# Second Order Reverse Mode of AD : A Vertex Elimination Perspective 

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## PURDUE

## Outline

- Second order reverse mode of Automatic Differentiation
- Vertex elimination for evaluating the Gradient and the Hessian
- The correspondence between second order reverse mode and vertex elimination
- Discussion and board picture


## AD Fundamentals

- Automatic Differentiation (AD) is a technique that augments a computer program so that the augmented program computes the derivatives as well as the values of the function defined by the original program.
Scalar Objective Function $f: \mathcal{R}^{n} \rightarrow \mathcal{R}^{1}$
Implemented as a computer program
The evaluation is on a sequence of decomposed elemental functions For $k=1,2$,


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v_{k}=\varphi_{k}\left(v_{i}\right)_{\left\{v_{i}: v_{i}<v_{k}\right\}}
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$\begin{aligned}-\mathrm{y} & =\operatorname{pow}(\operatorname{pow}(x \\ & \nabla v_{0} \ll=x\end{aligned}$
$\left(x>0, y=x^{4 x}\right)$

- $v_{1}=\varphi_{1}\left(v_{0}\right)=v_{0} * v_{0}$
- $v_{2}=\varphi_{2}\left(v_{1}\right)=\operatorname{pow}\left(v_{1}, 2.0\right)$
- $v_{3}=\varphi_{3}\left(v_{2}, v_{0}\right)=\operatorname{pow}\left(v_{2}, v_{0}\right)$
- $v_{3} \gg=y$


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- Indexing convention :
- Independent variables: $v_{1-n}, \cdots, v_{0}$
- Intermediate variables: $v_{1}, \cdots, v_{l-1}$
- Dependent variable : $v_{l}$


## Second Order Reverse Mode : Story Line

- First Proposed by Gower and Mello ${ }^{1}$
- Called Edge_Pushing initially
- From the closed form of second order derivative for composite functions

Wang, Gebremedhin, and Pothen provided a second perspective by
adopting live variable analysis ${ }^{2}$ from compiler theory.
Better complexity bound
Correct Implementation
Further improved with preaccumulation
The new proof can be extended into general high orders.

[^0]Wang, Mu, Assefaw Gebremedhin, and Alex Pothen.
Capitalizing on live variables: new algorithms for efficient Hessian computation via automatic differentiation.
Mathematical Programming Computation (2016): 1-41

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## Reverse Mode of AD

- Function evaluation : evaluate each elemental function

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\text { for } \begin{aligned}
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Reverse mode of AD : process sequence of elemental functions in reverse order

do something with $v_{k}=\varphi_{k}\left(v_{i}\right)_{\left\{v_{i} \nprec v_{k}\right.}$
Equivalent function $f_{k}\left(S_{k}\right)$ : a function defined by the elemental
functions $\varphi_{1}, \cdots, \varphi_{k}$ that have been processed at the end of step $k$,
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- Equivalent function $f_{k}\left(S_{k}\right)$ : a function defined by the elemental functions $\varphi_{1}, \cdots, \varphi_{k}$ that have been processed at the end of step $k$, in reverse mode
- $f=\underbrace{\varphi_{1} \circ \cdots \circ \varphi_{k}}_{f_{k}\left(S_{k}\right)} \circ \varphi_{k-1} \circ \cdots \circ \varphi_{1}$.
- The independent variables of $f_{k}$ are denoted by $S_{k}$.


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& f=\overbrace{\varphi_{1} \circ \cdots \circ \varphi_{k+1} \circ \varphi_{k} \circ \varphi_{k-1} \circ \cdots \circ \varphi_{1}}^{f_{k+1}\left(S_{k+1}\right)} \Rightarrow \underbrace{\varphi_{1} \circ \cdots \circ \varphi_{k+1} \circ \varphi_{k}}_{f_{k}\left(S_{k}\right)} \circ \varphi_{k-1} \circ \cdots \circ \varphi_{1}
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- First order chain rule : $\frac{\partial f_{k}}{\partial v_{i}}=\frac{\partial f_{k+1}}{\partial v_{i}}+\frac{\partial v_{k}}{\partial v_{i}} \frac{\partial f_{k+1}}{\partial v_{k}}$


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- Second order chain rule :

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\frac{\partial^{2} f_{k}}{\partial v_{i} \partial v_{j}} & =\frac{\partial^{2} f_{k+1}}{\partial v \partial u}+\frac{\partial v_{k}}{\partial v_{k}} \frac{\partial^{2} f_{k+1}}{\partial v_{j} \partial v_{k}}+\frac{\partial v_{k}}{\partial v_{j}} \frac{\partial^{2} f_{k+1}}{\partial v_{i} \partial v_{k}} \\
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- Adjoint variable $\bar{v}_{i}$ :
- Holds the value of $\frac{\partial f_{k}}{\partial v_{i}}$ after the step $k$
- Incremental updates in implementation


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- Incremental updates in implementation
- More implementation details for second order for exploiting sparsity and symmetry.


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- General high order chain rule $\rightarrow$ general high order reverse mode
- Taking advantage of symmetry becomes more important


## Reverse Mode of AD : Implementation

- Second order reverse mode : Initially implemented as LivarH in ADOL-C
- https://github.com/CSCsw/LivarH

ReverseAD : an operator overloading implementation of general high order reverse mode in $\mathrm{C}++11$.
https://github com/wangmu0701/ReverseAD
Available for experimentation
Monotonic indexing for variables on the trace $v_{i} \prec v_{k} \Longrightarrow \operatorname{index}\left(v_{i}\right)<\operatorname{index}\left(v_{j}\right)$

Not satisfied by ADOL-C
An immature fix was provided for LivarH

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## Reverse Mode of AD : Performance

- The FeasNewt Benchmark (T. S. Munson and P. D. Hovland, 2005)
- A mesh optimization problem with sparse Hessian matrix.
- Compared with compression-and-recovery approach implemented in ADOL-C + ColPack

| $n:$ |  | 2,598 | 12,597 | 39,379 |
| :---: | :---: | ---: | ---: | ---: |
| $\# n n z$ in $H:$ |  | 46,488 | 253,029 | 828,129 |
| Direct | \#colors : | 54 | 62 | 65 |
|  | runtime(S) : | 3.77 | 39.34 | 137.07 |
| Indirect | \#colors : | 31 | 30 | 31 |
|  | runtime(S) : | 3.56 | 31.07 | 119.04 |
| ReverseAD | runtime(S) : | 0.51 | 3.37 | 12.40 |

## From Analytical to Combinatorial

- The second (high) order reverse mode is derived from a purely analytical point of view.
- Same as the original derivation of Edge Pushing.

There are combinatorial models for AD algorithms based on the concept of Computational Graph $G$ of the objective function.

Edge Elimination
Vertex Elimination
Face Elimination
Closely related to the classical linear algebra problem of sparse Gaussian elimination

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## Computational Graph

- Computational graph: $G=(V, E)$
- Variables are vertices: $V=\left\{v_{i} \mid 1-n \leq i \leq I\right\}$
- Precedence relations are directed edges:

$$
E=\left\{v_{i} \rightarrow v_{k} \mid v_{i} \prec v_{k}, 1-n \leq i<k \leq I\right\}
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- Edge weights : $c(i, k) \doteq w\left(v_{i}, v_{k}\right)=\frac{\partial v_{k}}{\partial v_{i}}$


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$\Rightarrow v_{3}=\varphi_{3}\left(v_{2}, v_{0}\right)=\operatorname{pow}\left(v_{2}, v_{0}\right)$



## Vertex Elimination

## Repeat

- Pick intermediate node $v_{j}$
- For all $(i, k)$, s.t, $i \prec j \prec k$ do

$$
c(i, k)+=c(i, j) * c(j, k)
$$

$\checkmark$ Remove $v_{j}$ from $V$
Until $V$ has no intermediate vertices


- Proposed by Griewank and Reese, and studied extensively by Naumann and students


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- Proposed by Griewank and Reese, and studied extensively by Naumann and students
- Any elimination order will give the same final results.
- The time complexity (number of edge weights computed) varies with the ordering. Minimizing the space complexity also is likely to be intractable.
- NP-hard to determine the optimal ordering.


## Vertex Elimination for Hessian

- The vertex elimination algorithm applies on $G$, gives $\nabla \cdot f$. To evaluate the Hessian of $f$ we need the computational graph of the
gradient $G_{g}$, i.e, the computational graph of evaluating $\nabla \cdot f$.
$G_{g}$ can be constructed from first order non-incremental reverse mode

Function evaluation


First order (nonincremental) reverse mode
Initialize


$\bar{v}_{i}=\sum_{v_{i} \prec v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k}$

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Function evaluation:

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\text { for } \begin{aligned}
k & =1,2, \cdots, l \\
v_{k} & =\varphi_{k}\left(v_{i}\right)_{\left\{v_{i}: v_{i} \prec v_{k}\right\}}
\end{aligned}
$$

First order (nonincremental) reverse mode :
Initialize:

$$
\begin{aligned}
& \quad \bar{v}_{l}=1.0, \bar{v}_{l-1}=\cdots=0 \\
& \text { for } i=I-1, \cdots, 1,0, \cdots, 1-n \\
& \quad \bar{v}_{i}=\sum_{v_{i} \prec v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k} \\
& \quad \bar{v}_{i}=\bar{\varphi}_{i}\left(\cup_{v_{i} \prec v_{k}}\left\{v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right)
\end{aligned}
$$

## Computational Graph of the Gradient

Function evaluation :
$>$ for $k=1,2, \cdots, l$

$$
v_{k}=\varphi_{k}\left(v_{i}\right)_{\left\{v_{i}: v_{i} \prec v_{k}\right\}}
$$

First order (nonincremental) reverse mode :

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$$
\bar{v}_{l}=1.0, \bar{v}_{l-1}=\cdots=0
$$

$\Rightarrow$ for $i=I-1, \cdots, 1,0, \cdots, 1-n$

$$
\bar{v}_{i}=\sum_{v_{i} \prec v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k}
$$

$$
\bar{v}_{i}=\bar{\varphi}_{i}\left(\cup_{v_{i}} \prec v_{k}\left\{v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right)
$$

$\mathrm{V}_{2}$

$$
\overline{V_{g}}=V \cup \bar{V}, E_{g}=E_{G} \cup E_{\bar{G}} \cup E_{C}
$$

$v_{1}$

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$$

First order (nonincremental) reverse mode :

- Initialize:

$$
\bar{v}_{I}=1.0, \bar{v}_{I-1}=\cdots=0
$$

- for $i=I-1, \cdots, 1,0, \cdots, 1-n$

$$
\bar{v}_{i}=\sum_{v_{i}<v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k}
$$

$$
\bar{v}_{i}=\bar{\varphi}_{i}\left(\cup_{v_{i}} \prec v_{k}\left\{v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right)
$$

$V_{g}=V \cup \bar{V}, E_{g}=E_{G} \cup E_{\bar{G}} \cup E_{C}$
$E_{G}:\left(v_{i}, v_{k}\right) \in E_{g} \Longleftrightarrow v_{i} \prec v_{k}$

$$
\begin{aligned}
& v_{k}=\varphi_{k}\left(v_{i}\right)_{\left\{v_{i}: v_{i} \prec v_{k}\right\}} \\
& c(i, k)=\frac{\partial v_{k}}{\partial v_{i}}
\end{aligned}
$$

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Function evaluation :

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First order (nonincremental) reverse mode :

- Initialize:

$$
\bar{v}_{l}=1.0, \bar{v}_{I-1}=\cdots=0
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- for $i=I-1, \cdots, 1,0, \cdots, 1-n$ $\bar{v}_{i}=\sum_{v_{i} \prec v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k}$ $\bar{v}_{i}=\bar{\varphi}_{i}\left(\cup_{v_{i}} \prec v_{k}\left\{v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right)$

$$
\begin{aligned}
& \hline V_{g}=V \cup \bar{v}, E_{g}=E_{G} \cup E_{\bar{G}} \cup E_{C} \\
& E_{\bar{G}}:\left(\bar{v}_{k}, \bar{v}_{i}\right) \in E_{g} \Longleftrightarrow \bar{v}_{k} \prec \bar{v}_{i} \Longleftrightarrow v_{i} \prec v_{k} \\
& \bar{v}_{i}=\sum_{v_{i}}<v_{k} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k} \\
& =\bar{\varphi}_{i}\left(\cup_{v_{i}} \prec v_{k}\right. \\
& \left.\left.\qquad v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right) \\
& c(\bar{k}, \bar{i})=\frac{\partial v_{j}}{\partial v_{i}}
\end{aligned}
$$

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Function evaluation :
$\rightarrow$ for $k=1,2, \cdots, l$

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v_{k}=\varphi_{k}\left(v_{i}\right)_{\left\{v_{i}: v_{i} \prec v_{k}\right\}}
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First order (nonincremental) reverse mode :

- Initialize:

$$
\bar{v}_{l}=1.0, \bar{v}_{l-1}=\cdots=0
$$

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$$
\bar{v}_{i}=\sum_{v_{i}<v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k}
$$

$$
\bar{v}_{i}=\bar{\varphi}_{i}\left(\cup_{v_{i}} \prec v_{k}\left\{v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right)
$$

$$
\begin{aligned}
& V_{g}=V \cup \bar{v}, E_{g}=E_{G} \cup E_{\bar{G}} \cup E_{C} \\
& E_{C}:\left(v_{i}, \bar{v}_{j}\right) \in E_{g} \Longleftrightarrow \exists v_{k}, s . t, v_{i}, v_{j} \prec v_{k} \\
& \bar{v}_{i}=\sum_{v_{i} \prec v_{k}} \frac{\partial v_{k}}{\partial v_{i}} \bar{v}_{k} \\
& =\bar{\varphi}_{i}\left(\cup_{v_{i}} \prec v_{k}\left\{v_{j}: v_{j} \prec v_{k}\right\} \cup\left\{v_{k}\right\}\right) \\
& \qquad c(i, \bar{j})=\sum_{v_{i}, v_{j} \prec v_{k}} \frac{\partial^{2} v_{k}}{\partial v_{i} \partial v_{j}} \bar{v}_{k}
\end{aligned}
$$

## Equivalence

- Vertex elimination on the gradient graph $G_{g}$ gives the Hessian (combinatorial approach).
Second order reverse mode gives the Hessian (analytical approach)
Second order reverse mode:
Vertex Elimination on $G_{g}$
Initialize
Pick intermediate node $v_{j}$

Remove $v_{j}$ from $V$

> Theorem
> If vertex elimination is performed on $G_{g}$ in a symmetric reverse topological ordering, i.e, $\left(v_{k}, \bar{v}_{k}\right)$ are eliminated in pairs, in the order $k=I, I-1, \cdots, 1$, then the two algorithms correspond step-by-step.

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$\qquad$



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## Second order reverse mode:

- Initialize :

$$
\bar{v}_{I}=1.0, \bar{v}_{I-1}=\cdots=0
$$

$>$ for $k=I, I-1, \cdots, 1$

$$
\text { for each unordered pair }\left(v_{i}, v_{j}\right)
$$

$$
h_{k}\left(v_{i}, v_{j}\right)=h_{k+1}\left(v_{i}, v_{j}\right)
$$

$$
+\frac{\partial v_{k}}{\partial v_{i}} h_{k+1}\left(v_{j}, v_{k}\right)+\frac{\partial v_{k}}{\partial v_{j}} h_{k+1}\left(v_{i}, v_{k}\right)
$$

$$
+\frac{\partial v_{k}}{\partial v_{i}} \frac{\partial v_{k}}{\partial v_{j}} h_{k+1}\left(v_{k}, v_{k}\right)+\frac{\partial^{2} v_{k}}{\partial v_{i} \partial v_{j}} \bar{v}_{k}
$$

## Vertex Elimination on $G_{g}$

- Pick intermediate node $v_{j}$
- For all $(i, k)$, s.t, $i \prec j \prec k$ do $c(i, k)+=c(i, j) * c(j, k)$
$\Rightarrow$ Remove $v_{j}$ from $V$
- Repeat until $V$ has no intermediate vertices

If vertex elimination is performed on $G_{g}$ in a symmetric reverse topological ordering, i.e, $\left(v_{k}, \bar{v}_{k}\right)$ are eliminated in pairs, in the order $k=I, I-1, \cdots, 1$, then the two algorithms correspond step-by-step

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Second order reverse mode:

$$
\begin{aligned}
& \quad \text { Initialize : } \\
& \bar{v}_{l}=1.0, \bar{v}_{l-1}=\cdots=0 \\
& \text { for } k=I, I-1, \cdots, 1 \\
& \quad \text { for each unordered pair }\left(v_{i}, v_{j}\right) \\
& \quad h_{k}\left(v_{i}, v_{j}\right)=h_{k+1}\left(v_{i}, v_{j}\right) \\
& \quad+\frac{\partial v_{k}}{\partial v_{i}} h_{k+1}\left(v_{j}, v_{k}\right)+\frac{\partial v_{k}}{\partial v_{j}} h_{k+1}\left(v_{i}, v_{k}\right) \\
& \quad+\frac{\partial v_{k}}{\partial v_{i}} \frac{\partial v_{k}}{\partial v_{j}} h_{k+1}\left(v_{k}, v_{k}\right)+\frac{\partial^{2} v_{k}}{\partial v_{i} \partial v_{j}} \bar{v}_{k} \\
& \hline
\end{aligned}
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## Theorem

- The two algorithms perform the same computations, and thus maintain the same intermediate results after each step.
- With two minor tweaks of vertex elimination on $G_{g}$

Tweak one : parallel edges in $E_{C}$
Break the edge $c(i, \bar{j})=\sum_{v_{i}, v_{j}<v_{k}} \frac{\partial^{2} v_{k}}{\partial v_{i} v_{j}} \bar{v}_{k}$
Into parallel edges $c^{k}(i, \bar{j})=\frac{\partial^{2} v_{k}}{\partial v \partial_{v}} \bar{v}_{k}$
Tweak two : new set of edges $E_{H}$
Rule 1: all added edges are added into $E_{H}$
Rule 2 : After eliminating $\left(v_{k}, \bar{v}_{k}\right)$, move all $c^{k}(i, \bar{j})$ from $E_{C}$ to $E_{H}$
Claim : $E_{H}$ corresponds to the nonzeros in the Hessian of $f_{k}\left(S_{k}\right)$ after each step.

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# Tweak two : new set of edges $E_{H}$ <br> Rule 1: all added edges are added into $E_{H}$ <br> Rule 2 : After eliminating $\left(v_{k}, \bar{v}_{k}\right)$, move all $c^{k}(i, \bar{j})$ from $E_{C}$ to $E_{H}$ 

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## Discussion

- Second order reverse mode is equivalent to a special form of vertex elimination on the computational graph of the gradient $G_{g}$. May not be the optimal form of vertex elimination due to the structure of $G_{g}$. But, in practice it can be implemented with efficient
storage and memory access.
Second order reverse mode does not require the graph $G_{g}$ to be formed.
Can be implemented with a single reverse sweep.
Can incorporate checkpointing to overcome memory/disk limits
Possibilities of optimizing second order reverse mode by exploiting
structural properties
Out-of-order processing of $v_{k}=\varphi_{k}\left(v_{i}\right)\left\{v_{i}: v_{i} \& v_{k}\right\}$
Benefit must outweigh the optimization overhead


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## Future Work : Broad Picture

- This work reveals the correspondence between analytical and combinatorial points of view of AD algorithms.
- First order forward/reverse mode corresponds to edge elimination on $G$ with specific elimination ordering.
Second order reverse mode corresponds to vertex elimination on $G_{g}$
with reverse symmetric elimination ordering.
Is there a generalization to high orders?
The analytical form of the high order reverse mode is the implementation of high order chain rule. What is the generalization of the combinatorial form of high order reverse mode?
What is the computational graph of the Hessian $G_{H}$ ?
What is the elimination technique that we should perform on $G_{H}$ ?


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## References

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## Backup Slides

placeholder

## Vertex Elimination as Gaussian Elimination

- We can build a matrix as $\mathrm{C}=\left[c_{i j}\right]_{1-n \leq i, j \leq 1}$.
- $c_{i j}=\frac{\partial v_{i}}{\partial v_{j}}$ as the edge weight in $G$,when $v_{j} \prec v_{i}$
- $c_{i j}=-1$, diagonal elements
- Other elements are zero

$$
\left.\begin{array}{c}
n \\
\mathrm{C}=\begin{array}{c}
n \\
m
\end{array}
\end{array} \begin{array}{ccc}
n-m & m \\
-\mathrm{I} & 0 & 0 \\
\mathrm{~B} & \mathrm{~L}-\mathrm{I} & 0 \\
\mathrm{R} & \mathrm{~T} & -\mathrm{I}
\end{array}\right]
$$

- C is a lower triangular matrix
- The Jacobian $\nabla \cdot f=\mathrm{R}+\mathrm{T} \cdot(\mathrm{L}-\mathrm{I})^{-1} \cdot \mathrm{~B}$ is the Schur complement
- Can use a Gaussian elimination procedure to compute it.


## Adjacency Matrix for $G_{g}$

$$
\mathrm{H}=\begin{gathered}
\\
n \\
I-m \\
m \\
m \\
I-m \\
n
\end{gathered}\left[\begin{array}{cccccc}
n & I-m & m & m & I-m & n \\
-\mathrm{I} & 0 & 0 & & & \\
\mathrm{~B} & \mathrm{~L}-\mathrm{I} & 0 & & & \\
\mathrm{R} & \mathrm{~T} & -\mathrm{I} & & & \\
0 & 0 & 0 & -I & 0 & 0 \\
\mathrm{Z} & \mathrm{Y} & 0 & \mathrm{~T}^{\prime} & \mathrm{L}^{\prime}-\mathrm{I} & 0 \\
\mathrm{X} & \mathrm{Z}^{\prime} & 0 & \mathrm{R}^{\prime} & \mathrm{B}^{\prime} & -\mathrm{I}
\end{array}\right]
$$

- $\mathrm{C}^{\prime}$ is the transpose of C along the antidiagonal.
- The Hessian is the Schur complement of X with the rest of the matrix


[^0]:    ${ }^{1}$ Gower, Robert Mansel, and Margarida P. Mello. Hessian matrices via automatic differentiation. Universidade Estadual de Campinas, Instituto de Matemtica, Estatstica e Computao Cientfica, 2010.

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    ${ }^{2}$ Wang, Mu, Assefaw Gebremedhin, and Alex Pothen. "Capitalizing on live variables: new algorithms for efficient Hessian computation via automatic differentiation." Mathematical Programming Computation (2016): 1-41.

[^2]:    ${ }^{1}$ Gower, Robert Mansel, and Margarida P. Mello. Hessian matrices via automatic differentiation. Universidade Estadual de Campinas, Instituto de Matemtica, Estatstica e Computao Cientfica, 2010.
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