} B_j & A_j^{-1} B_j \end{bmatrix}$$

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# Factored sign function iteration for $A(SS^T) + (SS^T)A^T + BB^T = 0$

$$A_0 \leftarrow A, \quad A_{j+1} \leftarrow \frac{1}{2} \left( A_j + A_j^{-1} \right), B_0 \leftarrow B, \quad B_{j+1} \leftarrow \frac{1}{\sqrt{2}} \left[ B_j \quad A_j^{-1} B_j \right], \qquad j = 0, 1, 2, \dots.$$

#### Remarks:

• To get both Gramians, run in parallel

$$C_{j+1} \leftarrow \frac{1}{\sqrt{2}} \begin{bmatrix} C_j \\ C_j A_i^{-1} \end{bmatrix}.$$

- To avoid growth in numbers of columns of  $B_j$  (or rows of  $C_j$ ): column compression by RRLQ or truncated SVD.
- Several options to incorporate scaling, e.g., scale "A"-iteration only.
- Simple stopping cirterion:  $||A_j + I_n||_F \le tol$ .

# Numerical Methods for Solving Lyapunov Equations The ADI Method

Recall Peaceman Rachford ADI:

Consider Au = s where  $A \in \mathbb{R}^{n \times n}$  spd,  $s \in \mathbb{R}^n$ . ADI Iteration Idea:

Decompose A = H + V with  $H, V \in \mathbb{R}^{n \times n}$  such that

$$(H+pI)v = r$$
$$(V+pI)w = t$$

can be solved easily/efficiently.

### ADI Iteration

If  $H, V \text{ spd} \Rightarrow \exists p_k, k = 1, 2, \dots \text{ such that}$ 

$$u_{0} = 0$$

$$(H + p_{k}I)u_{k-\frac{1}{2}} = (p_{k}I - V)u_{k-1} + s$$

$$(V + p_{k}I)u_{k} = (p_{k}I - H)u_{k-\frac{1}{2}} + s$$

converges to  $u \in \mathbb{R}^n$  solving Au = s.

# Numerical Methods for Solving Lyapunov Equations The ADI Method

Recall Peaceman Rachford ADI:

Consider Au = s where  $A \in \mathbb{R}^{n \times n}$  spd,  $s \in \mathbb{R}^n$ . ADI Iteration Idea:

Decompose A = H + V with  $H, V \in \mathbb{R}^{n \times n}$  such that

$$(H+pI)v = r$$
$$(V+pI)w = t$$

can be solved easily/efficiently.

## **ADI** Iteration

If  $H, V \text{ spd} \Rightarrow \exists p_k, k = 1, 2, \dots \text{ such that}$ 

$$\begin{array}{rcl} u_0 & = & 0 \\ (H+p_k I)u_{k-\frac{1}{2}} & = & (p_k I-V)u_{k-1}+s \\ (V+p_k I)u_k & = & (p_k I-H)u_{k-\frac{1}{2}}+s \end{array}$$

converges to  $u \in \mathbb{R}^n$  solving Au = s.

The Lyapunov operator

$$\mathcal{L}: P \mapsto AX + XA^T$$

can be decomposed into the linear operators

$$\mathcal{L}_H: X \mapsto AX, \qquad \mathcal{L}_V: X \mapsto XA^T.$$

In analogy to the standard ADI method we find the

### ADI iteration for the Lyapunov equation

$$\begin{array}{rcl} X_0 & = & 0 \\ (A+p_kI)X_{k-\frac{1}{2}} & = & -W-X_{k-1}(A^T-p_kI) \\ (A+p_kI)X_k^T & = & -W-X_{k-\frac{1}{2}}^T(A^T-p_kI). \end{array}$$

#### Numerical Methods for Solving Lyapunov Equations Low-Rank ADI

Consider  $AX + XA^T = -BB^T$  for stable A;  $B \in \mathbb{R}^{n \times m}$  with  $m \ll n$ .

## ADI iteration for the Lyapunov equation

[Wachspress '95]

For  $k = 1, \ldots, k_{\text{max}}$ 

$$\begin{array}{rcl} X_0 & = & 0 \\ (A+p_kI)X_{k-\frac{1}{2}} & = & -BB^T - X_{k-1}(A^T - p_kI) \\ (A+p_kI)X_k^T & = & -BB^T - X_{k-\frac{1}{2}}^T(A^T - p_kI) \end{array}$$

Rewrite as one step iteration and factorize  $X_k = Z_k Z_k^T$ ,  $k = 0, \dots, k_{\text{max}}$ 

$$Z_{0}Z_{0}^{T} = 0$$

$$Z_{k}Z_{k}^{T} = -2p_{k}(A + p_{k}I)^{-1}BB^{T}(A + p_{k}I)^{-T} + (A + p_{k}I)^{-1}(A - p_{k}I)Z_{k-1}^{T}Z_{k-1}^{T}(A - p_{k}I)^{T}(A + p_{k}I)^{-T}$$

 $\ldots \leadsto$  low-rank Cholesky factor ADI

[PENZL '97/'00, LI/WHITE '99/'02, B./LI/PENZL '99/'08, GUGERCIN/SORENSEN/ANTOULAS '03

#### Numerical Methods for Solving Lyapunov Equations Low-Rank ADI

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... ~> low-rank Cholesky factor ADI

ze 50, 60, Geodelich, Southwell, Hillochie 60]

#### Numerical Methods for Solving Lyapunov Equations Low-Rank ADI

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[Wachspress '95]

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... --> low-rank Cholesky factor ADI

[Penzl '97/'00, Li/White '99/'02, B./Li/Penzl '99/'08, Gugercin/Sorensen/Antoulas '03]

# Solving Large-Scale Matrix Equations

**Numerical Methods for Solving Lyapunov Equations** 

$$Z_k = [\sqrt{-2p_k}(A + p_k I)^{-1}B, (A + p_k I)^{-1}(A - p_k I)Z_{k-1}]$$

[Penzl '00]

Observing that  $(A - p_i I)$ ,  $(A + p_k I)^{-1}$  commute, we rewrite  $Z_{k_{\text{max}}}$  as

$$Z_{k_{\max}} = [z_{k_{\max}}, P_{k_{\max}-1}z_{k_{\max}}, P_{k_{\max}-2}(P_{k_{\max}-1}z_{k_{\max}}), \dots, P_1(P_2 \cdots P_{k_{\max}-1}z_{k_{\max}})],$$

[Li/White '02]

where

$$z_{k_{\text{max}}} = \sqrt{-2p_{k_{\text{max}}}}(A + p_{k_{\text{max}}}I)^{-1}B$$

and

$$P_i := \frac{\sqrt{-2p_i}}{\sqrt{-2p_{i+1}}} \left[ I - (p_i + p_{i+1})(A + p_i I)^{-1} \right].$$

# Solving Large-Scale Matrix Equations

**Numerical Methods for Solving Lyapunov Equations** 

$$Z_k = [\sqrt{-2p_k}(A + p_k I)^{-1}B, (A + p_k I)^{-1}(A - p_k I)Z_{k-1}]$$

[Penzl '00]

Observing that  $(A - p_i I)$ ,  $(A + p_k I)^{-1}$  commute, we rewrite  $Z_{k_{\text{max}}}$  as

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[Li/White '02]

where

$$z_{k_{\text{max}}} = \sqrt{-2p_{k_{\text{max}}}} (A + p_{k_{\text{max}}} I)^{-1} B$$

and

$$P_i := \frac{\sqrt{-2p_i}}{\sqrt{-2p_{i+1}}} \left[ I - (p_i + p_{i+1})(A + p_i I)^{-1} \right].$$

# Algorithm [Penzl '97/'00, Li/White '99/'02, B. 04, B./Li/Penzl '99/'08]

$$V_1 \leftarrow \sqrt{-2\operatorname{re} p_1}(A+p_1I)^{-1}B, \quad Z_1 \leftarrow V_1$$

FOR  $k=2,3,\ldots$ 

$$V_k \leftarrow \sqrt{\frac{\operatorname{re} p_k}{\operatorname{re} p_{k-1}}} \left(V_{k-1} - (p_k + \overline{p_{k-1}})(A+p_kI)^{-1}V_{k-1}\right)$$

$$Z_k \leftarrow \left[\begin{array}{cc} Z_{k-1} & V_k \end{array}\right]$$

$$Z_k \leftarrow \operatorname{rrlq}(Z_k,\tau) \quad \text{column compression}$$

At convergence,  $Z_{k_{\text{max}}}Z_{k_{\text{max}}}^T \approx X$ , where (without column compression)

$$Z_{k_{\max}} = \left[ \begin{array}{ccc} V_1 & \dots & V_{k_{\max}} \end{array} \right], \quad V_k = \left[ \begin{array}{ccc} & & & & & \\ & & & & & \\ & & & & & \end{array} \right]$$

**Note:** Implementation in real arithmetic possible by combining two steps [B./Li/Penzl '99/'08] or using new idea employing the relation of 2 consecutive complex factors [B./Kürschner/Saak '11].

# Lyapunov equation $0 = AX + XA^T + BB^T$ .

# Algorithm [Penzl '97/'00, Li/White '99/'02, B. 04, B./Li/Penzl '99/'08]

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FOR  $k=2,3,\ldots$ 

$$V_k \leftarrow \sqrt{\frac{\operatorname{re} p_k}{\operatorname{re} p_{k-1}}} \left(V_{k-1} - (p_k + \overline{p_{k-1}})(A+p_kI)^{-1}V_{k-1}\right)$$

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$$Z_{k_{\max}} = \begin{bmatrix} V_1 & \dots & V_{k_{\max}} \end{bmatrix}, \quad V_k = \begin{bmatrix} \in \mathbb{C}^{n \times m}. \end{bmatrix}$$

Note: Implementation in real arithmetic possible by combining two steps [B./Li/Penzl '99/'08] or using new idea employing the relation of 2 consecutive complex factors [B./Kürschner/Saak '11].

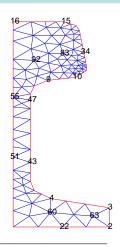
#### **Optimal Cooling of Steel Profiles**

 Mathematical model: boundary control for linearized 2D heat equation.

$$c \cdot \rho \frac{\partial}{\partial t} x = \lambda \Delta x, \quad \xi \in \Omega$$
$$\lambda \frac{\partial}{\partial n} x = \kappa (u_k - x), \quad \xi \in \Gamma_k, \ 1 \le k \le 7,$$
$$\frac{\partial}{\partial n} x = 0, \quad \xi \in \Gamma_7.$$

$$\implies m = 7, q = 6.$$

FEM Discretization, different models for initial mesh (n = 371),
 1, 2, 3, 4 steps of mesh refinement ⇒ n = 1357, 5177, 20209, 79841.



Source: Physical model: courtesy of Mannesmann/Demag.

Math. model: Tröltzsch/Unger 1999/2001, Penzl 1999, Saak 2003.

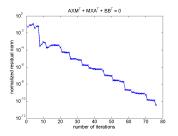
# Numerical Results for ADI

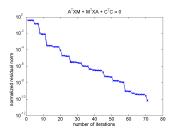
**Optimal Cooling of Steel Profiles** 

Solve dual Lyapunov equations needed for balanced truncation, i.e.,

$$APM^{T} + MPA^{T} + BB^{T} = 0, \quad A^{T}QM + M^{T}QA + C^{T}C = 0,$$
 for  $n = 79.841$ .

- 25 shifts chosen by Penzl heuristic from 50/25 Ritz values of A of largest/smallest magnitude, no column compression performed.
- No factorization of mass matrix required.
- Computations done on Core2Duo at 2.8GHz with 3GB RAM and 32Bit-MATLAB.





CPU times: 626 / 356 sec.

#### Scaling / Mesh Independence

#### Computations by Martin Köhler '10

- $A \in \mathbb{R}^{n \times n} \equiv \text{FDM}$  matrix for 2D heat equation on  $[0, 1]^2$  (LYAPACK benchmark demo\_11, m = 1).
- 16 shifts chosen by Penzl heuristic from 50/25 Ritz values of A of largest/smallest magnitude.
- Computations on 2 dual core Intel Xeon 5160 with 16 GB RAM using M.E.S.S. (http://svncsc.mpi-magdeburg.mpg.de/trac/messtrac/).

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#### CPII Times

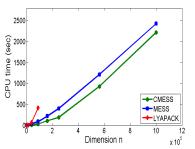
| n         | M.E.S.S. <sup>1</sup> (C) | LyaPack       | M.E.S.S. (MATLAB) |
|-----------|---------------------------|---------------|-------------------|
| 100       | 0.023                     | 0.124         | 0.158             |
| 625       | 0.042                     | 0.104         | 0.227             |
| 2,500     | 0.159                     | 0.702         | 0.989             |
| 10,000    | 0.965                     | 6.22          | 5.644             |
| 40,000    | 11.09                     | 71.48         | 34.55             |
| 90,000    | 34.67                     | 418.5         | 90.49             |
| 160,000   | 109.3                     | out of memory | 219.9             |
| 250,000   | 193.7                     | out of memory | 403.8             |
| 562,500   | 930.1                     | out of memory | 1216.7            |
| 1,000,000 | 2220.0                    | out of memory | 2428.6            |

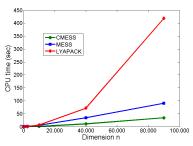
# Numerical Results for ADI

Scaling / Mesh Independence

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- Computations on 2 dual core Intel Xeon 5160 with 16 GB RAM using M.E.S.S. (http://svncsc.mpi-magdeburg.mpg.de/trac/messtrac/).





**Note:** for n=1,000,000, first sparse LU needs  $\sim 1,100$  sec., using UMFPACK this reduces to 30 sec.

Lyapunov equation  $0 = AX + XA^T + BB^T$ 

Projection-based methods for Lyapunov equations with  $A + A^T < 0$ :

- © Compute orthonormal basis range (Z),  $Z \in \mathbb{R}^{n \times r}$ , for subspace  $Z \subset \mathbb{R}^n$ , dim Z = r.
- Solve small-size Lyapunov equation  $\hat{A}\hat{X} + \hat{X}\hat{A}^T + \hat{B}\hat{B}^T = 0$ .
- ① Use  $X \approx Z\hat{X}Z^T$ .

### Examples:

• Krylov subspace methods, i.e., for m = 1:

$$\mathcal{Z} = \mathcal{K}(A, B, r) = \operatorname{span}\{B, AB, A^2B, \dots, A^{r-1}B\}$$

[Saad '90, Jaimoukha/Kasenally '94, Jbilou '02-'08].

• K-PIK [SIMONCINI '07],

$$\mathcal{Z} = \mathcal{K}(A, B, r) \cup \mathcal{K}(A^{-1}, B, r).$$

• Rational Krylov [Druskin/Simoncini '11] (→ exercises).

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- Solve small-size Lyapunov equation  $\hat{A}\hat{X} + \hat{X}\hat{A}^T + \hat{B}\hat{B}^T = 0$ .
- Use  $X \approx Z\hat{X}Z^T$ .

### Examples:

• ADI subspace [B./R.-C. LI/TRUHAR '08]:

$$\mathcal{Z} = \operatorname{colspan} \left[ \begin{array}{ccc} V_1, & \dots, & V_r \end{array} \right].$$

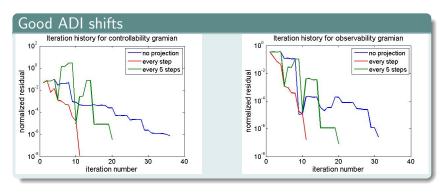
#### Note:

- ADI subspace is rational Krylov subspace [J.-R. LI/WHITE '02].
- Similar approach: ADI-preconditioned global Arnoldi method [JBILOU '08].

### Numerical Methods for Solving Lyapunov Equations Numerical examples for Galerkin-ADI

FEM semi-discretized control problem for parabolic PDE:

- optimal cooling of rail profiles,
- n = 20,209, m = 7, q = 6.

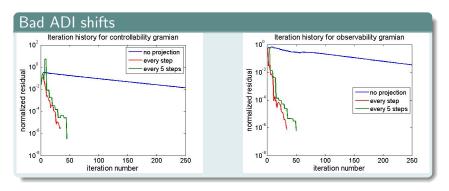


CPU times: 80s (projection every 5th ADI step) vs. 94s (no projection).

# Numerical Methods for Solving Lyapunov Equations Numerical examples for Galerkin-ADI

FEM semi-discretized control problem for parabolic PDE:

- optimal cooling of rail profiles,
- n = 20,209, m = 7, q = 6.

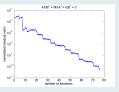


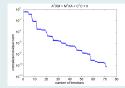
CPU times: 368s (projection every 5th ADI step) vs. 1207s (no projection).

## **Numerical Methods for Solving Lyapunov Equations**

Numerical examples for Galerkin-ADI: optimal cooling of rail profiles, n = 79,841.

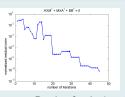
## M.E.S.S. w/o Galerkin projection and column compression

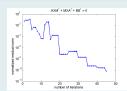




Rank of solution factors: 532 / 426

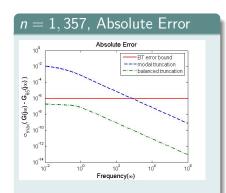
## M.E.S.S. with Galerkin projection and column compression





Rank of solution factors: 269 / 205

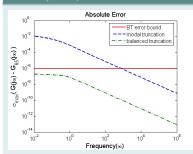
Numerical example for BT: Optimal Cooling of Steel Profiles



- BT model computed with sign function method,
- MT w/o static condensation, same order as BT model.

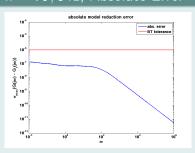
Numerical example for BT: Optimal Cooling of Steel Profiles

## n = 1,357, Absolute Error



- BT model computed with sign function method,
- MT w/o static condensation, same order as BT model.

## n = 79,841, Absolute Error



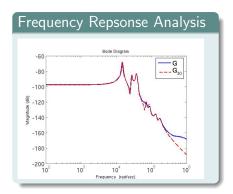
- BT model computed using M.E.S.S. in MATLAB,
- dualcore, computation time:<10 min.</li>

### Numerical example for BT: Microgyroscope (Butterfly Gyro)

- FEM discretization of structure dynamical model using quadratic tetrahedral elements (ANSYS-SOLID187)
  - $\rightarrow$  n = 34,722, m = 1, q = 12.
- Reduced model computed using SPARED, r = 30.

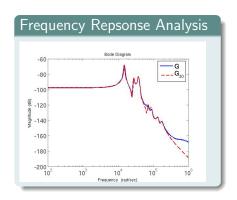
Numerical example for BT: Microgyroscope (Butterfly Gyro)

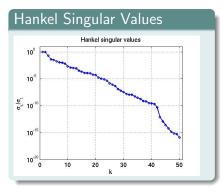
- Reduced model computed using SPARED, r = 30.



Numerical example for BT: Microgyroscope (Butterfly Gyro)

- Reduced model computed using SPARED, r = 30.





Theory [Lancaster/Rodman '95]

### Theorem

Consider the (continuous-time) algebraic Riccati equation (ARE)

$$0 = \mathcal{R}(X) = C^T C + A^T X + XA - XBB^T X,$$

with  $A \in \mathbb{R}^{n \times n}$ ,  $B \in \mathbb{R}^{n \times m}$ ,  $C \in \mathbb{R}^{q \times n}$ , (A, B) stabilizable, (A, C) detectable. Then:

- (a) There exists a unique stabilizing  $X_* \in \{X \in \mathbb{R}^{n \times n} \mid \mathcal{R}(X) = 0\}$ , i.e.,  $\Lambda(A BB^T X_*) \in \mathbb{C}^-$ .
- (b)  $X_* = X_*^T \ge 0$  and  $X_* \ge X$  for all  $X \in \{X \in \mathbb{R}^{n \times n} \mid \mathcal{R}(X) = 0\}$ .
- (c) If (A, C) observable, then  $X_* > 0$ .
- (d) span  $\left\{ \begin{bmatrix} I_n \\ -X_* \end{bmatrix} \right\}$  is the unique invariant subspace of the Hamiltonian matrix

$$H = \left[ \begin{array}{cc} A & BB^T \\ C^T C & -A^T \end{array} \right]$$

corresponding to  $\Lambda(H) \cap \mathbb{C}^-$ .

Numerical Methods [Bini/lannazzo/Meini '12]

## Numerical Methods (incomplete list)

- Invariant subspace methods (→ eigenproblem for Hamiltonian matrix):
  - Schur vector method (care) [LAUB '79]
  - Hamiltonian SR algorithm [Bunse-Gerstner/Mehrmann '86]
  - Symplectic URV-based method

[B./Mehrmann/Xu '97/'98, Chu/Liu/Mehrmann '07]

- Spectral projection methods
  - Sign function method [ROBERTS '71, BYERS '87]
  - Disk function method [Bai/Demmel/Gu '94, B. '97]
- (rational, global) Krylov subspace techniques
  [JAIMOUKHA/KASENALLY '94, JBILOU '03/'06, HEYOUNI/JBILOU '09]
- Newton's method
  - Kleinman iteration [Kleinman '68]
  - Line search acceleration [B./Byers '98]
  - Newton-ADI [B./J.-R. LI/PENZL '99/'08]
  - Inexact Newton [Feitzinger/Hylla/Sachs '09]

Newton's Method for AREs

[Kleinman '68, Mehrmann '91, Lancaster/Rodman '95, B./Byers '94/'98, B. '97, Guo/Laub '99]

- Consider  $0 = \mathcal{R}(X) = C^T C + A^T X + XA XBB^T X$ .
- Frechét derivative of  $\mathcal{R}(X)$  at X:

$$\mathcal{R}_{\mathbf{X}}^{'}: \mathbf{Z} \to (\mathbf{A} - \mathbf{B}\mathbf{B}^{\mathsf{T}}\mathbf{X})^{\mathsf{T}}\mathbf{Z} + \mathbf{Z}(\mathbf{A} - \mathbf{B}\mathbf{B}^{\mathsf{T}}\mathbf{X}).$$

Newton-Kantorovich method:

$$X_{j+1} = X_j - \left(\mathcal{R}'_{X_j}\right)^{-1} \mathcal{R}(X_j), \quad j = 0, 1, 2, \dots$$

## Newton's method (with line search) for AREs

FOR j = 0, 1, ...

- ② Solve the Lyapunov equation  $A_i^T N_j + N_j A_j = -\mathcal{R}(X_j)$ .
- $3 X_{j+1} \leftarrow X_j + t_j N_j.$

FND FOR i

Newton's Method for AREs

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END FOR i

## Newton's Method for AREs

Properties and Implementation

- Convergence for  $K_0$  stabilizing:
  - $A_i = A BK_i = A BB^T X_i$  is stable  $\forall j \geq 0$ .
  - $\lim_{i\to\infty} \|\mathcal{R}(X_i)\|_F = 0$  (monotonically).
  - $\lim_{i\to\infty} X_i = X_* \ge 0$  (locally quadratic).
- Need large-scale Lyapunov solver; here, ADI iteration: linear systems with dense, but "sparse+low rank" coefficient matrix A<sub>i</sub>:

$$A_j = A - B \cdot K_j$$
 $= \text{sparse} - m \cdot$ 

•  $m \ll n \Longrightarrow$  efficient "inversion" using Sherman-Morrison-Woodbury formula:

$$(A-BK_j+\rho_k^{(j)}I)^{-1}=(I_n+(A+\rho_k^{(j)}I)^{-1}B(I_m-K_j(A+\rho_k^{(j)}I)^{-1}B)^{-1}K_j)(A+\rho_k^{(j)}I)^{-1}.$$

## Newton's Method for AREs

Properties and Implementation

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 m ≪ n ⇒ efficient "inversion" using Sherman-Morrison-Woodbury formula:

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**Properties and Implementation** 

- Convergence for  $K_0$  stabilizing:
  - $A_i = A BK_i = A BB^T X_i$  is stable  $\forall i > 0$ .
  - $\lim_{i\to\infty} \|\mathcal{R}(X_i)\|_F = 0$  (monotonically).
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$$(A - BK_j + p_k^{(j)}I)^{-1} = (I_n + (A + p_k^{(j)}I)^{-1}B(I_m - K_j(A + p_k^{(j)}I)^{-1}B)^{-1}K_j)(A + p_k^{(j)}I)^{-1}.$$

## Low-Rank Newton-ADI for AREs

Re-write Newton's method for AREs

$$A_{j}^{T} N_{j} + N_{j} A_{j} = -\mathcal{R}(X_{j})$$

$$\iff$$

$$A_{j}^{T} \underbrace{(X_{j} + N_{j})}_{=X_{j+1}} + \underbrace{(X_{j} + N_{j})}_{=X_{j+1}} A_{j} = \underbrace{-C^{T}C - X_{j}BB^{T}X_{j}}_{=:-W_{j}W_{j}^{T}}$$

$$\text{Set } X_{j} = Z_{j}Z_{j}^{T} \text{ for } \operatorname{rank}(Z_{j}) \ll n \Longrightarrow$$

$$A_{j}^{T} \underbrace{(Z_{j+1}Z_{j+1}^{T})}_{=:-W_{j}W_{j}^{T}} + \underbrace{(Z_{j+1}Z_{j+1}^{T})}_{=:-W_{j}W_{j}^{T}} A_{j} = -W_{j}W_{j}^{T}$$

## Factored Newton Iteration [B./Li/Penzl 1999/2008]

Solve Lyapunov equations for  $Z_{j+1}$  directly by factored ADI iteration and use 'sparse + low-rank' structure of  $A_i$ .

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$$A_{j}^{T} (Z_{j+1}Z_{j+1}^{T}) + (Z_{j+1}Z_{j+1}^{T}) A_{j} = -W_{j}W_{j}^{T}$$

## Factored Newton Iteration [B./Li/Penzl 1999/2008]

Solve Lyapunov equations for  $Z_{i+1}$  directly by factored ADI iteration and use 'sparse + low-rank' structure of  $A_i$ .

#### Feedback Iteration

Optimal feedback

$$K_* = B^T X_* = B^T Z_* Z_*^T$$

can be computed by direct feedback iteration:

• *j*th Newton iteration:

$$K_j = B^T Z_j Z_j^T = \sum_{k=1}^{k_{\mathsf{max}}} (B^T V_{j,k}) V_{j,k}^T \quad \xrightarrow{j \to \infty} \quad K_* = B^T Z_* Z_*^T$$

•  $K_j$  can be updated in ADI iteration, no need to even form  $Z_j$ , need only fixed workspace for  $K_j \in \mathbb{R}^{m \times n}$ !

Related to earlier work by [Banks/Ito 1991].

Galerkin-Newton-ADI

### Basic ideas

- Hybrid method of Galerkin projection methods for AREs
   [Jaimoukha/Kasenally '94, Jbilou '06, Heyouni/Jbilou '09]
   and Newton-ADI, i.e., use column space of current Newton iterate
   for projection, solve projected ARE, and prolongate.

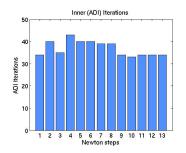
Galerkin-Newton-ADI

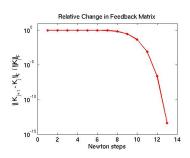
### Basic ideas

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   and Newton-ADI, i.e., use column space of current Newton iterate
   for projection, solve projected ARE, and prolongate.

# LQR Problem for 2D Geometry

- Linear 2D heat equation with homogeneous Dirichlet boundary and point control/observation.
- FD discretization on uniform  $150 \times 150$  grid.
- n = 22.500, m = p = 1, 10 shifts for ADI iterations.
- Convergence of large-scale matrix equation solvers:





Newton-ADI vs. Newton-ADI-Gelerkin

- FDM for 2D heat/convection-diffusion equations on  $[0,1]^2$  (LYAPACK benchmarks, m=p=1)  $\leadsto$  symmetric/nonsymmetric  $A \in \mathbb{R}^{n \times n}$ , n=10,000.
- 15 shifts chosen by Penzl's heuristic from 50/25 Ritz/harmonic Ritz values of A.
- Computations using Intel Core 2 Quad CPU of type Q9400 at 2.66GHz with 4 GB RAM and 64Bit-MATLAB.

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# Newton-ADI

| step | ) | rel. change | rel. residual | ADI |
|------|---|-------------|---------------|-----|
| 1    |   | 1           | 9.99e-01      | 200 |
| 2    |   | 9.99e-01    | 3.41e+01      | 23  |
| 3    |   | 5.25e-01    | 6.37e+00      | 20  |
| 4    |   | 5.37e-01    | 1.52e+00      | 20  |
| 5    |   | 7.03e-01    | 2.64e-01      | 23  |
| 6    | , | 5.57e-01    | 1.56e-02      | 23  |
| 7    |   | 6.59e-02    | 6.30e-05      | 23  |
| 8    |   | 4.02e-04    | 9.68e-10      | 23  |
| 9    | ) | 8.45e-09    | 1.09e-11      | 23  |
| 10   | ) | 1.52e–14    | 1.09e-11      | 23  |

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| step | rel. change | rel. residual | ADI |
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| 1    | 1           | 9.99e-01      | 200 |
| 2    | 9.99e-01    | 3.41e+01      | 23  |
| 3    | 5.25e-01    | 6.37e+00      | 20  |
| 4    | 5.37e-01    | 1.52e+00      | 20  |
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| 7    | 6.59e-02    | 6.30e-05      | 23  |
| 8    | 4.02e-04    | 9.68e-10      | 23  |
| 9    | 8.45e-09    | 1.09e-11      | 23  |
| 10   | 1.52e–14    | 1.09e-11      | 23  |

76.9 sec

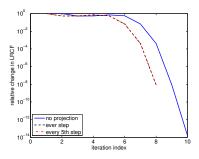
CPU time.

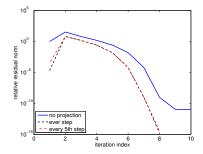
### Newton-Galerkin-ADI

| ste | ) | rel. change | rel. residual | ADI |
|-----|---|-------------|---------------|-----|
|     | Ĺ | 1           | 3.56e-04      | 20  |
| - 2 | 2 | 5.25e-01    | 6.37e+00      | 10  |
| 3   | 3 | 5.37e-01    | 1.52e+00      | 6   |
| 4   | 1 | 7.03e-01    | 2.64e-01      | 10  |
| į   | 5 | 5.57e-01    | 1.57e-02      | 10  |
| (   | ŝ | 6.59e-02    | 6.30e-05      | 10  |
|     | 7 | 4.03e-04    | 9.79e-10      | 10  |
| 8   | 3 | 8.45e-09    | 1.43e-15      | 10  |

CPU time: 38.0 sec

- FDM for 2D heat/convection-diffusion equations on  $[0,1]^2$  (LYAPACK benchmarks, m=p=1)  $\leadsto$  symmetric/nonsymmetric  $A \in \mathbb{R}^{n \times n}$ , n=10,000.
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Newton-ADI vs. Newton-ADI-Gelerkin

- FDM for 2D heat/convection-diffusion equations on  $[0,1]^2$  (LYAPACK benchmarks, m=p=1)  $\leadsto$  symmetric/nonsymmetric  $A \in \mathbb{R}^{n \times n}$ , n=10,000.
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- Computations using Intel Core 2 Quad CPU of type Q9400 at 2.66GHz with 4 GB RAM and 64Bit-MATLAB.

| Newton-ADI |             |               |     |
|------------|-------------|---------------|-----|
| step       | rel. change | rel. residual | ADI |
| 1          | 1           | 9.99e-01      | 200 |
| 2          | 9.99e-01    | 3.56e+01      | 60  |
| 3          | 3.11e-01    | 3.72e+00      | 39  |
| 4          | 2.88e-01    | 9.62e-01      | 40  |
| 5          | 3.41e-01    | 1.68e-01      | 45  |
| 6          | 1.22e-01    | 5.25e-03      | 42  |
| 7          | 3.88e-03    | 2.96e-06      | 47  |
| 8          | 2.30e-06    | 6.09e-13      | 47  |
|            | CPU time:   | 185.9 sec.    |     |

- FDM for 2D heat/convection-diffusion equations on  $[0,1]^2$  (LYAPACK benchmarks, m=p=1)  $\leadsto$  symmetric/nonsymmetric  $A \in \mathbb{R}^{n \times n}$ , n=10,000.
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## Newton-ADI

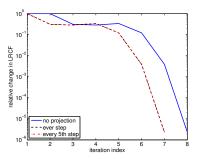
| step | rel. change | rel. residual | ADI |
|------|-------------|---------------|-----|
| 1    | 1           | 9.99e-01      | 200 |
| 2    | 9.99e-01    | 3.56e+01      | 60  |
| 3    | 3.11e-01    | 3.72e+00      | 39  |
| 4    | 2.88e-01    | 9.62e-01      | 40  |
| 5    | 3.41e-01    | 1.68e-01      | 45  |
| 6    | 1.22e-01    | 5.25e-03      | 42  |
| 7    | 3.88e-03    | 2.96e-06      | 47  |
| 8    | 2.30e-06    | 6.09e-13      | 47  |
|      | CPU time:   | 185.9 sec.    |     |

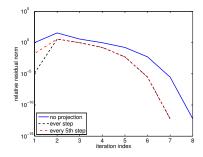
### Newton-Galerkin-ADI

| step | rel. change | rel. residual | ADI it. |
|------|-------------|---------------|---------|
| 1    | 1           | 1.78e-02      | 35      |
| 2    | 3.11e-01    | 3.72e+00      | 15      |
| 3    | 2.88e-01    | 9.62e-01      | 20      |
| 4    | 3.41e-01    | 1.68e-01      | 15      |
| 5    | 1.22e-01    | 5.25e-03      | 20      |
| 6    | 3.89e-03    | 2.96e-06      | 15      |
| 7    | 2.30e-06    | 6.14e-13      | 20      |
|      |             | ı             | •       |
|      |             |               |         |

CPU time: 75.7 sec.

- FDM for 2D heat/convection-diffusion equations on  $[0,1]^2$  (LYAPACK benchmarks, m=p=1)  $\leadsto$  symmetric/nonsymmetric  $A \in \mathbb{R}^{n \times n}$ , n=10,000.
- 15 shifts chosen by Penzl's heuristic from 50/25 Ritz/harmonic Ritz values of A.
- Computations using Intel Core 2 Quad CPU of type Q9400 at 2.66GHz with 4 GB RAM and 64Bit-MATLAB.





Example: LQR Problem for 3D Geometry

## Control problem for 3d Convection-Diffusion Equation

- FDM for 3D convection-diffusion equation on [0, 1]<sup>3</sup>
- proposed in [Simoncini '07], q = p = 1
- non-symmetric  $A \in \mathbb{R}^{n \times n}$  , n = 10648

### Test system:

INTEL Xeon 5160 3.00GHz; 16 GB RAM; 64Bit-MATLAB (R2010a) using threaded BLAS; stopping tolerance: 10<sup>-10</sup>

Example: LQR Problem for 3D Geometry

### Newton-ADI

| NWT | rel. change          | rel. residual        | ADI |
|-----|----------------------|----------------------|-----|
| 1   | $1.0 \cdot 10^{0}$   | $9.3 \cdot 10^{-01}$ | 100 |
| 2   | $3.7 \cdot 10^{-02}$ | $9.6 \cdot 10^{-02}$ | 94  |
| 3   | $1.4 \cdot 10^{-02}$ | $1.1 \cdot 10^{-03}$ | 98  |
| 4   | $3.5 \cdot 10^{-04}$ | $1.0 \cdot 10^{-07}$ | 97  |
| 5   | $6.4 \cdot 10^{-08}$ | $1.3 \cdot 10^{-10}$ | 97  |
| 6   | $7.5 \cdot 10^{-16}$ | $1.3 \cdot 10^{-10}$ | 97  |

CPU time: 4805.8 sec.

#### NG-ADI inner = 5, outer = 1

|   |                    | rel. residual        |    |
|---|--------------------|----------------------|----|
| 1 | $1.0 \cdot 10^{0}$ | $5.0 \cdot 10^{-11}$ | 80 |
|   |                    | : 497.6 sec.         |    |

#### NG-ADI inner = 1, outer = 1

| NWT                  |                    | rel. residual        | ADI |
|----------------------|--------------------|----------------------|-----|
| 1                    | $1.0 \cdot 10^{0}$ | $7.4 \cdot 10^{-11}$ | 71  |
| CPU time: 856.6 sec. |                    |                      |     |

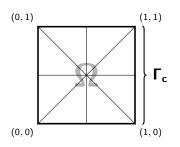
#### NG-ADI inner=0, outer=1

|                      | NWT | rel. change        | rel. residual        | ADI |
|----------------------|-----|--------------------|----------------------|-----|
| ĺ                    | 1   | $1.0 \cdot 10^{0}$ | $6.5 \cdot 10^{-13}$ | 100 |
| CPU time: 506.6 sec. |     |                    |                      |     |

#### Test system:

INTEL Xeon 5160 3.00GHz; 16 GB RAM; 64Bit-MATLAB (R2010a) using threaded BLAS; stopping tolerance: 10<sup>-10</sup>

Scaling of CPU times / Mesh Independence



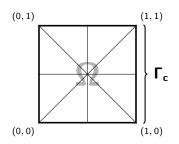
$$\begin{split} \partial_t x(\xi,t) &= \Delta x(\xi,t) &&\text{in } \Omega \\ \partial_\nu x &= b(\xi) \cdot u(t) - x &&\text{on } \Gamma_c \\ \partial_\nu x &= -x &&\text{on } \partial\Omega \setminus \Gamma_c \\ x(\xi,0) &= 1 \end{split}$$

#### Note:

Here  $b(\xi) = 4(1-\xi_2)\xi_2$  for  $\xi \in \Gamma_c$  and 0 otherwise, thus  $\forall t \in \mathbb{R}_{>0}$ , we have  $u(t) \in \mathbb{R}$ .

$$\Rightarrow B_h = M_{\Gamma,h} \cdot b.$$

Scaling of CPU times / Mesh Independence



$$\partial_t x(\xi, t) = \Delta x(\xi, t)$$
 in  $\Omega$ 

$$\partial_\nu x = b(\xi) \cdot u(t) - x$$
 on  $\Gamma_c$ 

$$\partial_\nu x = -x$$
 on  $\partial\Omega \setminus \Gamma_c$ 

$$x(\xi, 0) = 1$$

**Consider:** output equation y = Cx, where

$$\begin{array}{ccc} C: \mathcal{L}^2(\Omega) & \to \mathbb{R} \\ x(\xi,t) & \mapsto y(t) = \int_{\Omega} x(\xi,t) \, d\xi \end{array} \Rightarrow C_h = \underline{1} \cdot M_h.$$

RatInt Balanced Truncation Matrix Equations

## Numerical Results

Scaling of CPU times / Mesh Independence

## Simplified Low Rank Newton-Galerkin ADI

- generalized state space form implementation
- Penzl shifts (16/50/25) with respect to initial matrices
- projection acceleration in every outer iteration step
- projection acceleration in every 5-th inner iteration step

## Test system:

INTEL Xeon 5160 @ 3.00 GHz; 16 GB RAM; 64Bit-MATLAB (R2010a) using threaded BLAS,

stopping criterion tolerances:  $10^{-10}$ 

## **Computation Times**

| discretization level | problem size | time in seconds      |
|----------------------|--------------|----------------------|
| 3                    | 81           | $4.87 \cdot 10^{-2}$ |
| 4                    | 289          | $2.81 \cdot 10^{-1}$ |
| 5                    | 1 089        | $5.87 \cdot 10^{-1}$ |
| 6                    | 4 225        | 2.63                 |
| 7                    | 16 641       | $2.03 \cdot 10^{+1}$ |
| 8                    | 66 049       | $1.22 \cdot 10^{+2}$ |
| 9                    | 263 169      | $1.05 \cdot 10^{+3}$ |
| 10                   | 1 050 625    | $1.65 \cdot 10^{+4}$ |
| 11                   | 4 198 401    | $1.35 \cdot 10^{+5}$ |

### Test system:

INTEL Xeon 5160 @ 3.00 GHz; 16 GB RAM; 64Bit-MATLAB (R2010a) using threaded BLAS,

stopping criterion tolerances:  $10^{-10}$ 

## **Solving Large-Scale Matrix Equations**

Software

Lyapack [Penzl 2000]

MATLAB toolbox for solving

- Lyapunov equations and algebraic Riccati equations,
- model reduction and LQR problems.

Main work horse: Low-rank ADI and Newton-ADI iterations.

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- Extensions to bilinear and stochastic systems.
- Rational interpolation methods for nonlinear systems.
- Other MOR techniques like POD, RB.
- MOR methods for discrete-time systems.
- Extensions to descriptor systems  $E\dot{x} = Ax + Bu$ , E singular.
- Parametric model reduction:

$$\dot{x} = A(p)x + B(p)u, \quad y = C(p)x,$$

where  $p \in \mathbb{R}^d$  is a free parameter vector; parameters should be preserved in the reduced-order model.

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