SODA is jointly sponsored by the ACM Special Interest Group on Algorithms and Computation Theory and the SIAM Activity Group on Discrete Mathematics.

The SIAG on Discrete Mathematics focuses on combinatorics, graph theory, cryptography, discrete optimization, mathematical programming, coding theory, information theory, game theory, and theoretical computer science, including algorithms, complexity, circuit design, robotics, and parallel processing. This activity group provides an opportunity to unify pure discrete mathematics and areas of applied research such as computer science, operations research, combinatorics, and the social sciences. It organizes the SIAM Conference on Discrete Mathematics; co-sponsors, with ACM SIGACT, the annual Symposium on Discrete Algorithms; and sponsors minisymposia at SIAM meetings and conferences. The activity group also runs DM-Net, an electronic forum; publishes an electronic newsletter; and maintains a website and a member directory. Every two years, the activity group also awards the Dénes König Prize to a young researcher for outstanding research in the area of discrete mathematics.

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General Information

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- **Saturday, January 9**
  - 5:00 PM – 8:00 PM

- **Sunday, January 10**
  - 8:00 AM – 5:00 PM

- **Monday, January 11**
  - 8:00 AM – 5:00 PM

- **Tuesday, January 12**
  - 8:00 AM – 5:00 PM

**Hotel Information**

Crystal Gateway Marriott  
1700 Jefferson Davis Highway  
Arlington, Virginia  22202  
USA

**Hotel Telephone Number**

To reach an attendee or leave a message, call +1-703-920-3230. If the attendee is a hotel guest, the hotel operator can connect you with the attendee’s room.
Hotel Check-in and Check-out Times
Check-in time is 4:00 PM.
Check-out time is 12:00 PM.

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The Crystal Gateway Marriott recommends White House Nannies (http://www.whitehousenannies.com/) as a local child care provider.

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SIAM corporate members provide their employees with knowledge about, access to, and contacts in the applied mathematics and computational sciences community through their membership benefits. Corporate membership is more than just a bundle of tangible products and services; it is an expression of support for SIAM and its programs. SIAM is pleased to acknowledge its corporate members and sponsors. In recognition of their support, non-member attendees who are employed by the following organizations are entitled to the SIAM member registration rate.

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List current November 2015.

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Display copies of books and complimentary copies of journals are available on site. SIAM books are available at a discounted price during the conference. If a SIAM books representative is not available, completed order forms and payment (credit cards are preferred) may be taken to the SIAM registration desk. The books table will close at 2:00 PM on Tuesday, January 12.

Registration Fee Includes
• Admission to all technical sessions
• ANALCO/ALENEX Business Meeting
• Coffee breaks daily
• Continental Breakfast daily
• Luncheon on Sunday, January 10, 2016
• Proceedings (SODA USB distributed onsite; ALENEX, ANALCO and SODA posted online)
• Room set-ups and audio/visual equipment
• SODA Business Meeting
• Welcome Reception

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A space for emergency contact information is provided on the back of your name badge. Help us help you in the event of an emergency!

Get-togethers
Welcome Reception
Saturday, January 9
6:00 PM - 8:00 PM
Arlington Ballroom Salon III - 1st Floor

ALENEX/ANALCO Business Meeting
Sunday, January 10
6:45 PM – 7:45 PM
Rosslyn - Lobby Level

SODA Business Meeting and Awards Presentation
Monday, January 11
6:45 PM – 7:45 PM
Arlington Ballroom Salon IV - 1st Floor
Complimentary beer and wine will be served.

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Best Paper and Best Student Paper Awards

**Best Student Paper**
An Improved Distributed Algorithm for Maximal Independent Set
Mohsen Ghaffari

This paper will be presented in CP4: Session 2A on Sunday, January 10.
See page 14 for session details.

**Best Papers**
Awarded jointly to two papers

An Improved Distributed Algorithm for Maximal Independent Set
Mohsen Ghaffari
This paper will be presented in CP4: Session 2A on Sunday, January 10.
See page 14 for session details.

Near-Optimal Light Spanners
Shiri Chechik and Christian Wulff-Nilsen
This paper will be presented in CP13: Session 5A, Monday, January 11.
See page 19 for session details.
SIAM Activity Group on Discrete Mathematics (SIAG/DM)
www.siam.org/activity/dm

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- Dénes König Prize
- DM-Net, an electronic forum
- Website

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- Participation in the selection of SIAG/DM officers

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- Vice-Chair: Lenore Cowen, Tufts University
- Program Director: Guantao Chen, Georgia State University
- Secretary: Rick Brewster, Thompson Rivers University

TO JOIN:
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Invited Plenary Speakers

** All Invited Plenary Presentations will take place in Arlington Ballroom- Salon IV**

Sunday, January 10
11:30 AM - 12:30 PM

**IP1** Distributed Algorithms in Ant Colonies

**Deborah M. Gordon**, Stanford University, USA

Monday, January 11
11:30 AM - 12:30 PM

**IP2** Beyond Planarity: On Geometric Intersection Graphs

**Sariel Har-Peled**, University of Illinois at Urbana-Champaign, USA

Tuesday, January 12
11:30 AM - 12:30 PM

**IP3** Combinatorial Statistics

**Elchanan Mossel**, University of Pennsylvania, USA
SIAM Presents

An audio-visual archive, comprised of more than 1800 presentations posted in 28 searchable topics, including:

- algebraic geometry
- atmospheric and oceanographic science
- computational science
- data mining
- geophysical science
- optimization
- uncertainty quantification
and more...

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In addition, you can view brief video clips of speaker interviews and topic overviews from sessions at Annual Meetings starting in 2010, as well as the 2013 SIAM Conference on Computational Science and Engineering and the 2014 SIAM Conference on the Life Sciences.

Plans for adding more content from SIAM meetings abound, including presentations from six meetings in 2015.

New presentations are posted every few months as the program expands with sessions from additional SIAM meetings. Users can search for presentations by category, speaker name, and/or keywords.

The audio, slide, and video presentations are part of SIAM’s outreach activities to increase the public’s awareness of mathematics and computational science in the real world, and to bring attention to exciting and valuable work being done in the field. Funding from SIAM, the National Science Foundation, and the Department of Energy was used to support this project.

www.siam.org/meetings/presents.php
SODA16, ALENEX16 and ANALCO16 At-a-Glance

Saturday, January 9

5:00 PM - 8:00 PM
Registration
Arlington Registration - 1st Floor

6:00 PM - 8:00 PM
Welcome Reception
Arlington Ballroom Salon III - 1st Floor

Luncheon
12:30 PM - 2:00 PM
Arlington Ballroom Salon IV - 1st Floor

University, USA
Deborah M. Gordon, Stanford
Colonies
IP1 Distributed Algorithms in Ant
11:30 AM - 12:30 PM
Arlington Ballroom Foyer - 1st Floor

Coffee Break
11:05 AM - 11:30 AM
Arlington Ballroom Foyer - 1st Floor

Registration
8:00 AM - 5:00 PM
Arlington Registration - 1st Floor

Continental Breakfast
8:30 AM
Arlington Ballroom Foyer - 1st Floor

Intermission
11:05 AM - 11:50 AM
Arlington Ballroom Foyer - 1st Floor

Concurrent Sessions
ALENEX: Session 1
Rosslyn - Lobby Level
CP1 Session 1A
Arlington Ballroom Salon IV - 1st Floor
CP2 Session 1B
Arlington Ballroom Salon V - 1st Floor
CP3 Session 1C
Arlington Ballroom Salon VI - 1st Floor

Session 1A
CP9 Session 3C
Arlington Ballroom Salon VI - 1st Floor

Session 1B
CP7 Session 3A
Arlington Ballroom Salon IV - 1st Floor
CP8 Session 3B
Arlington Ballroom Salon V - 1st Floor

Session 1C
CP1 Session 3C
Arlington Ballroom Salon VI - 1st Floor

Coffee Break
4:05 PM - 4:30 PM
Arlington Ballroom Foyer - 1st Floor

3:00 PM - 4:05 PM
Concurrent Sessions
ALENEX: Session 2
Rosslyn - Lobby Level
CP12 Session 5A
Arlington Ballroom Salon IV - 1st Floor
CP13 Session 5B
Arlington Ballroom Salon V - 1st Floor

Session 2A
CP18 Session 5C
Arlington Ballroom Salon VI - 1st Floor

Session 2B
CP16 Session 4A
Arlington Ballroom Salon IV - 1st Floor
CP17 Session 4B
Arlington Ballroom Salon V - 1st Floor

Session 2C
CP14 Session 4C
Arlington Ballroom Salon VI - 1st Floor

Session 2D
CP15 Session 4D
Arlington Ballroom Salon I - 1st Floor

Monday, January 10

8:00 AM - 5:00 PM
Registration
Arlington Registration - 1st Floor

Continental Breakfast
8:30 AM
Arlington Ballroom Foyer - 1st Floor

Intermission
11:05 AM - 11:50 AM
Arlington Ballroom Foyer - 1st Floor

Concurrent Sessions
ALENEX: Session 1
Rosslyn - Lobby Level
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Arlington Ballroom Salon VI - 1st Floor

Session 1A
CP9 Session 3C
Arlington Ballroom Salon VI - 1st Floor

Session 1B
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Coffee Break
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Session 2C
CP14 Session 4C
Arlington Ballroom Salon VI - 1st Floor

Session 2D
CP15 Session 4D
Arlington Ballroom Salon I - 1st Floor

Tuesday, January 12

8:00 AM - 5:00 PM
Registration
Arlington Registration - 1st Floor

Continental Breakfast
8:30 AM
Arlington Ballroom Foyer - 1st Floor

Intermission
11:05 AM - 11:50 AM
Arlington Ballroom Foyer - 1st Floor

Concurrent Sessions
ALENEX: Session 1
Rosslyn - Lobby Level
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CP8 Session 3B
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Session 1C
CP1 Session 3C
Arlington Ballroom Salon VI - 1st Floor

Coffee Break
4:05 PM - 4:30 PM
Arlington Ballroom Foyer - 1st Floor

3:00 PM - 4:05 PM
Concurrent Sessions
ALENEX: Session 2
Rosslyn - Lobby Level
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Session 2B
CP16 Session 4A
Arlington Ballroom Salon IV - 1st Floor
CP17 Session 4B
Arlington Ballroom Salon V - 1st Floor

Session 2C
CP14 Session 4C
Arlington Ballroom Salon VI - 1st Floor

Session 2D
CP15 Session 4D
Arlington Ballroom Salon I - 1st Floor

Key to abbreviations and symbols

= Award Presentation
= Business Meeting
= Coffee Break
= Continental Breakfast and Luncheon
= Refreshments
SODA, ALENEX and ANALCO Program

ACM-SIAM Symposium on Discrete Algorithms
January 10-12, 2016
Crystal Gateway Marriott
Arlington, Virginia, USA

ANALCO16
Meeting on Analytic Algorithmics and Combinatorics
January 11, 2016
Crystal Gateway Marriott
Arlington, Virginia, USA

ALENEX16
Meeting on Algorithm Engineering & Experiments
January 10, 2016
Crystal Gateway Marriott
Arlington, Virginia, USA
Saturday, January 9

Registration
5:00 PM-8:00 PM
Room:Arlington Registration - 1st Floor

Welcome Reception
6:00 PM-8:00 PM
Room:Arlington Ballroom Salon III - 1st Floor

Sunday, January 10

Registration
8:00 AM-5:00 PM
Room:Arlington Registration - 1st Floor

Continental Breakfast
8:30 AM
Room:Arlington Ballroom Foyer - 1st Floor

Sunday, January 10

ALENEX: Session 1
9:00 AM-11:05 AM
Room:Rosslyn - Lobby Level
Chair: Darren Strash, Karlsruhe Institute of Technology, Germany

9:00-9:20 Engineering Oracles for Time-Dependent Road Networks
Spyros Kontogiannis, Computer Technology Institute and Press, Greece; George Michalopoulos, University of Patras, Greece; Georgia Papastavrou, University of Ioannina, Greece; Andreas Paraskevopoulos, University of Patras, Greece; Dorothea Wagner, Karlsruhe Institute of Technology, Germany; Christos Zaroliagis, University of Patras, Greece

9:25-9:45 Scalable Transfer Patterns
Hannah Bast, Matthias Hertel, and Sabine Storandt, Albert-Ludwigs Universität Freiburg, Germany

9:50-10:10 Experimental Evaluation of Distributed Node Coloring Algorithms for Wireless Networks
Fabian Fuchs, Karlsruhe Institute of Technology, Germany

10:15-10:35 Generating Massive Scale Free Networks under Resource Constraints
Manuel Penschuck and Ulrich Meyer, Goethe University Frankfurt am Main, Germany

10:40-11:00 k-Way Hypergraph Partitioning Via n-Level Recursive Bisection
Sebastian Schlag, Vitali Henne, Tobias Heuer, Henning Meyerhenke, Peter Sanders, and Christian Schulz, Karlsruhe Institute of Technology, Germany
Sunday, January 10

CP1
Session 1A
9:00 AM-11:05 AM
Room: Arlington Ballroom Salon IV - 1st Floor
Chair: Qin Zhang, Indiana University Bloomington, USA
9:00-9:20 Locality-Sensitive Hashing Without False Negatives
Rasmus Pagh, IT University of Copenhagen, Denmark
9:25-9:45 New Directions in Nearest Neighbor Searching with Applications to Lattice Sieving
Anja Becker, École Polytechnique Fédérale de Lausanne, Switzerland; Leo Ducas, Centrum voor Wiskunde en Informatica (CWI), Netherlands; Nicolas Gama, Université de Versailles Saint-Quentin-en-Yvelines; Thijs Laarhoven, Technische Universiteit Eindhoven, The Netherlands
9:50-10:10 Undirected Graph Exploration with \(\Theta(\log \log N)\) Pebbles
Yann Disser, Jan Hackfeld, and Max Klimm, Technische Universität Berlin, Germany
10:15-10:35 Phase Transitions in Group Testing
Volkan Cevher and Jonathan Scarlett, École Polytechnique Fédérale de Lausanne, Switzerland
10:40-11:00 Tight Bounds for the Distribution-Free Testing of Monotone Conjunctions
Xi Chen and Jinyu Xie, Columbia University, USA

CP2
Session 1B
9:00 AM-11:05 AM
Room: Arlington Ballroom Salon V - 1st Floor
Chair: Xiaotie Deng, Shanghai Jiao Tong University, China
9:00-9:20 Designing Networks with Good Equilibria under Uncertainty
George Christodoulou and Alkmini Sgouritsa, University of Liverpool, United Kingdom
9:25-9:45 An Improved Combinatorial Polynomial Algorithm for the Linear Arrow-Debreu Market
Ran Duan, Tsinghua University, P. R. China; Jugal Garg and Kurt Mehlhorn, Max Planck Institute for Informatics, Germany
9:50-10:10 Characterisation of Strongly Stable Matchings
Pratik Ghosal, Adam Kunysz, and Katarzyna Paluch, University of Wroclaw, Poland
10:15-10:35 Learning and Efficiency in Games with Dynamic Population
Thodoris Lykouris, Cornell University, USA; Vasilis Syrgkanis, Microsoft Research, USA; Eva Tardos, Cornell University, USA
10:40-11:00 The Complexity of All-Switches Strategy Improvement
John Fearnley and Rahul Savani, University of Liverpool, United Kingdom

CP3
Session 1C
9:00 AM-11:05 AM
Room: Arlington Ballroom Salon VI - 1st Floor
Chair: Amin Saberi, Stanford University, USA
9:00-9:20 Make-to-Order Integrated Scheduling and Distribution
Yossi Azar, Tel Aviv University, Israel; Amir Epstein, IBM Haifa Research Lab, Israel; Lukasz Jez, Tel Aviv University, Israel and University of Wroclaw, Poland; Adi Vardi, Tel Aviv University, Israel
9:25-9:45 An \(O(\log m)\)-Competitive Algorithm for Online Machine Minimization
Lin Chen and Nicole Megow, Technische Universität München, Germany; Kevin Schewior, Technische Universität Berlin, Germany
9:50-10:10 Online Degree-Bounded Steiner Network Design
Sina Dehghani, Soheil Ehsani, and MohammadTaghi Hajiaghayi, University of Maryland, College Park, USA; Vahid Liasghat, Stanford University, USA; Saeed Seddighin, University of Maryland, College Park, USA
10:15-10:35 Scheduling Parallel Dag Jobs Online to Minimize Average Flow Time
Kunal Agrawal, Jing Li, Kefu Lu, and Benjamin Moseley, Washington University in St. Louis, USA
10:40-11:00 Online Pricing with Impatient Bidders
Marek Cygan, Marcin Mucha, Piotr Sankowski, and Qiang Zhang, University of Warsaw, Poland

Coffee Break
11:05 AM-11:30 AM
Room: Arlington Ballroom Foyer - 1st Floor
Like many distributed systems, both natural and engineered, ant colonies operate without any central control. No ant can assess what needs to be done. Each ant responds to its interactions with other ants nearby and in the aggregate, these stochastic, dynamical networks of interaction regulate colony behavior. Ant species use distributed algorithms that differ according to environmental constraints such as operating costs, time pressure, patchy distribution of resources in space, risk of rupture, and the intensity of competition. I will discuss algorithms used to regulate foraging activity, nestmate recognition and colony security, collective search, and stable trail systems.

Deborah M. Gordon
Stanford University, USA
Sunday, January 10

CP5
Session 2B
2:00 PM-3:40 PM
Room: Arlington Ballroom Salon V - 1st Floor
Chair: Eric Price, University of Texas at Austin, USA
2:00-2:20 Nearly Tight Oblivious Subspace Embeddings by Trace Inequalities
Michael B. Cohen, Massachusetts Institute of Technology, USA
2:25-2:45 The Restricted Isometry Property of Subsampled Fourier Matrices
Ishay Haviv, Academic College of Tel-Aviv Yaffo, Israel; Oded Regev, New York University, USA
2:50-3:10 Nearly Optimal Deterministic Algorithm for Sparse Walsh-Hadamard Transform
Mahdi Cheraghchi, Imperial College London, United Kingdom; Piotr Indyk, Massachusetts Institute of Technology, USA
3:15-3:35 Nearly-Optimal Bounds for Sparse Recovery in Generic Norms, with Applications to k-Median Sketching
Arturs Backurs, Piotr Indyk, and Ilya Razenshteyn, Massachusetts Institute of Technology, USA; David Woodruff, IBM Almaden Research Center, USA

Sunday, January 10

CP6
Session 2C
2:00 PM-3:40 PM
Room: Arlington Ballroom Salon VI - 1st Floor
Chair: Barna Saha, University of Massachusetts, Amherst, USA
2:00-2:20 Simpler, Faster and Shorter Labels for Distances in Graphs
Stephen Alstrup, University of Copenhagen, Denmark; Cyril Gavoille, Bordeaux University, France; Esben Bistrup Halvorsen, University of Copenhagen, Denmark; Holger Petersen, Unaffiliated
2:25-2:45 Approximate Distance Oracles for Planar Graphs with Improved Query Time-Space Tradeoff
Christian Wulff-Nilsen, University of Copenhagen, Denmark
2:50-3:10 New Bounds for Approximating Extremal Distances in Undirected Graphs
Massimo Cairo, Scuola Normale Superiore, Pisa, Italy; Roberto Grossi, Universita di Pisa, Italy; Romeo Rizzi, Universita di Verona, Italy
3:15-3:35 Approximation and Fixed Parameter Subquadratic Algorithms for Radius and Diameter in Sparse Graphs
Amir Abboud, Virginia Vassilevska Williams, and Joshua R. Wang, Stanford University, USA

Coffee Break
4:05 PM-4:30 PM
Room: Arlington Ballroom Foyer - 1st Floor

Sunday, January 10

ALENEX: Session 3
4:30 PM-6:35 PM
Room: Rosslyn - Lobby Level
Chair: Ulrich Meyer, Goethe University Frankfurt am Main, Germany
4:30-4:50 Real-Time k-Bounded Preemptive Scheduling
Baruch Schieber, IBM T.J. Watson Research Center, USA; Sivan Albagli-Kim and Hadas Shachnai, Technion, Israel; Tami Tamir, The Interdisciplinary Center, Israel
4:55-5:15 Finding Near-Optimal Independent Sets at Scale
Sebastian Lamm, Peter Sanders, Christian Schulz, and Darren Strash, Karlsruhe Institute of Technology, Germany; Renato F. Werneck, Unaffiliated
5:20-5:40 A Novel Dual Ascent Algorithm for Solving the Min-Cost Flow Problem
Ruben Becker, Maximilian Fickert, and Andreas Karrenbauer, Max Planck Institute for Informatics, Germany
5:45-6:05 A General Framework for Dynamic Succinct and Compressed Data Structures
Patrick K. Nicholson, Bell Labs, Ireland; Patrick Klitzke, Universität des Saarlandes, Germany
6:10-6:30 An Empirical Comparison of Graph Laplacian Solvers
Kevin Deweese and John R. Gilbert, University of California, Santa Barbara, USA; Erik G. Boman, Sandia National Laboratories, USA
Sunday, January 10

CP7
Session 3A
4:30 PM-6:35 PM
Room: Arlington Ballroom Salon IV - 1st Floor
Chair: Eden Chlamtác, Ben Gurion University Negev, Israel
4:30-4:50 Deterministic Algorithms for Submodular Maximization Problems
Niv Buchbinder, Tel Aviv University, Israel; Moran Feldman, École Polytechnique Fédérale de Lausanne, Switzerland
4:55-5:15 Improved Approximation Algorithms for k-Submodular Function Maximization
Satoru Iwata, University of Tokyo, Japan; Shin-ichi Tanigawa, Kyoto University, Japan; Yuichi Yoshida, National Institute of Informatics, Japan
5:20-5:40 Locally Adaptive Optimization: Adaptive Seeding for Monotone Submodular Functions
Ashwinkumar Badanidiyuru, Google, Inc., USA; Christos Papadimitriou and Avidi Rubinstein, University of California, Berkeley, USA; Yaron Singer and Lior Seeman, Harvard University, USA
5:45-6:05 Exact and Approximation Algorithms for Weighted Matroid Intersection
Chien-Chung Huang, Chalmers University of Technology, Sweden; Naonori Kakimura, University of Tokyo, Japan; Naoyuki Kamiyama, Kyushu University, Japan
6:10-6:30 A Fast Approximation for Maximum Weight Matroid Intersection
Chandra Chekuri and Kent Quanrud, University of Illinois at Urbana-Champaign, USA

CP8
Session 3B
4:30 PM-6:35 PM
Room: Arlington Ballroom Salon V - 1st Floor
Chair: Mordecai Golin, Hong Kong University of Science and Technology, Hong Kong
4:30-4:50 The Complexity of Approximately Counting in 2-Spin Systems on k-Uniform Bounded-Degree Hypergraphs
Andreas Galanis and Leslie Ann Goldberg, University of Oxford, United Kingdom
4:55-5:15 Bounds for Random Constraint Satisfaction Problems Via Spatial Coupling
Dimitris Achlioptas, University of California, Santa Cruz, USA; Hamed Hassani, Nicolas Macris, and Rudiger Urbanke, École Polytechnique Fédérale de Lausanne, Switzerland
5:20-5:40 Evolutionary Dynamics in Finite Populations Mix Rapidly
Ioannis Panageas, Georgia Institute of Technology, USA; Piyush Srivastava, California Institute of Technology, USA; Nisheeth K. Vishnoi, École Polytechnique Fédérale de Lausanne, Switzerland
5:45-6:05 Random-Cluster Model in Z2
Antonio Blanca and Alistair Sinclair, University of California, Berkeley, USA
6:10-6:30 Canonical Paths for Mcmc: from Art to Science
Lingxiao Huang, Tsinghua University, P. R. China; Pinyan Lu, Microsoft Research Asia; Chihao Zhang, Shanghai Jiao Tong University, China

CP9
Session 3C
4:30 PM-6:35 PM
Room: Arlington Ballroom Salon VI - 1st Floor
Chair: Aravindan Vijayaraghavan, New York University, USA
4:30-4:50 Towards Optimal Algorithms for Prediction with Expert Advice
Nick Gravin and Yuval Peres, Microsoft Research, USA; Balasubramanian Sivan, Google Research, USA
4:55-5:15 Sparse Approximation Via Generating Point Sets
Avrim Blum, Carnegie Mellon University, USA; Sariel Har-Peled, University of Illinois at Urbana-Champaign, USA; Benjamin A. Recht, University of Texas at Dallas, USA
5:20-5:40 Weighted SGD for \ell_p Regression with Randomized Preconditioning
Yinlam Chow, Stanford University, USA; Michael Mahoney, University of California, Berkeley, USA; Christopher Re and Jiyan Yang, Stanford University, USA
5:45-6:05 Non-Convex Compressed Sensing with the Sum-of-Squares Method
Tasuku Soma, University of Tokyo, Japan; Yuichi Yoshida, National Institute of Informatics, Japan
6:10-6:30 Jointly Private Convex Programming
Justin Hsu, University of Pennsylvania, USA; Zhiyi Huang, University of Hong Kong, Hong Kong; Aaron Roth and Steven Wu, University of Pennsylvania, USA

Intermission
6:35 PM-6:45 PM

ALENEX and ANALCO Business Meeting
6:45 PM-7:45 PM
Room: Rosslyn - Lobby Level
Monday, January 11

CP10
Session 4A
9:00 AM-11:05 AM
Room: Arlington Ballroom Salon IV- 1st Floor
Chair: David Kempe, University of Southern California, USA

9:00-9:20 Time Vs. Information Tradeoffs For Leader Election in Anonymous Trees
Christian Glacet, CNR, Italy; Avery Miller, Tel Aviv University, Israel; Andrzej Pelc, Université du Québec en Outaouais, Canada

9:25-9:45 Markovian Hitlers and the Complexity of Blind Rendezvous
Sixia Chen, Central Connecticut State University, USA; Matthew Dippel, Northeastern University, USA; Alexander Russell, University of Connecticut, USA; Abhishek Samanta, Google, Inc., USA; Ravi Sundaram, Northeastern University, USA

9:50-10:10 Stabilizing Consensus with Many Opinions
Luca Becchetti, University of Rome La Sapienza, Italy; Andrea Clementi, University of Rome II, Tor Vergata, Italy; Emanuele Natale, University of Rome La Sapienza, Italy; Francesco Pasquale, University of Rome II, Tor Vergata, Italy; Luca Trevisan, University of California, Berkeley, USA

10:15-10:35 How to Scale Exponential Backoff: Constant Throughput, Polylog Access Attempts, and Robustness
Michael A. Bender, Stony Brook University, USA; Jeremy T. Fineman, Georgetown University, USA; Seth Gilbert, National University of Singapore, Singapore; Maxwell Young, Mississippi State University, USA

10:40-11:00 Balanced Allocation: Patience is not a Virtue
John Augustine and William K. Moses Jr., Indian Institute of Technology Madras, India; Amanda Redlich, Bowdoin College, USA; Eli Upfal, Brown University, USA

ANALCO: Session 1
9:00 AM-11:05 AM
Room: Rosslyn - Lobby Level
Chair: James Fill, Johns Hopkins University, USA

9:00-9:20 More Analysis of Double Hashing for Balanced Allocations
Michael Mitzenmacher, Harvard University, USA

9:25-9:45 A New Approach to Analyzing Robin Hood Hashing
Michael Mitzenmacher, Harvard University, USA

9:50-10:10 Randomized Strategies for Cardinality Robustness in the Knapsack Problem
Yusuke Kobayashi, University of Tsukuba, Japan; Kenjiro Takazawa, Kyoto University, Japan

10:15-10:35 Graphs with Degree Constraints
Elie De Panafieu, Université Pierre et Marie Curie - Paris VI, France; Lander Ramos, Universitat Politecnica de Catalunya, Spain

10:40-11:00 Sampling Weighted Perfect Matchings on the Square-Octagon Lattice
Prateek Bhakta and Dana Randall, Georgia Institute of Technology, USA

Monday, January 11

CP11
Session 4B
9:00 AM-11:05 AM
Room: Arlington Ballroom Salon V- 1st Floor
Chair: Danupon Nanongkai, KTH Royal Institute of Technology, Sweden

9:00-9:20 Weighted Dynamic Finger in Binary Search Trees
John Iacono, New York University, USA; Stefan Langerman, Université Libre de Bruxelles, Belgium

9:25-9:45 Faster Fully Dynamic Matchings with Small Approximation Ratios
Aaron Bernstein, Columbia University, USA

9:50-10:10 Dynamic (1+\(\varepsilon\))-Approximate Matchings: A Density-Sensitive Approach
David Peleg, Weizmann Institute of Science, Israel; Shay Solomon, Tel Aviv University, Israel

10:15-10:35 Dynamic Dfs in Undirected Graphs: Breaking the \(O(m)\) Barrier
Surender Baswana, Shreijit Chaudhury, Keerti Choudhary, and Shahbaz Khan, Indian Institute of Technology, Kanpur, India

10:40-11:00 On Dynamic Approximate Shortest Paths for Planar Graphs with Worst-Case Costs
Ittai Abraham, Vmware Research, USA; Shiri Chechik, Tel Aviv University, Israel; Daniel Delling, Andrew Goldberg, and Renato F. Werneck, Unaffiliated
Monday, January 11

**CP12**

**Session 4C**

9:00 AM-11:05 AM  
*Room: Arlington Ballroom Salon VI - 1st Floor*

**Chair:** Eden Chlamtác, Ben Gurion University Negev, Israel

9:00-9:20 Reducing Curse of Dimensionality: Improved PTAS for TSP (with Neighborhoods) in Doubling Metrics  
T-H. Hubert Chan and Shaofeng H.-C. Jiang, University of Hong Kong, Hong Kong

9:25-9:45 Clustering Time Series under the Frechet Distance  
Anne Driemel, Technische Universiteit Eindhoven, The Netherlands; Amer Krivosija and Christian Sohler, Technische Universität Dortmund, Germany

9:50-10:10 Approximating Capacitated k-Median with (1+ε)k Open Facilities  
Shi Li, Toyota Technological Institute at Chicago, USA

10:15-10:35 Simple and Fast Rounding Algorithms for Directed and Node-Weighted Multiway Cut  
Chandra Chekuri and Vivek Madan, University of Illinois at Urbana-Champaign, USA

10:40-11:00 Constant Factor Approximation for Subset Feedback Set Problems Via a New Lp Relaxation  
Chandra Chekuri and Vivek Madan, University of Illinois at Urbana-Champaign, USA

**Coffee Break**  
11:05 AM-11:30 AM  
*Room: Arlington Ballroom Foyer - 1st Floor*

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Monday, January 11

**IP2**

**Beyond Planarity: On Geometric Intersection Graphs**

11:30 AM-12:30 PM  
*Room: Arlington Ballroom Salon IV - 1st Floor*

**Chair:** Robert Krauthgamer, Weizmann Institute of Science, Israel

Many efficient algorithms had been developed over the years for planar graphs and more general graphs such as low genus graphs. Intersection graphs of geometric objects (in low dimensions) with some additional properties, such as fatness or low density, provide yet another family of graphs for which one can design better algorithms. This family is a vast extension of planar graphs, and yet is still algorithmically tractable for many problems. In this talk, we will survey this class of graphs, and some algorithms and intractability results known for such graphs, and outline open problems for further research.

**Sariel Har-Peled**  
University of Illinois at Urbana-Champaign, USA

**Lunch Break**

12:30 PM-2:00 PM  
*Attendees on their own*

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Monday, January 11

**ANALCO:**

**Session 2**

2:00 PM-4:05 PM  
*Room: Rosslyn - Lobby Level*

**Chair:** Mark Daniel Ward, Purdue University, USA

2:00-2:20 Scaling Limit of Random k-Trees  
Michael Drmota, Technische Universitat Vienna, Austria; Emma Yu Jin, Vienna University of Technology, Austria

2:25-2:45 Betweenness Centrality in Random Trees  
Kevin Durant and Stephan Wagner, Stellenbosch University, South Africa

2:50-3:10 On Terminal Nodes and the Degree Profile of Preferential Dynamic Attachment Circuits  
Panpan Zhang, George Washington University, USA

3:15-3:35 Prime Factorization of the Kirchhoff Polynomial: Compact Enumeration of Arborescences  
Przemyslaw Uznanski, Aalto University, Finland; Matus Mihalak, Maastricht University, The Netherlands; Pencho Yordanov, ETH Zürich, Switzerland

3:40-4:00 Ricci-Ollivier Curvature of the Rooted Phylogenetic Subtree-Prune-Regraft Graph  
Chris Whidden and Frederick Matsen, Fred Hutchinson Cancer Research Center, USA
Monday, January 11

**CP13**

**Session 5A**

2:00 PM-4:05 PM

Room: Arlington Ballroom Salon IV - 1st Floor

Chair: Fabrizio Grandoni, University of Lugano, Switzerland

2:00-2:20 *Approximating Low-Stretch Spanners*

Michael Dinitz and Zeyu Zhang, Johns Hopkins University, USA

2:25-2:45 *Error Amplification for Pairwise Spanner Lower Bounds*

Amir Abboud and Greg Bodwin, Stanford University, USA

2:50-3:10 *Better Distance Preservers and Additive Spanners*

Greg Bodwin and Virginia Williams, Stanford University, USA

3:15-3:35 *On Notions of Distortion and An Almost Minimum Spanning Tree with Constant Average Distortion*

Yair Bartal, Hebrew University, Israel; Arnold Filtser, Ben Gurion University, Israel; Ofer Neiman, Ben-Gurion University, Israel

3:40-4:00 *Near-Optimal Light Spanners*

Shiri Chechik, Tel Aviv University, Israel; Christian Wulff-Nilsen, University of Copenhagen, Denmark

**CP14**

**Session 5B**

2:00 PM-3:40 PM

Room: Arlington Ballroom Salon V - 1st Floor

Chair: Robert Krauthgamer, Weizmann Institute of Science, Israel

2:00-2:20 *Efficient Quantum Algorithms for Computing Class Groups and Solving the Principal Ideal Problem in Arbitrary Degree Number Fields*

Jean-François Biasse, University of South Florida, USA; Fang Song, University of Waterloo, Canada

2:25-2:45 *Efficient Quantum Algorithms for (Gapped) Group Testing and Junta Testing*

Andris Ambainis, University of Latvia, Latvia; Aleksandrs Belovs, Centrum voor Wiskunde en Informatica (CWI), Netherlands; Ronald de Wolf, University of Amsterdam, Netherlands; Oded Regev, New York University, USA

2:50-3:10 *Permutation Patterns Are Hard to Count*

Igor Pak and Scott Garrabrant, University of California, Los Angeles, USA

3:15-3:35 *Computing in Continuous Space with Self-Assembling Polygonal Tiles*

Jacob Hendricks, University of Wisconsin-River Falls, USA; Gilbert Oscar, Matthew Patitz, and Trent Rogers, University of Arkansas, USA

**CP15**

**Session 5C**

2:00 PM-3:40 PM

Room: Arlington Ballroom Salon VI - 1st Floor

Chair: Vida Dujmovic, University of Ottawa, Canada

2:00-2:20 *Beyond the Richter-Thomassen Conjecture*

János Pach, EPFL, Switzerland; Natan Rubin, Ben-Gurion University, Israel; Gábor Tardos, Renyi Institute, Hungary

2:25-2:45 *Treetopes and Their Graphs*

David Eppstein, University of California, Irvine, USA

2:50-3:10 *Windrose Planarity: Embedding Graphs with Direction-Constrained Edges*

Patrizio Angelini, Universitàt Tübingen, Germany; Giordano Da Lozzo, Università degli Studi Roma Tre, Italy; Giuseppe Di Battista, Universita di Roma Tre, Italy; Valentino Di Donato, Università degli Studi Roma Tre, Italy; Philipp Kindermann, Fern Universität in Hagen, Germany; Günter Rote, Freie Universität Berlin, Germany; Ignaz Rutter, Karlsruhe Institute of Technology, Germany

3:15-3:35 *Multiscale Mapper: Topological Summarization Via Codomain Covers*

Tamal Dey, Facundo Memoli, and Yusu Wang, Ohio State University, USA

**Coffee Break**

4:05 PM-4:30 PM

Room: Arlington Ballroom Foyer - 1st Floor
### Monday, January 11

#### ANALCO: Session 3

**4:30 PM-5:45 PM**

**Room:** Rosslyn - Lobby Level  
**Chair:** Hosam M. Mahmoud, George Washington University, USA  
**4:30-4:50** The Clt Analogue for Cyclic Urns  
**Ralph Neininger,** J.W. Goethe-Universität, Germany  
**4:55-5:15** On Connectivity in a General Random Intersection Graph  
**Jun Zhao,** Carnegie Mellon University, USA; **Panpan Zhang,** George Washington University, USA  
**5:20-5:40** Total Variation Discrepancy of Deterministic Random Walks for Ergodic Markov Chains  
**Takeharu Shiraga,** Yukiko Yamauchi, Shuji Kijima, and Masafumi Yamashita, Kyushu University, Japan

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#### Monday, January 11

#### CP16

**Session 6A**  
**4:30 PM-6:35 PM**  
**Room:** Arlington Ballroom Salon IV- 1st Floor  
**Chair:** Mohit Singh, Microsoft Research, USA  
**4:30-4:50** Online Contention Resolution Schemes  
**Moran Feldman,** The Open University of Israel, Israel; **Ola Svensson,** École Polytechnique Fédérale de Lausanne, Switzerland; **Rico Zenklusen,** ETH Zürich, Switzerland  
**4:55-5:15** Finding a Stable Allocation in Polymatroid Intersection  
**Satoru Iwata** and **Yu Yokoi,** University of Tokyo, Japan  
**5:20-5:40** An Improved Approximation Guarantee for the Maximum Budgeted Allocation Problem  
**Christos Kalaitzis,** École Polytechnique Fédérale de Lausanne, Switzerland  
**5:45-6:05** The Matching Problem Has No Small Symmetric Sdp  
**Gabor Braun,** Georgia Institute of Technology, USA; **Jonah Brown-Cohen,** University of California, Berkeley, USA; **Arefin Huq** and **Sebastian Pokutta,** Georgia Institute of Technology, USA; **Prasad Raghavendra,** University of California, Berkeley, USA; **Aurko Roy,** Georgia Institute of Technology, USA; **Benjamin Weitz,** University of California, Berkeley, USA; **Daniel Zink,** Georgia Institute of Technology, USA  
**6:10-6:30** On the Integrality Gap of Degree-4 Sum of Squares for Planted Clique  
**Samuel Hopkins,** Cornell University, USA; **Pravesh Kothari,** University of Texas at Austin, USA; **Aaron Potechin,** Simons Institute, USA; **Prasad Raghavendra** and **Tselil Schramm,** University of California, Berkeley, USA

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#### Monday, January 11

#### CP17

**Session 6B**  
**4:30 PM-6:35 PM**  
**Room:** Arlington Ballroom Salon V- 1st Floor  
**Chair:** Philip Klein, Brown University, USA  
**4:30-4:50** Linear Recognition of Almost Interval Graphs  
**Yixin Cao,** Hong Kong Polytechnic University, China  
**4:55-5:15** Subexponential Parameterized Algorithm for Interval Completion  
**Ivan Bliznets,** Steklov Institute of Mathematics, Russia; **Fedor Fomin,** University of Bergen, Norway; **Marcin Pilipczuk** and **Michal Pilipczuk,** University of Warsaw, Poland  
**5:20-5:40** Lower Bounds for the Parameterized Complexity of Minimum Fill-In and Other Completion Problems  
**Ivan Bliznets,** Steklov Institute of Mathematics, Russia; **Marek Cygan** and **Pawel Komosa,** University of Warsaw, Poland; **Lukas Mach,** University of Warwick, United Kingdom; **Michal Pilipczuk,** University of Warsaw, Poland  
**5:45-6:05** Raising The Bar For Vertex Cover: Fixed-Parameter Tractability Above A Higher Guarantee  
**Shivam Garg,** Indian Institute of Technology Bombay, India; **Geevarghese Philip,** Chennai Mathematical Institute, India  
**6:10-6:30** Directed Multicut Is W(1)-Hard, Even for Four Terminal Pairs  
**Marcin Pilipczuk,** University of Warsaw, Poland; **Magnus Wahlström,** Royal Holloway, University of London, United Kingdom
Monday, January 11

CP18

Session 6C
4:30 PM-6:35 PM

Room: Arlington Ballroom Salon VI - 1st Floor

Chair: Mordecai Golin, Hong Kong University of Science and Technology, Hong Kong

4:30-4:50 An Efficient Algorithm for Computing High-Quality Paths Amid Polygonal Obstacles
Pankaj Agarwal and Kyle Fox, Duke University, USA; Oren Salzman, Tel Aviv University, Israel

4:55-5:15 Approximating the K-Level in Three-Dimensional Plane Arrangements
Sariel Har-Peled, University of Illinois at Urbana-Champaign, USA

5:20-5:40 Improved Deterministic Algorithms for Linear Programming in Low Dimensions
Timothy M. Chan, University of Waterloo, Canada

5:45-6:05 A Fast and Simple Algorithm for Computing Approximate Euclidean Minimum Spanning Trees
Sunil Arya, Hong Kong University of Science and Technology, Hong Kong; David M. Mount, University of Maryland, USA

6:10-6:30 Persistent Homology and Nested Dissection
Michael Kerber, Graz University of Technology, Austria; Donald Sheehy, University of Connecticut, USA; Primoz Skraba, Jozef Stefan Institute, Slovenia

Intermission
6:35 PM-6:45 PM

SODA Business Meeting and Awards Presentation
6:45 PM-7:45 PM

Room: Arlington Ballroom Salon IV - 1st Floor

Complimentary wine and beer will be served.

Tuesday, January 12

CP19

Session 7A
9:00 AM-11:05 AM

Room: Arlington Ballroom Salon IV - 1st Floor

Chair: Fabrizio Grandoni, University of Lugano, Switzerland

9:00-9:20 Deterministic APSP, Orthogonal Vectors, and More: Quickly Derandomizing Razborov-Smolensky
Timothy Chan, University of Waterloo, Canada; Ryan Williams, Stanford University, USA

9:25-9:45 Subtree Isomorphism Revisited
Amir Abboud, Stanford University, USA; Arturs Backurs, Massachusetts Institute of Technology, USA; Thomas Dueholm Hansen, Aarhus University, Denmark; Virginia Vassilevska Williams, Stanford University, USA; Or Zamir, Tel Aviv University, Israel

9:50-10:10 Higher Lower Bounds from the 3sum Conjecture
Tsvi Kopelowitz, University of Michigan, USA; Ely Porat, Bar-Ilan University, Israel; Seth Pettie, University of Michigan, Ann Arbor, USA

10:15-10:35 A Faster Subquadratic Algorithm for Finding Outlier Correlations
Matti Karppa, Petteri Kaski, and Jukka Kohonen, Aalto University, Finland

Tuesday, January 12

CP18

Session 6C
4:30 PM-6:35 PM

Room: Arlington Ballroom Salon VI - 1st Floor

Chair: Mordecai Golin, Hong Kong University of Science and Technology, Hong Kong

4:30-4:50 An Efficient Algorithm for Computing High-Quality Paths Amid Polygonal Obstacles
Pankaj Agarwal and Kyle Fox, Duke University, USA; Oren Salzman, Tel Aviv University, Israel

4:55-5:15 Approximating the K-Level in Three-Dimensional Plane Arrangements
Sariel Har-Peled, University of Illinois at Urbana-Champaign, USA

5:20-5:40 Improved Deterministic Algorithms for Linear Programming in Low Dimensions
Timothy M. Chan, University of Waterloo, Canada

5:45-6:05 A Fast and Simple Algorithm for Computing Approximate Euclidean Minimum Spanning Trees
Sunil Arya, Hong Kong University of Science and Technology, Hong Kong; David M. Mount, University of Maryland, USA

6:10-6:30 Persistent Homology and Nested Dissection
Michael Kerber, Graz University of Technology, Austria; Donald Sheehy, University of Connecticut, USA; Primoz Skraba, Jozef Stefan Institute, Slovenia

Intermission
6:35 PM-6:45 PM

SODA Business Meeting and Awards Presentation
6:45 PM-7:45 PM

Room: Arlington Ballroom Salon IV - 1st Floor

Complimentary wine and beer will be served.
Tuesday, January 12

**CP20**

**Session 7B**

9:00 AM-11:05 AM

*Room:* Arlington Ballroom Salon V - 1st Floor

*Chair:* Qin Zhang, Indiana University Bloomington, USA

9:00-9:20 Kernelization via Sampling with Applications to Finding Matchings and Related Problems in Dynamic Graph Streams

Rajesh Chitnis, Weizmann Institute of Science, Israel; Graham Cormode, University of Warwick, United Kingdom; Hossein Esfandiari and MohammadTaghi Hajiaghayi, University of Maryland, College Park, USA; Andrew McGregor, University of Massachusetts, Amherst, USA; Morteza Monemizadeh, Charles University, Czech Republic; Sofya Vorotnikova, University of Massachusetts, Amherst, USA

9:25-9:45 Maximum Matchings in Dynamic Graph Streams and the Simultaneous Communication Model

Sepehr Assadi, Sanjeev Khanna, Yang Li, and Grigory Yaroslavtsev, University of Pennsylvania, USA

9:50-10:10 Incidence Geometries and the Pass Complexity Of Semi-Streaming Set Cover

Amit Chakrabarti, Dartmouth College, USA; Anthony Wirth, University of Melbourne, Australia

10:15-10:35 Clustering Problems on Sliding Windows

Vladimir Braverman, Harry Lang, and Keith Levin, Johns Hopkins University, USA; Morteza Monemizadeh, Charles University, Czech Republic

10:40-11:00 Gowers Norm, Function Limits, and Parameter Estimation

Yuichi Yoshida, National Institute of Informatics, Japan

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Tuesday, January 12

**CP21**

**Session 7C**

9:00 AM-11:05 AM

*Room:* Arlington Ballroom Salon VI - 1st Floor

*Chair:* Nicole Immorlica, Microsoft Research, USA

9:00-9:20 On the Economic Efficiency of the Combinatorial Clock Auction

Nicolas Bouquet and Yang Cai, McGill University, Canada; Christoph Hunkenschröder, University of Bonn, Germany; Adrian Vetta, McGill University, Canada

9:25-9:45 Approximately Efficient Double Auctions with Strong Budget Balance

Riccardo Colini-Baldeschi, Bart De Keijzer, and Stefano Leonardi, University of Rome La Sapienza, Italy; Stefano Turchetta, University of Oxford, United Kingdom

9:50-10:10 Interpolating Between Truthful and Non-Truthful Mechanisms for Combinatorial Auctions

Mark Braverman, Jiemeing Mao, and Seth M. Weinberg, Princeton University, USA

10:15-10:35 On the Complexity of Dynamic Mechanism Design

Christos Papadimitriou, George Pierrakos, Alexandros Psomas, and Aviad Rubinstein, University of California, Berkeley, USA

10:40-11:00 Simple Pricing Schemes for Consumers With Evolving Values

Shuichi Chawla, University of Wisconsin, Madison, USA; Nikhil R. Devanur, Microsoft Research, USA; Anna R. Karlin, University of Washington, USA; Balasubramanian Sivan, Google Research, USA

Coffee Break

11:05 AM-11:30 AM

*Room:* Arlington Ballroom Foyer - 1st Floor

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Tuesday, January 12

**CP22**

**Session 7D**

9:00 AM-11:05 AM

*Room:* Rosslyn - Lobby Level

*Chair:* Mohit Singh, Microsoft Research, USA

9:00-9:20 On Approximating Strip Packing with a Better Ratio Than 3/2

Giorgi Nadiradze, Saarland University, Germany; Andreas Wieße, MPII Saarbrücken, Germany

9:25-9:45 Packing Small Vectors

Yossi Azar, Ilan R. Cohen, Amos Fiat, and Alan Roytman, Tel Aviv University, Israel

9:50-10:10 Approximation Schemes for Machine Scheduling with Resource (in-)dependent Processing Times

Klaus Jansen, Marten Maack, and Malin Rau, University of Kiel, Germany

10:15-10:35 Integrality Gaps and Approximation Algorithms for Dispersers and Bipartite Expanders

Xue Chen, University of Texas at Austin, USA

10:40-11:00 Improved Approximation for Vector Bin Packing

Nikhil Bansal and Marek Elias, Eindhoven University of Technology, Netherlands; Arindam Khan, Georgia Institute of Technology, USA

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Tuesday, January 12

Coffee Break

11:05 AM-11:30 AM

*Room:* Arlington Ballroom Foyer - 1st Floor
**Tuesday, January 12**

**IP3**

**Combinatorial Statistics**

11:30 AM - 12:30 PM

Room: Arlington Ballroom Salon IV - 1st Floor

Chair: David Gamarnik, Massachusetts Institute of Technology, USA

Combinatorial Statistics studies combinatorial algorithms for inference of statistical models that are parametrized by discrete parameters such as trees, graphs or permutations. The talk will outline some key examples of results and techniques that are used in this area as well as future research directions and open problems.

Elchanan Mossel

*University of Pennsylvania, USA*

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**Lunch Break**

12:30 PM - 2:00 PM

Attendees on their own

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**Tuesday, January 12**

**CP23**

**Session 8A**

2:00 PM - 4:05 PM

Room: Arlington Ballroom Salon IV - 1st Floor

Chair: Xiaotie Deng, Shanghai Jiao Tong University, China

2:00-2:20 Connectivity in Bridge-Addable Graph Classes: the McDiarmid-Steger-Welsh Conjecture

Guillaume Chapuy, CNRS, Université Paris Diderot, France and Université de Montréal, Canada; Guillem Perarnau, McGill University, Canada

2:25-2:45 Recovery and Rigidity in a Regular Stochastic Block Model

Gerandy Brito, Ioana Dumitriu, Shirshendu Ganguly, and Christopher Hoffman, University of Washington, USA; Linh V. Tran, Vietnam National University at Ho Chi Minh City, Vietnam

2:50-3:10 Distance in the Forest Fire Model How Far Are You from Eve?

Frederik Mallmann-Trenn, Ecole Normale Supérieure, France

3:15-3:35 Species Trees from Gene Trees Despite a High Rate of Lateral Genetic Transfer: A Tight Bound

Sebastien Roch, University of Wisconsin, Madison, USA

3:40-4:00 The Power of Two Choices with Simple Tabulation

Søren Dahlgaard, Mathias Knudsen, and Eva Rotenberg, University of Copenhagen, Denmark; Mikkel Thorup, AT&T Labs Research, USA and University of Copenhagen, Denmark

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**Tuesday, January 12**

**CP24**

**Session 8B**

2:00 PM - 4:05 PM

Room: Arlington Ballroom Salon V - 1st Floor

Chair: Vida Dujmovic, University of Ottawa, Canada

2:00-2:20 Tight Bounds for Graph Homomorphism and Subgraph Isomorphism

Marek Cygan, University of Warsaw, Poland; Fedor Fomin, University of Bergen, Norway; Aleksandr Golovnev, New York University, USA; Alexander Kulikov, Steklov Institute of Mathematics, Russia; Ivan Mihajlin, University of California, San Diego, USA; Jakub Pachocki, Carnegie Mellon University, USA; Arkadiusz Socała, University of Warsaw, Poland

2:25-2:45 Tight Conditional Lower Bounds for Counting Perfect Matchings on Graphs of Bounded Treewidth, Cliquewidth, and Genus

Radu Curticapean, University of California, Berkeley, USA; Dániel Marx, Hungarian Academy of Sciences, Hungary

2:50-3:10 Discovering Archipelagos of Tractability for Constraint Satisfaction and Counting

Robert Ganian, Ramanujan M. Sridharan, and Stefan Szeider, Technische Universität Wien, Austria

3:15-3:35 Blocking Optimal k-Arborescences

Attila Bernáth and Tamás Király, Eötvös University, Hungary

3:40-4:00 Constructive Algorithm for Path-Width of Matroids

Jisu Jeong, KAIST, Korea; Eun Jung Kim, CNRS, France; Sang-Il Oum, KAIST, Korea
Tuesday, January 12

CP25
Session 8C
2:00 PM-4:05 PM

Room: Arlington Ballroom Salon VI - 1st Floor
Chair: Barna Saha, University of Massachusetts, Amherst, USA

2:00-2:20 Approximation of Non-boolean 2csp
Guy Kindler, Hebrew University, Israel; Alexandra Kolla, University of Illinois at Urbana-Champaign, USA; Luca Trevisan, University of California, Berkeley, USA

2:25-2:45 An FPTAS for Minimizing Indefinite Quadratic Forms over Integers in Polyhedra
Robert Hildebrand, Robert Weismantel, and Kevin Zemmer, ETH Zürich, Switzerland

2:50-3:10 Nearly Optimal NP-Hardness of Unique Coverage
Venkatesan Guruswami and Euiwoong Lee, Carnegie Mellon University, USA

3:15-3:35 Algorithms and Adaptivity Gaps for Stochastic Probing
Anupam Gupta, Carnegie Mellon University, USA; Viswanath Nagarajan, University of Michigan, USA; Sahil Singla, Carnegie Mellon University, USA

3:40-4:00 Discrete Gaussian Sampling Reduces to CVP and SVP
Noah Stephens-Davidowitz, New York University, USA

Coffee Break
4:05 PM-4:30 PM
Room: Arlington Ballroom Foyer - 1st Floor

Tuesday, January 12

CP26
Session 8D
2:00 PM-4:05 PM

Room: Rosslyn - Lobby Level
Chair: Yusuke Kobayashi, University of Tsukuba, Japan

2:00-2:20 An Algorithmic Hypergraph Regularity Lemma
Brendan Nagle, University of South Florida, USA; Vojtech Rödl, Emory University, USA; Mathias Schacht, Universität Hamburg, Germany

2:25-2:45 Obstructions for Three-Coloring Graphs with One Forbidden Induced Subgraph
Maria Chudnovsky, Princeton University, USA; Jan Goedgebeur, Ghent University, Belgium; Oliver Schaudt, University of Cologne, Germany; Mingxian Zhong, Columbia University, USA

2:50-3:10 Independence and Efficient Domination on P6-Free Graphs
Daniel Lokshtanov, University of Bergen, Norway; Marcin Pilipczuk, University of Warsaw, Poland; Erik Jan van Leeuwen, Max Planck Institute for Informatics, Germany

3:15-3:35 Sparsity and Dimension
Gwenaël Joret, Université Libre de Bruxelles, Belgium; Piotr Micek, Jagiellonian University, Poland; Veit Wiechert, Technical University Berlin, Germany

3:40-4:00 Finding Perfect Matchings in Bipartite Hypergraphs
Chidambaram Annamalai, ETH Zürich, Switzerland

Tuesday, January 12

CP27
Session 9A
4:30 PM-6:35 PM

Room: Arlington Ballroom Salon IV - 1st Floor
Chair: Aravindan Vijayaraghavan, New York University, USA

4:30-4:50 Using Optimization to Obtain a Width-Independent, Parallel, Simpler, and Faster Positive Sdp Solver
Zeyuan Allen-Zhu, Princeton University, USA; Yin Tat Lee, Massachusetts Institute of Technology, USA; Lorenzo Orecchia, Boston University, USA

4:55-5:15 How to Round Subspaces: A New Spectral Clustering Algorithm
Ali K. Sinop, University of Illinois at Urbana-Champaign, USA

5:20-5:40 Improved Cheeger’s Inequality Using Vertex Expansion and Expansion Profile
Tsz Chiu Kwok, École Polytechnique Fédérale de Lausanne, Switzerland; Lap Chi Lau, University of Waterloo, Canada; Yin Tat Lee, Massachusetts Institute of Technology, USA

5:45-6:05 Approximate Undirected Maximum Flows in O(m Polylog(n)) Time
Richard Peng, Georgia Institute of Technology, USA

6:10-6:30 Natural Algorithms for Flow Problems
Damian Straszak and Nisheeth K. Vishnoi, École Polytechnique Fédérale de Lausanne, Switzerland
Tuesday, January 12

**CP28**

**Session 9B**

4:30 PM-6:35 PM

*Room:* Arlington Ballroom Salon V - 1st Floor

Chair: Eric Price, University of Texas at Austin, USA

- **4:30-4:50** Efficient Low-Redundancy Codes for Correcting Multiple Deletions
  Joshua Brakensiek, Venkatesan Guruswami, and Samuel Zbarsky, Carnegie Mellon University, USA

- **4:55-5:15** An Improved Bound on the Fraction of Correctable Deletions
  Boris Bukh and Venkatesan Guruswami, Carnegie Mellon University, USA

- **5:20-5:40** Communication Complexity of Permutation-Invariant Functions
  Badih Ghazi and Pritish Kamath, Massachusetts Institute of Technology, USA; Madhu Sudan, Harvard University, USA

- **5:45-6:05** Towards Optimal Deterministic Coding for Interactive Communication
  Ran Gelles, Princeton University, USA; Bernhard Haeupler, Carnegie Mellon University, USA; Gillat Kol, Institute for Advanced Study, USA; Noga Ron-Zewi, Institute for Advanced Study and Rutgers University, USA; Avi Wigderson, Institute for Advanced Studies, USA

- **6:10-6:30** Robust Positioning Patterns
  Ross Berkowitz and Swastik Kopparty, Rutgers University, USA

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**Tuesday, January 12**

**CP29**

**Session 9C**

4:30 PM-6:35 PM

*Room:* Arlington Ballroom Salon VI - 1st Floor

Chair: David Kempe, University of Southern California, USA

- **4:30-4:50** Sampling on Lattices with Free Boundary Conditions Using Randomized Extensions
  Sarah Cannon and Dana Randall, Georgia Institute of Technology, USA

- **4:55-5:15** On the Switch Markov Chain for Perfect Matchings
  Martin E. Dyer, University of Leeds, United Kingdom; Mark Jerrum, Queen Mary, University of London, United Kingdom; Haiko Muller, University of Leeds, United Kingdom

- **5:20-5:40** Partial Resampling to Approximate Covering Integer Programs
  David Harris, University of Maryland, USA

- **5:45-6:05** Algorithmic and Enumerative Aspects of the Moser-Tardos Distribution
  David Harris, University of Maryland, USA

- **6:10-6:30** Focused Stochastic Local Search and the Lovász Local Lemma
  Dimitris Achlioptas, University of California, Santa Cruz, USA; Fotis Iliopoulos, University of California, Berkeley, USA

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**Tuesday, January 12**

**CP30**

**Session 9D**

4:30 PM-6:35 PM

*Room:* Rosslyn - Lobby Level

Chair: Philip Klein, Brown University, USA

- **4:30-4:50** The k-Mismatch Problem Revisited
  Raphael Clifford and Allyx Fontaine, University of Bristol, United Kingdom; Ely Porat, Bar-Ilan University, Israel; Benjamin G. Sach and Tatiana Starikovskaya, University of Bristol, United Kingdom

- **4:55-5:15** Range Predecessor and Lempel-Ziv Parsing
  Djamal Belazzougui and Simon J. Puglisi, University of Helsinki, Finland

- **5:20-5:40** Communication with Contextual Uncertainty
  Badih Ghazi, Massachusetts Institute of Technology, USA; Ilan Komargodski, Weizmann Institute of Science, Israel; Pravesh Kothari, University of Texas at Austin, USA; Madhu Sudan, Harvard University, USA

- **5:45-6:05** Weak Duality for Packing Edge-Disjoint Odd (u,v)-Trails
  Ross Churchley and Bojan Mohar, Simon Fraser University, Canada; Hehui Wu, University of Mississippi, USA

- **6:10-6:30** On the Maximum Quartet Distance Between Phylogenetic Trees
  Noga Alon, Tel Aviv University, Israel; Humberto Naves, University of Minnesota, USA; Benjamin Sudakov, ETH Zürich, Switzerland
Abstracts for ALENEX and ANALCO appear in the beginning of the section, under “CP0”, and are alphabetized by the speaker’s last name.

Abstracts are printed as submitted by the authors.
IP1
Distributed Algorithms in Ant Colonies

Like many distributed systems, both natural and engineered, ant colonies operate without any central control. No ant can assess what needs to be done. Each ant responds to its interactions with other ants nearby and in the aggregate, these stochastic, dynamical networks of interaction regulate colony behavior. Ant species use distributed algorithms that differ according to environmental constraints such as operating costs, time pressure, patchy distribution of resources in space, risk of rupture, and the intensity of competition. I will discuss algorithms used to regulate foraging activity, nestmate recognition and colony security, collective search, and stable trail systems.

Deborah M. Gordon
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IP2
Beyond Planarity: On Geometric Intersection Graphs

Many efficient algorithms had been developed over the years for planar graphs and more general graphs such as low genus graphs. Intersection graphs of geometric objects (in low dimensions) with some additional proprieties, such as fatness or low density, provide yet another family of graphs for which one can design better algorithms. This family is a vast extension of planar graphs, and yet is still algorithmically tractable for many problems. In this talk, we will survey this class of graphs, and some algorithms and intractability results known for such graphs, and outline open problems for further research.

Sariel Har-Peled
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IP3
Combinatorial Statistics

Combinatorial Statistics studies combinatorial algorithms for inference of statistical models that are parametrized by discrete parameters such as trees, graphs or permutations. The talk will outline some key examples of results and techniques that are used in this area as well as future research directions and open problems.

Elchanan Mossel
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CP0
A Novel Dual Ascent Algorithm for Solving the Min-Cost Flow Problem

We present a novel algorithm for the min-cost flow problem that is competitive with recent third-party implementations of well-known algorithms for this problem and even outperforms them on certain realistic instances. We formally prove correctness of our algorithm and show that the worst-case running time is in $O(||b||_1(m + n \log n))$ where $b$ is the vector of demands. Combined with standard scaling techniques, this pseudo-polynomial bound can be made polynomial in a straightforward way. Furthermore, we evaluate our approach experimentally. Our empirical findings indeed suggest that the running time does not significantly depend on the costs and that a linear dependence on $||b||_1$ is overly pessimistic.

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Centrality indices are widely used analytic measures for the importance of nodes in a network. Closeness centrality is very popular among these measures. For a single node $v$, it takes the sum of the distances of $v$ to all other nodes into account. The currently best algorithms in practical applications for computing the closeness for all nodes exactly in unweighted graphs are based on breadth-first search (BFS) from every node. Thus, even for sparse graphs, these algorithms require quadratic running time in the worst case, which is prohibitive for large networks. In many relevant applications, however, it is unnecessary to compute closeness values for all nodes. Instead, one requires only the $k$ nodes with the highest closeness values in descending order. Thus, we present a new algorithm for computing this top-$k$ ranking in unweighted graphs. Following the rationale of previous work, our algorithm significantly reduces the number of traversed edges. It does so by computing upper bounds on the closeness and stopping the current BFS search when $k$ nodes already have higher closeness than the bounds computed for the other nodes. In our experiments with real-world and synthetic instances of various types, one of these new bounds is good for small-world graphs with low diameter (such as social networks), while the other one excels for graphs with high diameter (such as road networks). Combining them yields an algorithm that is faster than the state of the art for top-$k$ centrality in unweighted graphs. Finally, we prove that the quadratic worst-case complexity cannot be improved on directed, disconnected graphs, under reasonable complexity assumptions.

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Henning Meyerhenke
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Sampling Weighted Perfect Matchings on the Square-Octagon Lattice

We consider perfect matchings of the square-octagon lattice, also known as ‘fortresses’. There is a natural local Markov chain on the set of perfect matchings that is known to be ergodic. However, unlike seemingly related Markov chains on perfect matchings on the square and hexagonal lattices, corresponding to domino and lozenge tilings, respectively, the Markov chain on the square-octagon lattice appears to converge slowly. To understand why, we consider a weighted version of the problem. As with domino and lozenge tilings, it will be useful to view perfect matchings on the square-octagon lattice in terms of sets of paths and cycles on a corresponding lattice region; here, the paths and cycles lie on the Cartesian lattice and are required to turn left or right at every step. For input parameters $\lambda$ and $\mu$, we define the weight of a configuration to be $\lambda E(\sigma) + \mu V(\sigma)$, where $E(\sigma)$ is the total number of edges on the paths and cycles of $\sigma$ and $V(\sigma)$ is the number of vertices that are not incident to any of the paths or cycles in $\sigma$. Weighted paths already come up in the reduction from perfect matchings to turning lattice paths, corresponding to the case when $\lambda = 1$ and $\mu = 2$. First fixing $\mu = 1$, we show that there are choices of $\lambda$ for which the chain converges slowly and another for which it is fast, suggesting a phase change in the mixing time. More precisely, the chain requires exponential time (in the size of the lattice region) when $\lambda < 1/(2\sqrt{\sigma})$ or $\lambda > 2\sqrt{\sigma}$, while it is polynomially mixing at $\lambda = 1$. When $\mu > 1$, we show that the Markov chain is slowly mixing if $\lambda < \sqrt{\mu}/(2\sqrt{\sigma})$ or $\lambda > 2\mu\sqrt{\sigma}$. These are the first rigorous proofs explaining why the natural local Markov chain can be slow for the weighted fortress model, or perfect matchings on the square-octagon lattice.

Prateek Bhakta, Dana Randall
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Fast Algorithms for Pseudoarboricity

We show that the pseudoarboricity $p$ of a graph can be computed in $\tilde{O}(\sqrt{|E|/\log p})$ time, and give better runtime estimates for certain asymptotic bounds on $p$. These results are achieved by accelerating a binary search with an approximation scheme, and a runtime analysis of Dinitz’s algorithm on ‘almost unit capacity’ networks. We experimentally compare several algorithms for the problem.

Markus Blumenstock
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Graphs with Degree Constraints

Given a set $D$ of nonnegative integers, we derive the asymptotic number of graphs with a given number of vertices, edges, and such that the degree of every vertex is in $D$. This generalizes existing results, such as the enumeration of graphs with a given minimum degree, and establishes new ones, such as the enumeration of Euler graphs, $i.e.$ where all vertices have an even degree. Those results are derived using analytic combinatorics.

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An Empirical Comparison of Graph Laplacian Solvers

Solving Laplacian linear systems is an important task in a variety of practical and theoretical applications. This problem is known to have solutions that perform in linear times polylogarithmic work in theory, but these algorithms are difficult to implement in practice. We examine existing solution techniques in order to determine the best methods currently available and for which types of problems they are useful.

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Betweenness Centrality in Random Trees

The betweenness centrality of a node is a quantity that is frequently used in network theory to measure how ‘central’ a node $v$ is. It is defined as the summed proportion of shortest paths between pairs of vertices that pass through $v$. We study the average and limiting distribution of the betweenness centrality in random trees and, related, sub-critical graph families.

Kevin Durant
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Experimental Evaluation of Distributed Node Coloring Algorithms for Wireless Networks

In this paper we evaluate distributed node coloring algorithms for wireless networks using the network simulator Sinalgo. All considered algorithms operate in the realistic signal-to-interference-and-noise-ratio (SINR) model of interference. We evaluate two recent coloring algorithms,
**CP0**

**Scaling Limit of Random \( k \)-Trees**

We consider a random \( k \)-tree \( G_{n,k} \) that is uniformly selected from the class of labelled \( k \)-trees with \( n + k \) vertices. Since 1-trees are just trees, it is well-known that \( G_{n,1} \) (after scaling the distances by 1/\( 2\sqrt{n} \)) converges to the Continuum Random Tree \( T_e \). Our main result is that for \( k \neq 1 \), the random \( k \)-tree \( G_{n,k} \), scaled by \( (kH_{k-1} + 1)/(2\sqrt{n}) \) where \( H_{k-1} \) is the \((k - 1)\)th Harmonic number, converges to the Continuum Random Tree \( T_e \), too. In particular this shows that the diameter as well as the typical distance of two vertices in a random \( k \)-tree \( G_{n,k} \) are of order \( \sqrt{n} \).

**CP0**

**A General Framework for Dynamic Succinct and Compressed Data Structures**

Succinct data structures are becoming increasingly popular in big data processing applications due to their low memory consumption. In this paper we design, implement, and test a general framework that allows for practical dynamic succinct structures. We first describe implementations of compressed modifiable bit vectors, and extended compressed random access memory. Then, we implement and test our data structures using several popular compression libraries, and both synthetic data and a real-world temporal graph.

**DA16 Abstracts**

**Random Tree**

**Random k-Trees**

**Randomized Strategies for Cardinality Robustness in the Knapsack Problem**

We address randomized strategies for the following zero-sum game related to the knapsack problem. Given an instance of the knapsack problem, Alice chooses a knapsack solution and Bob, knowing Alice’s solution, chooses a cardinality \( k \). Then, Alice obtains a payoff equal to the ratio of the profit of the best \( k \) items in her solution to that of the best solution of size at most \( k \). We first show an instance such that the payoff of an arbitrary randomized strategy is both \( O(\log \log \mu/\log \mu) \) and \( O(\log \log \rho/\log \rho) \), where \( \mu \) is the exchangeability of the independence system and \( \rho \) is the ratio of the size of a maximum feasible set to that of minimum infeasible set minus one. We then design two randomized strategies with payoff \( \Omega(1/\log \mu) \) and \( \Omega(1/\log \rho) \) guaranteed, respectively, which substantially improve upon that of deterministic strategies and almost attain the above upper bounds.

**Engineering Oracles for Time-Dependent Road Networks**

We implement and experimentally evaluate landmark-based oracles for min-cost paths in large-scale time-dependent road networks. We exploit parallelism and lossless compression, combined with a novel travel-time approximation technique, to severely reduce preprocessing space and time. We significantly improve the FLAT oracle, improving the previous query-time by 30% and doubling the Dijkstra-rank speedup. We also implement and experimentally evaluate a novel oracle (HORN), based on a landmark hierarchy, achieving even better performance wrt to FLAT.

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CP0

An Algorithm for Online K-Means Clustering

Abstract not available.

Edo Liberty
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CP0

More Analysis of Double Hashing for Balanced Allocations

With double hashing, for a key $x$, one generates two hash values $f(x)$ and $g(x)$, and then uses combinations $(f(x) + ig(x)) \mod n$ for $i = 0, 1, 2, \ldots$ to generate multiple hash values in the range $[0, n-1]$ from the initial two. For balanced allocations, keys are hashed into a hash table where each bucket can hold multiple keys, and each key is placed in the least loaded of $d$ choices. It has been shown previously that asymptotically the performance of double hashing and fully random hashing is the same in the balanced allocation paradigm using fluid limit methods. Here we extend a coupling argument used by Lueker and Molodowitch to show that double hashing and ideal uniform hashing are asymptotically equivalent in the setting of open address hash tables to the balanced allocation paradigm, providing further insight into this phenomenon. We also discuss the potential for and bottlenecks limiting the use this approach for other multiple choice hashing schemes.

Michael Mitzenmacher
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CP0

A New Approach to Analyzing Robin Hood Hashing

Robin Hood hashing is a variation on open addressing hashing designed to reduce the maximum search time as well as the variance in the search time for elements in the hash table. While the case of insertions only using Robin Hood hashing is well understood, the behavior with deletions has remained open. Here we show that Robin Hood hashing with random hash functions can be analyzed under the framework of finite-level finite-dimensional jump Markov chains. This framework allows us to re-derive some past results for the insertion-only case with some new insight, as well as provide new analyses for a standard deletion model, where we alternate between deleting a random old key and inserting a new one.

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CP0

The Clt Analogue for Cyclic Urns

A cyclic urn is an urn model for balls of types $0, \ldots, m-1$ where in each draw the ball drawn, say of type $j$, is returned to the urn together with a new ball of type $j+1 \mod m$. The case $m = 2$ is the well-known Friedman urn. The composition vector, i.e., the vector of the numbers of balls of each type after $n$ steps is, after normalization, known to be asymptotically normal for $2 \leq m \leq 6$. For $m \geq 7$ the normalized composition vector does not converge. However, there is an almost sure approximation by a periodic random vector. In this paper the asymptotic fluctuations around this periodic random vector are identified. We show that these fluctuations are asymptotically normal for all $m \geq 7$. However, they are of maximal dimension $m-1$ only when 6 does not divide $m$. For $m$ being a multiple of 6 the fluctuations are supported by a two-dimensional subspace.

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CP0

Geometry Helps to Compare Persistence Diagrams

Exploiting geometric structure to improve the asymptotic complexity of discrete assignment problems is a well-studied subject. In contrast, the practical advantages of using geometry for such problems have not been explored. We implement geometric variants of the Hopcroft–Karp algorithm for bottleneck matching (based on previous work by Efrat et al.), and of the auction algorithm by Bertsekas for Wasserstein distance computation. Both implementations use k-d trees to replace a linear scan with a geometric proximity query. Our interest in this problem stems from the desire to compute distances between persistence diagrams, a problem that comes up frequently in topological data analysis. We show that our geometric matching algorithms lead to a substantial performance gain, both in running time and in memory consumption, over their purely combinatorial counterparts. Moreover, our implementation significantly outperforms the only other implementation available for comparing persistence diagrams.

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CP0

Generating Massive Scale Free Networks under Resource Constraints

Preferential attachment is a popular building block in generators for scale-free random graphs. We present the first two I/O-efficient node sampling primitives for preferen-
tial attachment with linear fitness functions and evaluate their performance for Barabási-Albert graphs. One approach is parallelizable: Our GPGPU-accelerated generator is 14.4 times faster than a sequential linear-time scheme by Batagelj and Brandes for large graphs in main memory and scales well beyond the RAM size. Both schemes support arbitrary seed graphs exceeding main memory, ran-

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CP0
k-Way Hypergraph Partitioning Via n-Level Recursive Bisection

We develop a multilevel algorithm for hypergraph partitioning that contracts the vertices one at a time. Using several caching and lazy-evaluation techniques during coarsening and refinement, we reduce the running time by up to two-orders of magnitude compared to a naive n-level algorithm that would be adequate for ordinary graph partitioning. The overall performance is even better than the widely used hMetis hypergraph partitioner that uses a classical multilevel algorithm with few levels. Aided by a portfolio-based approach to initial partitioning and adaptive budgeting of imbalance within recursive bipartitioning, we achieve very high quality. We assembled a large benchmark set with 310 hypergraphs stemming from application areas such VLSI, SAT solving, social networks, and scientific computing. Experiments indicate that our algorithm is the method of choice for a wide range of hypergraph partitioning tasks. The algorithm presented in this work forms the basis of our hypergraph partitioning framework KaHyPar (Karlsruhe Hypergraph Partitioning).

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CP0
Total Variation Discrepancy of Deterministic Random Walks for Ergodic Markov Chains

Motivated by a derandomization of Markov chain Monte Carlo, this paper investigates deterministic random walks, which is a deterministic process analogous to a random walk. This paper gives upper bounds of the total variation discrepancy between the expected number of tokens in a Markov chain and the number of tokens in its corresponding deterministic random walk, in terms of the mixing time of the Markov chain. We also present some lower bounds.

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CP0
Scalable Transfer Patterns

We consider the problem of Pareto-optimal route planning in public-transit networks of a whole country, a whole continent, or even the whole world. On such large networks, existing approaches suffer from either a very large space consumption, a very long preprocessing time or slow query processing. Transfer Patterns, a state-of-the-art technique for route planning in transit networks, achieves excellent query times, but the space consumption is large and the preprocessing time is huge. In this paper, we introduce a new scheme for the Transfer Pattern precomputation and query graph construction that reduces both the necessary preprocessing time and space consumption by an order of magnitude and more. Average query times are below 1ms for local queries, independent of the size of the network,
around 30ms for non-local queries on the complete transit network of Germany, and an estimated 200ms for a fictitious transit network covering the currently available data of the whole world.

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CP0

Finding Near-Optimal Independent Sets at Scale

The independent set problem is NP-hard and particularly difficult to solve in large sparse graphs which typically take exponential time to solve exactly using the best-known exact algorithms. In this work, we develop an advanced evolutionary algorithm, which incorporates kernelization techniques to compute large independent sets in huge sparse networks. A recent exact algorithm has shown that large networks can be solved exactly by employing a branch-and-reduce technique that recursively kernelizes the graph and performs branching. However, one major drawback of their algorithm is that, for huge graphs, branching still can take exponential time. To avoid this problem, we recursively choose vertices that are likely to be in a large independent set (using an evolutionary approach), then further kernelize the graph. We show that identifying and removing vertices likely to be in large independent sets opens up the reduction space—which not only speeds up the computation of large independent sets drastically, but also enables us to compute high-quality independent sets on much larger instances than previously reported in the literature.

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CP0

Graph Bisection with Pareto-Optimization

We introduce FlowCutter, a novel algorithm to compute a set of edge cuts or node separators that optimize cut size and balance in the Pareto-sense. Using the computed Pareto-set we can identify cuts with a particularly good trade-off between cut size and balance that can be used to compute contraction and minimum fill-in orders, which can be used in Customizable Contraction Hierarchies, a speed-up technique for shortest path computations.

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CP0

Prime Factorization of the Kirchhoff Polynomial: Compact Enumeration of Arborescences

We study the problem of enumerating all arborescences of a directed graph $G$. They can be expressed as the Kirchhoff polynomial, a polynomial over the edge labels of $G$, in which each monomial represents a single arborescence. We show how to compute, in linear time, a compact representation of the Kirchhoff polynomial—their prime factorization, and how it relates to combinatorial properties of digraphs such as strong connectivity and vertex domination.

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CP0

Ricci-Ollivier Curvature of the Rooted Phylogenetic Subtree-Prune-Regraft Graph

We investigated the Ricci-Ollivier curvature of the graph induced by subtree-prune-regraft moves on rooted tree topologies, such as those explored by statistical phylogenetic inference methods. We developed fast new algorithms for constructing and sampling these graphs. We confirm using simulation that access time distributions of random walks on these graphs depend on distance, degree, and curvature. We then study degree changes in these graphs theoretically and give bounds on the curvature.

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On Terminal Nodes and the Degree Profile of Preferential Dynamic Attachment Circuits

We investigate terminal nodes and the degree profile in preferential dynamic attachment circuits. We study the distribution of the number of terminal nodes, which are the nodes that have not recruited other nodes, as the circuit ages. The expectation and variance of the number of terminal nodes are both linear with respect to the age of the circuit. We show via martingale that the number of terminal nodes asymptotically follows a Gaussian law. We also study the exact distribution of the degree of a specific node as the circuit grows. The exact expectation and variance of the degree of a node are determined via a series of Polya-Eggenberger urn models with “hiccups” in between and recurrence methods. Phase transitions of these degrees are discussed briefly.

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On Connectivity in a General Random Intersection Graph

There has been growing interest recently in studies of general random intersection graphs. In this paper, we consider a general random intersection graph \( G(n, \vec{a}, \vec{K}, P_n) \) defined on a set \( V_n \) comprising \( n \) vertices, where \( \vec{a} \) is a probability vector \((a_1, a_2, \ldots, a_m)\) and \( \vec{K} = (K_1, n, K_2, n, \ldots, K_m, n) \). Suppose there is a pool \( P_n \) consisting of \( P_n \) distinct objects. The \( n \) vertices in \( V_n \) are divided into \( m \) groups \( A_1, A_2, \ldots, A_m \). Each vertex \( v \) is independently assigned to exactly a group according to the probability distribution with \( P[v \in A_i] = a_i \), where \( i = 1, 2, \ldots, m \). Afterwards, each vertex in group \( A_i \) independently chooses \( K_{i,n} \) objects uniformly at random from the object pool \( P_n \). Finally, an undirected edge is drawn between two vertices in \( V_n \) that share at least one object. This graph model \( G(n, \vec{a}, \vec{K}, P_n) \) has applications in secure sensor networks and social networks. We investigate connectivity in this general random intersection graph \( G(n, \vec{a}, \vec{K}, P_n) \) and present a sharp zero-one law.

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Undirected Graph Exploration with \( \Theta(\log \log N) \) Pebbles

We consider the fundamental problem of exploring an undirected and initially unknown graph by an agent with little memory. The vertices of the graph are unlabeled, and the edges incident to a vertex have locally distinct labels. In this setting, it is known that \( \Theta(\log n) \) bits of memory are necessary and sufficient to explore any graph with at most \( n \) vertices. We show that this memory requirement can be decreased significantly by making a part of the memory distributable in the form of pebbles. A pebble is a device that can be dropped to mark a vertex and can be collected when the agent returns to the vertex. We show that for an agent \( \Omega(\log \log n) \) distinguishable pebbles and bits of memory are sufficient to explore any bounded-degree graph with at most \( n \) vertices. We match this result with a lower bound exhibiting that for any agent with sub-logarithmic memory, \( \Omega(\log \log n) \) distinguishable pebbles are necessary for exploration.

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New Directions in Nearest Neighbor Searching with Applications to Lattice Sieving

To solve the approximate nearest neighbor search problem (NNS) on the sphere, we propose a method using locality-sensitive filters (LSF), with the property that nearby vectors have a higher probability of surviving the same filter than vectors which are far apart. We instantiate the filters using spherical caps of height \( 1-\alpha \), where a vector survives a filter if it is contained in the corresponding spherical cap, and where ideally each filter has an independent, uniformly random direction. For small \( \alpha \), these filters are very similar to the spherical locality-sensitive hash (LSH) family previously studied by Andoni et al. For larger \( \alpha \) bounded away from 0, these filters potentially achieve a superior performance, provided we have access to an efficient oracle for finding relevant filters. Whereas existing LSH schemes are limited by a performance parameter of \( \rho \geq 1/(2c^2-1) \) to solve approximate NNS with approximation factor \( c \), these filters potentially achieve smaller asymptotic values of \( \rho \), depending on the density of the data set. For sparse data sets where the dimension is super-logarithmic in the size of the data set, we asymptotically obtain \( \rho \sim 1/(4kc^2) \), which is better than vectors which are far apart. We instantiate the filters, we replace the independent filters by filters taken from certain structured random product codes. We show that the additional structure in these concatenation codes allows us to decode efficiently using techniques similar to lattice enumeration, and we can find the relevant filters with low overhead, while at the same time not significantly changing the collision probabilities of the filters. We finally apply spherical LSF to sieving algorithms for solving the shortest vector problem (SVP) on lattices, and show that this leads to a time complexity for solving SVP in dimension \( n \) of \((3/2)^{n/2+o(n)}\).

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CP1
Locality-Sensitive Hashing Without False Negatives

We consider a new construction of locality-sensitive hash functions for Hamming space that is covering in the sense that it guarantees to produce a collision for every pair of vectors within a given radius $r$. The construction is efficient in the sense that the expected number of hash collisions between vectors at distance $cr$, for a given $c > 1$, comes close to that of the best possible data independent LSH without the covering guarantee, namely, the seminal LSH construction of Indyk and Motwani (FOCS ’98). The efficiency of the new construction essentially matches their bound if $cr = \log(n)/k$, where $n$ is the number of points in the data set and $k \in \mathbb{N}$, and differs from it by at most a factor $\ln(4) < 1.4$ in the exponent for general values of $cr$. As a consequence, LSH-based similarity search in Hamming space can avoid the problem of false negatives at little or no cost in efficiency.

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CP1
Phase Transitions in Group Testing

We study the fundamental limits of any group testing procedure regardless of its computational complexity. In the noiseless case with the number of defective items $k$ scaling with the total number of items $p$ as $O(p^\theta)$ ($\theta \in (0, 1)$), we show that the probability of reconstruction error tends to one when $n \leq k \log_2 \frac{1}{\theta}(1 + o(1))$, but vanishes when $n \geq c(\theta)k \log_2 \frac{1}{\theta}(1 + o(1))$, for some explicit constant $c(\theta)$. For $\theta \leq \frac{1}{3}$, we show that $c(\theta) = 1$, thus proving a phase transition. Analogous conditions are derived for noisy and partial recovery settings.

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CP1
Tight Bounds for the Distribution-Free Testing of Monotone Conjunctions

We study the distribution-free testing of monotone conjunctions. We present an $\tilde{O}(n^{1/3}/\epsilon^5)$-query algorithm that tests whether an unknown $f$ is a monotone conjunction versus $\epsilon$-far from monotone conjunctions with respect to an unknown distribution. This improves the $\tilde{O}(n^{1/2}/\epsilon)$ upper bound by Dolev and Ren. We also prove a lower bound of $\Omega(n^{1/3})$, improving the $\tilde{\Omega}(n^{1/5})$ lower bound by Glasner and Servedio. Our bounds are tight, up to a polylogarithmic factor, when $\epsilon$ is a constant.

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CP2
The Complexity of All-Switches Strategy Improvement

Strategy improvement is a widely-used and well-studied class of algorithms for solving graph-based infinite games. We show that two natural problems are PSPACE-complete for the all-switches variant of strategy improvement: the problem of determining whether a given edge is switched, and the problem of determining which optimal strategy is found. These results hold for parity games games, mean-payoff games, discounted-payoff games, and simple-stochastic games. We also show related results for acyclic unique sink orientations.

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CP2
Characterisation of Strongly Stable Matchings

We present a characterisation of the set of all strongly stable matchings, thus solving an open problem already stated in 1989 in the book by Gusfield and Irving. Although the number of strongly stable matchings can be exponential, we show that there exists a partial order with $O(m)$ elements representing all strongly stable matchings. We give two algorithms that construct two such representations: one in $O(nm^2)$ time and the other in $O(nm)$ time.

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CP2
An Improved Combinatorial Polynomial Algorithm for the Linear Arrow-Debreu Market

We present an improved combinatorial algorithm for the computation of equilibrium prices in the linear Arrow-Debreu model. For a market with $n$ agents and integral utilities bounded by $U$, the algorithm runs in $O(n^2 \log^3(nU))$ time. This improves upon the previously best algorithm of Ye by a factor of $\tilde{\Omega}(n)$. The algorithm refines the algorithm described by Duan and Mehlhorn and improves it by a factor of $\tilde{\Omega}(n^3)$. The improvement comes from a better understanding of the iterative price adjustment process, the improved balanced flow computation for nondegenerate instances, and a novel perturbation technique for achieving nondegeneracy.

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Previous work has left open the question of whether an online algorithm can be $O(1)$-competitive with $O(1)$-speed for average flow time in the DAG setting. In this work, we answer this question positively by giving a scalable algorithm which is $(1 + \epsilon)$-speed $O(\frac{1}{\epsilon})$-competitive for any $\epsilon > 0$. We further introduce the first greedy algorithm for scheduling parallelizable jobs — our algorithm is a generalization of the shortest jobs first algorithm. Greedy algorithms are among the most useful in practice due to their simplicity. We show that this algorithm is $(2 + \epsilon)$-speed $O(\frac{1}{\epsilon})$-competitive for any $\epsilon > 0$. 

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CP3
Online Pricing with Impatient Bidders

In this paper we consider the following online pricing problem. An auctioneer is selling identical items in unlimited supply, whereas each bidder from a given set is interested in purchasing a single copy of the item. Each bidder is characterized by a budget and a time interval, in which he is considering to buy the item. Bidders are willing to buy the item at the earliest time provided it is within their time intervals and the price at that time is within their budgets. We call such bidders impatient bidders. The problem is considered in the online setting, i.e., each bidder arrives at the start of his time interval, and only then an algorithm learns of his existence and his budget. The goal of the seller is to set the price of the item over time so that the total revenue is maximized. We study two versions of the impatient bidders problem: the one introduced by Bansal et al. [TALG’10], and a more restricted setting in which the deadline of each bidder remains unknown until it is hit. We give tight bounds for both settings. Rather surprisingly, in both cases the optimum competitive ratios are the same. In particular we prove that the competitive ratio of an optimum deterministic algorithm is $\Theta(\log h/\log \log h)$, whereas for randomized algorithms it is $\Theta(\log \log h)$.

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CP3
An $\mathcal{O}(\log m)$-Competitive Algorithm for Online Machine Minimization

We consider the online machine minimization problem in which jobs with hard deadlines arrive online over time at their release dates. The task is to determine a feasible preemptive schedule on a minimum number of machines. Our main result is a general $\mathcal{O}(\log m)$-competitive algorithm for the online problem, where $m$ is the optimal number of machines used in an offline solution. This is the first improvement on an intriguing problem in nearly two decades. To date, the best known result is a $\mathcal{O}(\log(p_{\text{max}}/p_{\text{min}}))$-competitive algorithm which is $(1 + \epsilon)$-speed $O(\frac{1}{\epsilon})$-competitive for any $\epsilon > 0$.
competitive algorithm by Phillips et al. (STOC 1997) that depends on the ratio of maximum and minimum job sizes, $p_{\text{max}}$ and $p_{\text{min}}$. Even for $m = 2$ no better algorithm was known. Our algorithm is in this case constant-competitive. When applied to laminar or agreeable instances, our algorithm achieves a competitive ratio of $O(1)$ even independently of $m$.

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CP3
Online Degree-Bounded Steiner Network Design

We initiate the study of degree-bounded network design problems in the online setting. The degree-bounded Steiner tree problem, which asks for a subgraph with minimum degree that connects a given set of vertices, is perhaps one of the most representative problems in this class. We design an intuitive greedy-like algorithm that achieves a competitive ratio of $O(\log n)$ where $n$ is the number of vertices. We show that no (randomized) algorithm can achieve a (multiplicative) competitive ratio $o(\log n)$; thus our result is asymptotically tight. Our results carry over to the more general degree-bounded Steiner forest problem as well.

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CP3
Make-to-Order Integrated Scheduling and Distribution

Production and distribution are fundamental operational functions in supply chains. The main challenge is to design algorithms that optimize operational performance by jointly scheduling production and delivery of customer orders. In this paper we study a model of scheduling customer orders on multiple identical machines and their distribution to customers afterwards. The goal is to minimize the total time from release to distribution plus total distribution cost to the customers. We design the first poly-logarithmic competitive algorithm for the problem, improving upon previous algorithms with linear competitive ratios. Our model generalizes two fundamental problems: scheduling of jobs on multiple identical machines (where the goal function is to minimize the total flow time) as well as the TCP Acknowledgment problem.

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CP4
Expanders Via Local Edge Flips

Designing distributed and scalable algorithms to improve network connectivity is a central topic in peer-to-peer networks. In this paper we focus on the following well-known problem: given an $n$-node $d$-regular network, we want to design a decentralized, local algorithm that transforms the graph into one that has good connectivity properties (low diameter, expansion, etc.). To this end, Mahlmann and Schindelhauer introduced the random “flip” transformation, where in each time step, a random pair of vertices that have an edge decide to ‘swap a neighbor’. They conjectured that performing $poly(d)\times n \log n$ such flips at random would convert any connected $d$-regular graph into a $d$-regular expander graph, with high probability. However, the best known upper bound for the number of steps is roughly $O(n^{17}d^{23})$, obtained via a delicate Markov chain comparison argument. Our main result is to prove that a natural instantiation of the random flip produces an expander in at most $O(n^{2d^2\sqrt{\log n}})$ steps, with high probability. Our argument uses a potential-function analysis based on the matrix exponential, together with the recent beautiful results on the higher-order Cheeger inequality of graphs. We also show that our technique can be used to analyze another well-studied random process known as the ‘random switch’, and show that it produces an expander in $O(nd)$ steps for $d = \Theta(\log n)$, with high probability.

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CP4
An Improved Distributed Algorithm for Maximal Independent Set

The Maximal Independent Set (MIS) problem is one of the basics in the study of locality in distributed graph algorithms. This paper presents an extremely simple randomized algorithm providing a near-optimal local complexity
for this problem, which incidentally, when combined with some known techniques, also leads to a near-optimal global complexity. Classical MIS algorithms of Luby [STOC’85] and Alon, Babai and Itai [JALG’86] provide the global complexity guarantee that, with high probability, all nodes terminate after $O(\log n)$ rounds. In contrast, our initial focus is on the local complexity, and our main contribution is to provide a very simple algorithm guaranteeing that each particular node $v$ terminates after $O(\log \deg(v) + \log 1/\epsilon)$ rounds, with probability at least $1 - \epsilon$. This degree-dependency is optimal, due to a lower bound of Kuhn, Moscibrodzka, and Wattenhofer [PODC’04]. Interestingly, this local complexity smoothly transitions to a global complexity: by adding techniques of Barenboim, Elkin, Pettie, and Schneider [FOCS’12; arXiv: 1202.1983v3], we get a randomized MIS algorithm with a high probability global complexity of $O(\log \Delta) + 2^{O(\sqrt{\log \log \Delta})}$, where $\Delta$ denotes the maximum degree. This improves over the $O(\log^2 \Delta) + 2^{O(\sqrt{\log \log \Delta})}$ result of Barenboim et al., and gets close to the $\Omega(\min(\log \Delta, \sqrt{\log n}))$ lower bound of Kuhn et al. Corollaries include improved algorithms for MIS in graphs of upper-bounded arboricity, or lower-bounded girth, for Ruling Sets, for MIS in the Local Computation Algorithms (LCA) model, and a faster distributed algorithm for the Lovász Local Lemma.

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CP4

Distributed Algorithms for Planar Networks II: Low-Congestion Shortcuts, MST, and Min-Cut

This paper introduces the concept of low-congestion shortcuts for (near-)planar networks, and demonstrates their power by using them to obtain near-optimal distributed algorithms for problems such as Minimum Spanning Tree (MST) or Minimum Cut, in planar networks Consider a graph $G = (V, E)$ and a partitioning of $V$ into subsets of nodes $S_1, \ldots, S_N$, each inducing a connected subgraph $G[S_i]$. We define an $\alpha$-congestion shortcut with dilation $\beta$ to be a set of subgraphs $H_1, \ldots, H_N \subseteq G$, one for each subset $S_i$, such that:

1. For each $i \in [1, N]$, the diameter of the subgraph $G[S_i] + H_i$ is at most $\beta$.
2. For each edge $e \in E$, the number of subgraphs $G[S_i] + H_i$ containing $e$ is at most $\alpha$.

We prove that any partition of a $D$-dimensional planar graph into individually-connected parts admits an $O(D \log D)$-congestion shortcut with dilation $O(D \log D)$, which we prove to be nearly best-possible, and we also present a distributed construction of it in $\tilde{O}(D)$ rounds. Finally, we use low-congestion shortcuts, and their efficient distributed construction, to derive $O(D)$-round distributed algorithms for MST and Min-Cut, in planar networks. This complexity nearly matches the trivial lower bound of $\Omega(D)$. We remark that this is the first result bypassing the well-known $\Omega(D + \sqrt{n})$ existential lower bound of general graphs (see Peleg and Rubinovich [FOCS’99]; Elkin [STOC’04]; and Das Sarma et al. [STOC’11]) in a family of graphs of interest.

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CP4

The Adversarial Noise Threshold for Distributed Protocols

We consider the problem of implementing distributed protocols, despite adversarial channel errors, on synchronous-messaging networks with arbitrary topology. Our main result: Any $n$-party $T$-round protocol on an undirected communication network with $m$ edges can be compiled into an $O\left(\frac{m \log n T}{1 - \epsilon}\right)$-round simulation protocol that runs on an $O(n)$-edge subnetwork and tolerates an (optimal) adversarial error rate of $\Omega\left(\frac{\epsilon}{n}\right)$.

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CP4

Local-on-Average Distributed Tasks

A distributed task is local if its time complexity is (nearly) constant, otherwise it is global. Unfortunately, local tasks are relatively scarce, and most distributed tasks require time at least logarithmic in the network size (and often higher than that). In a dynamic setting, i.e., when the network undergoes repeated and frequent topological changes, such as vertex and edge insertions and deletions, it is desirable to be able to perform a local update procedure around the modified part of the network, rather than running a static global algorithm from scratch following each change. This paper makes a step towards establishing the hypothesis that many (statically) non-local distributed tasks are local-on-average in the dynamic setting, namely, their amortized time complexity is $O(\log^* n)$. Towards establishing the plausibility of this hypothesis, we propose a strategy for transforming static $O(\text{polylog}(n))$-time algorithms into dynamic $O(\log^* n)$ amortized time update procedures. We then demonstrate the usefulness of our strategy by applying it to several fundamental problems whose static time complexity is logarithmic, including forest-decomposition, edge-orientation and coloring sparse graphs, and show that their amortized time complexity in the dynamic setting is indeed $O(\log^* n)$.

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CP5

Nearly-Optimal Bounds for Sparse Recovery in Generic Norms, with Applications to k-Median Sketching

We initiate the study of trade-offs between sparsity and the number of measurements in sparse recovery schemes for generic norms. Specifically, for a norm $\| \cdot \|$, sparsity parameter $k$, approximation factor $K > 0$, and probability of failure $P > 0$, we ask: what is the minimal value of $m$ so...
that there is a distribution over $m \times n$ matrices $A$ with the property that for any $x$, given $Ax$, we can recover a $k$-sparse approximation to $x$ in the given norm with probability at least $1 - P$. We give a partial answer to this problem, by showing that for norms that admit efficient linear sketches, the optimal number of measurements $m$ is closely related to the doubling dimension of the metric induced by the norm $|| \cdot ||$ on the set of all $k$-sparse vectors. By applying our result to specific norms, we cast known measurement bounds in our general framework (for the $\ell_p$ norms, $p \in [1,2]$) as well as provide new, measurement-efficient schemes (for the Earth-Mover Distance norm). The latter result directly implies more succinct linear sketches for the well-studied planar $k$-median clustering problem. Finally, our lower bound for the doubling dimension of the EMD norm enables us to resolve the open question of [Frahling-Sohler, STOC’05] about the space complexity of clustering problems in the dynamic streaming model.

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CP5
Nearly Optimal Deterministic Algorithm for Sparse Walsh-Hadamard Transform

For every fixed constant $\alpha > 0$, we design an algorithm for computing the $k$-sparse Walsh-Hadamard transform (i.e., Discrete Fourier Transform over the Boolean cube) of an $N$-dimensional vector $x \in \mathbb{R}^N$ in time $k^{1+o}(\log N)^{O(1)}$. Specifically, the algorithm is given query access to $x$ and computes a $k$-sparse $\hat{x} \in \mathbb{R}^N$ satisfying $||\hat{x} - x||_1 \leq c||\hat{x} - H_k(\hat{x})||_1$, for an absolute constant $c > 0$, where $\hat{x}$ is the transform of $x$ and $H_k(\hat{x})$ is its best $k$-sparse approximation. Our algorithm is fully deterministic and only uses non-adaptive queries to $x$ (i.e., all queries are determined and performed in parallel when the algorithm starts). An important technical tool that we use is a construction of nearly optimal and linear lossless condensers which is a careful instantiation of the GUV condenser (Guruswami, Umans, Vadhan, JACM 2009). Moreover, we design a deterministic and non-adaptive $\ell_1/\ell_1$ compressed sensing scheme based on general lossless condensers that is equipped with a fast reconstruction algorithm running in time $k^{1+o}(\log N)^{O(1)}$ (for the GUV-based condenser) and is of independent interest. Our scheme significantly simplifies and improves an earlier expander-based construction due to Berinde, Gilbert, Indyk, Karloff, Strauss (Allerton 2008). Our methods use linear lossless condensers in a black box fashion; therefore, any future improvement on explicit constructions of such condensers would immediately translate to improved parameters in our framework (potentially leading to $k(\log N)^{O(1)}$ reconstruction time with a reduced exponent in the poly-logarithmic factor, and eliminating the extra parameter $\alpha$). By allowing the algorithm to use randomness, while still using non-adaptive queries, the running time of the algorithm can be improved to $O(k \log^3 N)$.

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CP5
Nearly Tight Oblivious Subspace Embeddings by Trace Inequalities

We present a new analysis of sparse oblivious subspace embeddings, based on the “matrix Chernoff” technique. These are probability distributions over (relatively) sparse matrices such that for any $d$-dimensional subspace of $\mathbb{R}^n$, the norms of all vectors in the subspace are simultaneously approximately preserved by the embedding with high probability—typically with parameters depending on $d$ but not on $n$. The families of embedding matrices considered here are essentially the same as those in [Nelson-Nguyen ‘13], but with better parameters (sparsity and embedding dimension). Because of this, our analysis essentially serves as a “drop-in replacement” for Nelson-Nguyen’s, improving bounds on its many applications to problems such as least squares regression and low-rank approximation. This new method is based on elementary tail bounds combined with matrix trace inequalities (Golden-Thompson or Lieb’s theorem), and does not require combinatorics, unlike the Nelson-Nguyen approach. There are also variants of this method that are even simpler, at the cost of worse parameters. Furthermore, the bounds obtained are much tighter than previous ones, matching known lower bounds up to a single log(d) factor in embedding dimension (previous results had more log factors and also had suboptimal trade-offs with sparsity).

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CP5
The Restricted Isometry Property of Subsampled Fourier Matrices

A matrix $A \in \mathbb{C}^{N \times N}$ satisfies the restricted isometry property of order $k$ with constant $\epsilon$ if it preserves the $\ell_2$ norm of all $k$-sparse vectors up to a factor of $1 \pm \epsilon$. We prove that a matrix $A$ obtained by randomly sampling $q = O(k \log^2 k \log N)$ rows from an $N \times N$ Fourier matrix satisfies the restricted isometry property of order $k$ with a fixed $\epsilon$ with high probability. This improves on Rudelson and Vershynin (Comm. Pure Appl. Math., 2008), its subsequent improvements, and Bourgain (GAFA Seminar Notes, 2014).

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CP6

Simpler, Faster and Shorter Labels for Distances in Graphs

We present a distance labeling scheme for undirected graphs with \( n \) nodes having labels of length \( \frac{\log n}{2} \) bits and constant decoding time. This outperforms all existing results with respect to both size and decoding time, including Winkler's (Combinatorica 1983) decade-old result, which uses labels of size \( \log 3n \) and \( O(n/\log n) \) decoding time, and Gavoille et al. (SODA'01), which uses labels of size \( 11n + o(n) \) and \( O(\log \log n) \) decoding time.

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CP6

New Bounds for Approximating Extremal Distances in Undirected Graphs

We provide new approximation bounds for the extremal distances (the diameter, the radius, and the eccentricities of all nodes) of an undirected graph with \( n \) nodes and \( m \) edges. First we show under the Strong Exponential Time Hypothesis (SETH) of Impagliazzo, Paturi and Zane [JCSS01] that it is impossible to get a \((3/2 - \varepsilon)\)-approximation of the diameter and a \((5/3 - \varepsilon)\)-approximation of all the eccentricities in \( O(m^{2-\delta}) \) time for any \( \varepsilon, \delta > 0 \), even allowing a constant additive term in the approximation. Second, we present an algorithmic scheme that gives a \((2 - 1/2^k)\)-approximation of the diameter and the radius and a \((3 - 4/2^k + 1)\)-approximation of all eccentricities in \( O(n m^{\frac{1}{2^k-1}}) \) expected time for any \( k \geq 0 \). For \( k \geq 2 \), this gives a family of previously unknown bounds, and approaches near-linear running time as \( k \) grows. Third, we observe a connection between the approximation of the diameter and the \( h \)-dominating sets, which are subsets of nodes at distance \( \leq h \) from every other node. We give bounds for the size of these sets, related with the diameter.

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CP6

Approximation and Fixed Parameter Subquadratic Algorithms for Radius and Diameter in Sparse Graphs

The radius and diameter are fundamental graph parameters, with several natural definitions for directed graphs. All versions of diameter and radius can be solved via solving all-pairs shortest paths (APSP), followed by a fast postprocessing step. However, solving APSP on \( n \)-node graphs requires \( \Omega(n^2) \) time even in sparse graphs. We study the question: when can diameter and radius in sparse graphs be solved in truly subquadratic time, and when is such an algorithm unlikely? Motivated by our conditional lower bounds on computing these measures exactly in truly subquadratic time, we search for approximation and fixed parameter subquadratic algorithms, and alternatively, for reasons why they do not exist. We find that:
- Most versions of Diameter and Radius can be solved in truly subquadratic time with \textit{optimal} approximation guarantees, under plausible assumptions. For example, there is a \( 2\)-approximation algorithm for directed Radius with one-way distances that runs in \( O(m \sqrt{n}) \) time, while a \((2 - \delta)\)-approximation algorithm in \( O(n^{2-\epsilon}) \) time is considered unlikely. - On graphs with treewidth \( k \), we can solve all versions in \( 2^{O(k \log k)} n^{1+o(1)} \) time. We show that these algorithms are near \textit{optimal} since even a \((3/2 - \delta)\)-approximation algorithm that runs in time \( 2^\Theta(k) n^{2-\epsilon} \) would refute plausible assumptions. Two conceptual contributions of this work that we hope will incite future work are: the introduction of a \textit{Fixed Parameter Tractability} in \textit{P} framework, and the statement of a differently-quantified variant of the Orthogonal Vectors Conjecture, which we call the \textit{Hitting Set Conjecture}.

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CP6

Approximate Distance Oracles for Planar Graphs with Improved Query Time-Space Tradeoff

We consider approximate distance oracles for edge-weighted \( n \)-vertex undirected planar graphs. Given fixed \( \epsilon > 0 \), we present a \((1+\epsilon)\)-approximate distance oracle with \( O(n (\log \log n)^2) \) space and \( O((\log \log n)^3) \) query time. This improves the previous best product of query time and space of the oracles of Thorup (FOCS 2001, J. ACM 2004) and Klein (SODA 2002) from \( O(n (\log n)) \) to \( O(n (\log \log n)^5) \).

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CP7

Deterministic Algorithms for Submodular Maximization Problems

Randomization is a fundamental tool used in many theoretical and practical areas of computer science. We study here the role of randomization in the area of submodular function maximization. In this area most algorithms
are randomized, and in almost all cases the approximation ratios obtained by current randomized algorithms are superior to the best results obtained by known deterministic algorithms. Derandomization of algorithms for general submodular function maximization seems hard since the access to the function is done via a value oracle. This makes it hard, for example, to apply standard derandomization techniques such as conditional expectations. Therefore, an interesting fundamental problem in this area is whether randomization is inherently necessary for obtaining good approximation ratios. In this work we give evidence that randomization is not necessary for obtaining good algorithms by presenting a new technique for derandomization of algorithms for submodular function maximization. Our high level idea is to maintain explicitly a (small) distribution over the states of the algorithm, and carefully update it using marginal values obtained from an extreme point solution of a suitable linear formulation. We demonstrate our technique on two recent algorithms for unconstrained submodular maximization and for maximizing submodular function subject to a cardinality constraint. In particular, for unconstrained submodular maximization we obtain an optimal deterministic 1/2-approximation showing that randomization is unnecessary for obtaining optimal results for this setting.

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CP7
Exact and Approximation Algorithms for Weighted Matroid Intersection

In this paper, we propose new exact and approximation algorithms for the weighted matroid intersection problem. Our exact algorithm is faster than previous algorithms when the largest weight is relatively small. Our approximation algorithm delivers a $(1 - \epsilon)$-approximate solution with a running time significantly faster than known exact algorithms. The core of our algorithms is a decomposition technique: we decompose an instance of the weighted matroid intersection problem into a set of instances of the unweighted matroid intersection problem. The computational advantage of this approach is that we can make use of fast unweighted matroid intersection algorithms as a black box for designing algorithms. Precisely speaking, we prove that we can solve the weighted matroid intersection problem via solving $W$ instances of the unweighted matroid intersection problem, where $W$ is the largest given weight. Furthermore, we can find a $(1 - \epsilon)$-approximate solution via solving $O((1/\epsilon)^2 \log r)$ instances of the unweighted matroid intersection problem, where $r$ is the smallest rank of the given two matroids. Our algorithms are simple and flexible: they can be adapted to special cases of the weighted matroid intersection problem.

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CP7
A Fast Approximation for Maximum Weight Matroid Intersection

Given two matroids defined over a common ground set of $n$ elements, let $k$ be the rank of the matroid intersection and let $Q$ denote the cost of an independence query for either matroid. We present a $(1 - \epsilon)$ approximation algorithm for the maximum weight matroid intersection problem with running time $O(nk \log^2(1/\epsilon)Q/\epsilon^2)$.

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CP7
Locally Adaptive Optimization: Adaptive Seeding for Monotone Submodular Functions

The Adaptive Seeding problem is an algorithmic challenge motivated by influence maximization in social networks: one seeks to select among certain accessible nodes in a network, and then select, adaptively, among neighbors of those nodes as they become accessible in order to maximize a global objective function. More generally, adaptive seeding is a stochastic optimization framework where the choices in the first stage affect the realizations in the second stage, over which we aim to optimize. Our main result is a $(1 - 1/\epsilon)^2$-approximation for the adaptive seeding problem for any monotone submodular function. While adaptive policies are often approximated via non-adaptive policies, our algorithm is based on a novel method we call locally-adaptive policies. These policies combine a non-adaptive global structure with local adaptive optimizations. This method enables the $(1 - 1/\epsilon)^2$-approximation for general monotone submodular functions and circumvents some of the impossibilities associated with non-adaptive policies. We also introduce a fundamental problem in submodular optimization that may be of independent interest: given a ground set of elements where every element appears with some small probability, find a set of expected size at most $k$ that has the highest expected value over the realization of the elements. We show a surprising result: there are classes of monotone submodular functions (including coverage) that can be approximated almost optimally as the probability vanishes. For general monotone submodular functions, we show via a reduction from Planted-Clique that approximations for this problem are not likely to be obtainable. This optimization problem is an important tool for adaptive seeding via non-adaptive policies, and its hardness motivates the introduction of locally-adaptive policies we use in the main result.

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Improved Approximation Algorithms for \( k \)-Submodular Function Maximization

This paper presents a polynomial-time \( 1/2 \)-approximation algorithm for maximizing nonnegative \( k \)-submodular functions. This improves upon the previous \( \max\{1/3, 1/(1 + a)\} \)-approximation by Ward and Živný, where \( a = \max\{1, \sqrt{(k-1)/4}\} \). We also show that for monotone \( k \)-submodular functions there is a polynomial-time \( k/(2k - 1) \)-approximation algorithm while for any \( \varepsilon > 0 \) a \( (k + 1)/(2k + \varepsilon) \)-approximation algorithm for maximizing monotone \( k \)-submodular functions would require exponentially many queries. In particular, our hardness result implies that our algorithms are asymptotically tight. We also extend the approach to provide constant factor approximation algorithms for maximizing skew-bisubmodular functions, which were recently introduced as generalizations of bisubmodular functions.

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CP8  
Random-Cluster Model in \( Z^2 \)

The random-cluster model has been widely studied as a unifying framework for random graphs, spin systems and electrical networks, but its dynamics have so far largely resisted analysis. In this paper we analyze the Glauber dynamics of the random-cluster model in the canonical case where the underlying graph is an \( n \times n \) box in the Cartesian lattice \( Z^2 \). Our main result is a \( O(n^2 \log n) \) upper bound for the mixing time at all values of the model parameter \( p \) except the critical point \( p_c(q) \), and for all values of the second model parameter \( q \geq 1 \). We also provide a matching lower bound proving that our result is tight. Our analysis takes as its starting point the recent breakthrough by Beffara and Duminil-Copin on the location of the random-cluster phase transition in \( Z^2 \). It is reminiscent of similar results for spin systems such as the Ising and Potts models, but requires the reworking of several standard tools in the context of the random-cluster model, which is not a spin system in the usual sense.

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CP8  
The Complexity of Approximately Counting in 2-

Spin Systems on \( k \)-Uniform Bounded-Degree Hypergraphs

An important development in approximate counting is the classification of antiferromagnetic 2-spin systems on bounded-degree graphs, based on a beautiful connection to the uniqueness phase transition. For 2-spin models on hypergraphs, the connection between uniqueness and counting breaks down. Nevertheless, we show that for every non-trivial symmetric \( k \)-ary Boolean function \( f \), for \( \Delta \) sufficiently large, it is NP-hard to approximate the partition function of the 2-spin system associated with \( f \) on \( k \)-uniform hypergraphs of max degree \( \Delta \), even within an exponential factor.

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CP8  
Bounds for Random Constraint Satisfaction Problems Via Spatial Coupling

We report on a novel technique called spatial coupling and its application in the analysis of random constraint satisfaction problems (CSP). Spatial coupling was invented as an engineering construction in the area of error correcting codes where it has resulted in efficient capacity-achieving codes for a wide range of channels. However, this technique is not limited to problems in communications, and can be applied in the much broader context of graphical models. We describe here a general methodology for applying spatial coupling to random constraint satisfaction problems and obtain lower bounds for their rough satisfiability threshold, defined as the largest constraint density for which all but a vanishing fraction of the constraints can be satisfied. The main idea is to construct a distribution of geometrically structured random \( k \)-SAT instances - namely the spatially coupled ensemble - which has the same rough satisfiability threshold, and is at the same time algorithmically easier to solve. Then by running well-known algorithms on the spatially coupled ensemble we obtain a lower bound on the rough satisfiability threshold of the original ensemble. The method is versatile because one can choose the CSP, there is a certain amount of freedom in the construction of the bigger ensemble, and also in the choice of the algorithm. In this work we focus on random \( k \)-SAT but we have also checked that the method is successful for graph coloring, NAE-SAT and XOR-SAT. We choose Unit Clause propagation for the algorithm which is analyzed over the spatially coupled instances. For \( k = 3 \), for instance, our lower bound is equal to 3.67 which is better than the current bounds in the literature. Similarly, for graph 3-colorability we get a bound of 4.44 which is also better than the current bounds in the literature.

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CP8  
Canonical Paths for Mcmc: from Art to Science

Markov Chain Monte Carlo (MCMC) method is a widely
CP8
Evolutionary Dynamics in Finite Populations Mix Rapidly

In this paper we prove that the mixing time of a broad class of evolutionary dynamics in finite, unstructured populations is roughly logarithmic in the size of the state space. An important special case of such a stochastic process is the Wright-Fisher model from evolutionary biology (with selection and mutation) on a population of size $N$ over $m$ genotypes. Our main result implies that the mixing time of this process is $O(\log N)$ for all mutation rates and fitness landscapes, and solves the main open problem from [N. Dixit, P. Srivastava and N. K. Vishnoi, A finite population model of molecular evolution: Theory and computation]. In particular, it significantly extends the main result in [N. K. Vishnoi, The speed of evolution] who proved this for $m = 2$. Biologically, such models have been used to study the evolution of viral populations with applications to drug design strategies countering them. Here the time it takes for the population to reach a steady state is important both for the estimation of the steady-state structure of the population as well as in the modeling of the treatment strength and duration. Our result, that such populations exhibit rapid mixing, makes both of these approaches sound. Technically, we make a novel connection between Markov chains arising in evolutionary dynamics and dynamical systems on the probability simplex. This allows us to use the local and global stability properties of the fixed points of such dynamical systems to construct a contractive coupling in a fairly general setting. We expect that our mixing time result would be useful beyond the evolutionary biology setting, and the techniques used here would find applications in bounding the mixing times of Markov chains which have a natural underlying dynamical system.

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CP9
Sparse Approximation Via Generating Point Sets

For a set $P$ of $n$ points in the unit ball $b \subseteq \mathbb{R}^d$, consider the problem of finding a small subset $T \subseteq P$ such that its convex hull $\varepsilon$-approximates the convex hull of the original set. Specifically, the Hausdorff distance between the convex hull of $T$ and the convex hull of $P$ should be at most $\varepsilon$. We present an efficient algorithm to compute such an $\varepsilon'$-approximation of size $k_{alg}$, where $\varepsilon'$ is a function of $\varepsilon$, and $k_{alg}$ is a function of the minimum size $k_{opt}$ of such an $\varepsilon$-approximation. Surprisingly, there is no dependence on the dimension $d$ in either of the bounds. Furthermore, every point of $P$ can be $\varepsilon$-approximated by a convex-combination of points of $T$ that is $O(1/\varepsilon^2)$-sparse. Our result can be viewed as a method for sparse, convex autoencoding: approximately representing the data in a compact way using sparse combinations of a small subset $T$ of the original data. The new algorithm can be kernelized, and it preserves sparsity in the original input.

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CP9
Towards Optimal Algorithms for Prediction with Expert Advice

We study the classical problem of prediction with expert advice in the adversarial setting with a geometric stopping time. In 1965, Cover gave the optimal algorithm for the case of 2 experts. In this paper, we design the optimal algorithm, adversary and regret for the case of 3 experts. Further, we show that the optimal algorithm for 2 and 3 experts is a probability matching algorithm (analogous to Thompson sampling) against a particular randomized adversary.

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CP9
Non-Convex Compressed Sensing with the Sum-of-Squares Method

In this paper, we consider stable signal recovery in \( \ell_q \) quasi-norm for \( 0 < q \leq 1 \). In this problem, given a measurement vector \( \mathbf{y} = A\mathbf{x} \) for some unknown signal vector \( \mathbf{x} \in \mathbb{R}^n \) and a known matrix \( A \in \mathbb{R}^{m \times n} \), we want to recover \( \mathbf{z} \in \mathbb{R}^n \) with \( \|\mathbf{x} - \mathbf{z}\| = O(\|\mathbf{x} - \mathbf{z}_1\|) \) from a measurement vector, where \( \mathbf{z}_1 \) is the \( s \)-sparse vector closest to \( \mathbf{x} \) in \( \ell_q \) quasi-norm. Although smaller \( q \) is favorable for measuring the distance to sparse vectors, previous methods for \( q < 1 \) involve \( \ell_q \) quasi-norm minimization and thus are intractable. In this paper, we overcome this issue by using the sum-of-squares method, and give the first polynomial-time stable recovery scheme for a large class of \( A \) in \( \ell_q \) quasi-norm for any fixed constant \( 0 < q \leq 1 \).

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CP9
Jointly Private Convex Programming

We present a general method for approximately solving convex programs defined by private information from agents, when the solution can be naturally partitioned among the agents. This class of problems includes multi-commodity flow problems, general allocation problems, and multi-dimensional knapsack problems, among other examples. The accuracy of our algorithm depends on the number of coupling constraints, which bind multiple agents. On the other hand, our accuracy is nearly independent of the number of variables, and in many cases, actually improves as the number of agents increases. A special case of our result (solving general allocation problems beyond “Cross Substitute” preferences) resolves the main open problem from [Hsu et al. STOC 2014]. We also consider strategic agents who have preferences over their part of the solution. For any convex program in our class that maximizes social welfare, we show how to create an \emph{approximately dominant strategy truthful} mechanism, approximately maximizing welfare. The central idea is to charge agents prices based on the approximately optimal dual variables, which are themselves computed under differential privacy. Our results substantially broaden the class of problems that are known to be solvable under privacy and/or incentive constraints.

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CP9
Weighted SGD for \( \ell_p \) Regression with Randomized Preconditioning

In recent years, stochastic gradient descent (SGD) methods and randomized linear algebra (RLA) algorithms have been applied to many large-scale problems in machine learning and data analysis. We aim to bridge the gap between these two classes of methods in solving constrained overdetermined linear regression problems—e.g., \( \ell_2 \) and \( \ell_1 \) regression problems.

- We propose a hybrid algorithm named pwSGD that uses RLA techniques for preconditioning and constructing an importance sampling distribution, and then performs an SGD-like iterative process with weighted sampling on the preconditioned system.
- By rewriting the \( \ell_p \) regression problem into a stochastic optimization problem, we connect pwSGD to several existing \( \ell_p \) solvers including RLA methods with algorithmic leveraging (RLA for short).
- Particularly, when solving \( \ell_1 \) regression with size \( n \) by \( d \), pwSGD returns an approximate solution with \( \epsilon \) relative error on the objective value in \( O(n \cdot \text{nnz}(A) + \text{poly}(d)/\epsilon^2) \) time. This complexity is uniformly better than that of RLA methods in terms of both \( \epsilon \) and \( d \) when the problem is unconstrained.
- For \( \ell_2 \) regression, pwSGD returns an approximate solution with \( \epsilon \) relative error on the objective value and solution vector in prediction norm in \( O(n \cdot \text{nnz}(A) + \text{poly}(d) \log(1/\epsilon)/\epsilon) \) time. We show that when solving unconstrained \( \ell_2 \) regression, this complexity is comparable to that of RLA and is asymptotically better over several state-of-the-art solvers in the regime where the desired accuracy \( \epsilon \), high dimension \( n \) and low dimension \( d \) satisfy \( d \geq 1/\epsilon \) and \( n \geq d^2/\epsilon \).

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CP10
Time Vs. Information Tradeoffs For Leader Election in Anonymous Trees

In anonymous leader election, every node of the network must output a simple path so that all these paths end at a common node: the leader. We establish tradeoffs between the number of communication rounds and the amount of...
information that has to be given a priori to the nodes such that they can deterministically solve leader election in anonymous trees. This is achieved using the algorithms with advice framework.

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CP10
Balanced Allocation: Patience is not a Virtue

We present a load-balancing algorithm, FirstDiff[d], that combines the simplicity of Azar et al.’s Greedy[d] and the improved balance of Vöcking’s Left[d]. Using the standard balls-and-bins framework, we show that FirstDiff[d] requires d probes on average per ball and closely matches the maximum load ensured by Left[d] in both the standard and heavily-loaded settings. Unlike Left[d], FirstDiff[d] uses no additional structure on the bins. Experimentally, we show that FirstDiff[d] equals Left[d] in practice.

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CP10
Stabilizing Consensus with Many Opinions

We consider the following distributed consensus problem: Each node in a complete communication network of size n initially holds an opinion, which is chosen arbitrarily from a finite set Σ. The system must converge toward a consensus state in which all, or almost all nodes, hold the same opinion. Moreover, this opinion should be valid, i.e., it should be one among those initially present in the system. This condition should be met even in the presence of a malicious adversary who can modify the opinions of a bounded subset of nodes, adaptively chosen in every round. We consider the 3-majority dynamics: At every round, every node pulls the opinion from three random neighbors and sets his new opinion to the majority one (ties are broken arbitrarily). Let k be the number of valid opinions. We show that, if k ≤ n α, where α is a suitable positive constant, the 3-majority dynamics converges in time polynomial in k and log n with high probability even in the presence of an adversary who can affect up to ϵ(n) nodes at each round. Previously, the convergence of the 3-majority protocol was known for |Σ| = 2 only, with an argument that is robust to adversarial errors. On the other hand, no anonymous, uniform-gossip protocol that is robust to adversarial errors was known for |Σ| > 2.

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CP10
Markovian Hitters and the Complexity of Blind Rendezvous

We define and construct a novel pseudorandom tool, the Markovian hitter. Given an input sequence of n independent random bits, a Markovian hitter produces a sequence of pseudorandom samples in {0, 1} k, in an online fashion, that hits any subset W ⊂ {0, 1} k of size ϵ2 k with probability ≈ 1 − 2 −(n−k) ϵ. This is comparable to the behavior of truly random samples or classical pseudorandom hitting sets. A Markovian hitter has an additional “Markovian” property of interest: each pseudorandom sample is a function of only the most recent bits of the input sequence (of random bits). In particular, we apply Markovian hitters to obtain a new algorithm for the well-studied blind rendezvous problem for cognitive radios. This is the problem faced by two parties equipped with radios that can access different channels. Their challenge is to discover each other (by tuning their radios to the same channel at the same time) as quickly as possible. In prior work [?] it was shown that deterministic schedules have a lower bound for rendezvous time of Ω(|S1| · |S2|). We beat this quadratic barrier by utilizing a public source of randomness in conjunction with a Markovian hitter to achieve rendezvous in expected time

O \left( \log n + \frac{|S_1 \cup S_2|}{|S_1 \cap S_2|} \right).

We counterbalance this result by establishing two lower bounds on expected rendezvous time: an bound for the setting with public randomness, and an Ω(|S1| · |S2|) bound in the setting with private randomness but no public randomness, which is a strengthening of the result for deterministic schedules.

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CP10
How to Scale Exponential Backoff: Constant Throughput, Polylog Access Attempts, and Robustness

Randomized exponential backoff is a widely deployed technique for coordinating access to a shared resource. A good backoff protocol should, arguably, satisfy three natural properties: (i) it should provide constant throughput, wasting as little time as possible; (ii) it should require few failed access attempts, minimizing the amount of wasted effort; and (iii) it should be robust, continuing to work efficiently even if some of the access attempts fail for spurious reasons. Unfortunately, exponential backoff has some well-known limitations in two of these areas: it provides poor (sub-constant) throughput (in the worst case), and is not robust (to adversarial disruption). The goal of this paper is to “fix” exponential backoff by making it scalable, particularly focusing on the case where processes arrive in an on-line, worst-case fashion. We present a relatively simple backoff protocol, Re-Backoff, that has, at its heart, a version of exponential backoff. It guarantees expected constant throughput with dynamic process arrivals and requires only an expected polylogarithmic number of access attempts per process. Re-Backoff is also robust to periods where the shared resource is unavailable for a period of time. If it is unavailable for $D$ time slots, Re-Backoff provides the following guarantees. When the number of packets is a finite $n$, the average expected number of access attempts for successfully sending a packet is $O(\log^2(n + D))$. In the infinite case, the average expected number of access attempts for successfully sending a packet is $O(\log^2(\eta + D))$ where $\eta$ is the maximum number of processes that are ever in the system concurrently.

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CP11
On Dynamic Approximate Shortest Paths for Planar Graphs with Worst-Case Costs

Given a base weighted planar graph $G_{\text{input}}$ on $n$ nodes and parameters $M, \epsilon$ we present a dynamic distance oracle with $1 + \epsilon$ stretch and worst case update and query costs of $\epsilon^{-3}M^{4}\cdot\text{poly-log}(n)$. We allow arbitrary edge weight updates as long as the shortest path metric induced by the updated graph has stretch of at most $M$ relative to the shortest path metric of the base graph $G_{\text{input}}$. For example, on a planar road network, we can support fast queries and dynamic traffic updates as long as the shortest path from any source to any target (including using arbitrary detours) is between, say, 80 and 3 miles-per-hour. As a warm-up we also prove that graphs of bounded treewidth have exact distance oracles in the dynamic edge model. To the best of our knowledge, this is the first dynamic distance oracle for a non-trivial family of dynamic changes to planar graphs with worst case costs of $O(n^{1/2})$ both for query and for update operations.

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CP11
Faster Fully Dynamic Matchings with Small Approximation Ratios

Maximum cardinality matching is a fundamental algorithmic problem with many algorithms and applications. The fully dynamic version, in which edges are inserted and deleted over time has also been the subject of much attention. Existing algorithms for dynamic matching (in general $n$-vertex $m$-edge graphs) fall into two groups: there are fast (mostly randomized) algorithms that achieve a 2-approximation or worse, and there are slow algorithms with $\Omega(\sqrt{m})$ update time that achieve a better-than-2 approximation. Thus the obvious question is whether we can design an algorithm that achieves a tradeoff between these two: a $o(\sqrt{m})$ update time and a better-than-2 approximation simultaneously. We answer this question in the affirmative. Previously, such bounds were only known for the special case of bipartite graphs. Our main result is a fully dynamic deterministic algorithm that maintains a $(3/2 + \epsilon)$-approximation in amortized update time $O(m^{1/4} \epsilon^{-2.5})$. In addition to achieving the tradeoff described above, our algorithm manages to be polynomially faster than all existing deterministic algorithms...
(excluding an existing log \(n\)-approximation of Onak and Rubinfeld), while still maintaining a better-than-2 approximation.

We also give stronger results for graphs whose arboricity is at most \(\alpha\). We show how to maintain a \((1+\varepsilon)\)-approximate fractional matching or a \((3/2+\varepsilon)\)-approximate integral matching in worst-case time \(O(\alpha(\alpha+\log n))\) for constant \(\varepsilon\). When the arboricity is constant, this bound is \(O(\log n)\) and when the arboricity is polylogarithmic the update time is also polylogarithmic. Previous results for small arboricity non-bipartite graphs could only maintain a maximal matching (2-approximation).

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CP11  
Weighted Dynamic Finger in Binary Search Trees

It is shown that the online binary search tree data structure GreedyASS performs asymptotically as well as on a sufficiently long sequence of searches as any static binary search tree where each search begins from the previous search (rather than the root). This bound is known to be equivalent to assigning each item \(i\) in the search tree a positive weight \(w_i\) and bounding the search cost of an item in the search sequence \(s_1, \ldots, s_m\) by

\[
O\left(1 + \log \frac{\min(s_{i-1}, s_i) \leq z \leq \max(s_{i-1}, s_i)}{\min(w_{s_{i-1}}, w_{s_{i-1}})}\right) \text{ amortized.}
\]

This result is the strongest finger-type bound to be proven for binary search trees. By setting the weights to be equal, one observes that our bound implies the dynamic finger bound. Compared to the previous proof of the dynamic finger bound for Splay trees, our result is significantly shorter, stronger, simpler, and has reasonable constants.

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CP11  
Dynamic Dfs in Undirected Graphs: Breaking the \(O(m)\) Barrier

Given an undirected graph \(G = (V,E)\) on \(n\) vertices and \(m\) edges, we address the problem of maintaining a DFS tree when the graph is undergoing updates (insertion and deletion of vertices or edges). We present the following results for this problem.

1. **Fault tolerant DFS tree**: There exists a data structure of size \(O(m)\) such that given any set \(F\) of failed vertices or edges, a DFS tree of the graph \(G \setminus F\) can be reported in \(O(n|F|)\) time.

2. **Fully dynamic DFS tree**: There exists a fully dynamic algorithm for maintaining a DFS tree that takes worst case \(O(\sqrt{mn})\) time per update for any arbitrary online sequence of updates.

3. **Incremental DFS tree**: Given an arbitrary online sequence of edge insertions, we can maintain a DFS tree in \(O(n)\) worst case time per edge insertion.

These are the first \(O(m)\) worst case time results for maintaining a DFS tree in a dynamic environment. Moreover, our fully dynamic algorithm provides, in a seamless manner, the first deterministic algorithm with \(O(1)\) query time and \(O(m)\) worst case update time for the dynamic subgraph connectivity, biconnectivity, and 2-edge connectivity.

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1 The arboricity of a graph is the minimum number of edge-disjoint forests into which it can be partitioned, and it is close to the density of its densest subgraph.
CP12
Clustering Time Series under the Fréchet Distance

The Fréchet distance is a popular distance measure for curves. We study the problem of clustering time series under the Fréchet distance. In particular, we give $(1+\varepsilon)$-approximation algorithms for variations of the following problem with parameters $k$ and $\ell$. Given $n$ univariate time series $P$, each of complexity at most $m$, we find $k$ time series, not necessarily from $P$, which we call cluster centers and which each have complexity at most $\ell$, such that (a) the maximum distance of an element of $P$ to its nearest cluster center or (b) the sum of these distances is minimized. Our algorithms have running time near-linear in the input size for constant $\varepsilon$, $k$ and $\ell$. To the best of our knowledge, our algorithms are the first clustering algorithms for the Fréchet distance which achieve an approximation factor of $(1+\varepsilon)$ or better.

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CP12
Approximating Capacitated K-Median with $(1+\varepsilon)k$
Open Facilities

In this paper, we give the first constant approximation for the capacitated k-median problem, that only violates the cardinality constraint by a factor of $1+\varepsilon$. This generalizes the result of [Li15], which only works for the case when all capacities are the same. Our algorithm is based on a novel configuration LP relaxation for the problem.

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CP12
Simple and Fast Rounding Algorithms for Directed
and Node-Weighted Multiway Cut

In Multiway Cut problem, input is an edge/node-weighted graph and a set of terminals $S = \{s_1,\ldots, s_k\}$, and the goal is to remove minimum weight set of edges to disconnect all terminals. We present a very simple 2-approximation for Directed Multiway Cut and $2(1-1/k)$ approximation for Node-weighted Multiway Cut.

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CP12
Constant Factor Approximation for Subset Feedback Set Problems Via a New LP Relaxation

In the Subset Feedback Edge/Vertex Set problem (SFES/SFVS) the input is edge/node-weighted graph $G = (V,E)$ and a set $S = \{s_1,\ldots, s_k\} \subset V$ of $k$ terminals and the goal is to remove a minimum weight set of edges/nodes such that no interesting containing a terminal remains. One reason for the difficulty in addressing these problems has been the lack of LP relaxations with constant factor integrality gaps. We give first LP relaxation for SFVS and SFES with constant integrality gap.

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CP13
Error Amplification for Pairwise Spanner Lower Bounds

A pairwise spanner of a graph $G = (V,E)$ and a “pair set” $P \subseteq V \times V$ is a subgraph $H$ that preserves all pairwise distances in $P$, up to some additive error term $+\beta$. When $\beta = 0$ the object is called a pairwise distance preserver. A large and growing body of work has considered upper bounds for these objects, but lower bounds have been elusive. The only known lower bound results are (1) Coppersmith and Elkin (SODA’05) against preservers, and (2) considerably weaker bounds by Woodruff (FOCS’06) against spanners. Our main result is an amplification theorem: we prove that lower bounds against pairwise distance preservers imply lower bounds against pairwise spanners. In other words, to prove lower bounds against any constant error spanners, it is enough to consider only subgraphs that are not allowed any error at all! We apply this theorem to obtain drastically improved lower bounds. Some of these include:

- Linear size pairwise spanners with up to $+(2k-1)$ error cannot span $|P| = \omega(n^{(1+k)/(1+k)})$ pairs. This is a large improvement over Woodruff’s $|P| = \omega(n^{2-2/k})$ ($|P|$ is now linear, rather than quadratic, as $k$ gets large).
- $|E(H)| = \Omega(n^{1+1/k})$ edges are required for a $+(2k-1)$ spanner of $|P| = \Omega(n^{1+1/k})$ pairs - another large improvement over Woodruff’s $|P| = \Omega(n^2)$.
- The first tight bounds for pairwise spanners: for $+2$ error and $P = \Theta(n^{3/2})$ we show that $\Theta(n^{3/2})$ edges are necessary and sufficient (this also reflects a new upper bound: we construct $+2$ pairwise spanners on $O(n|P|^{1/3})$ edges, removing a log factor from a prior algorithm).

We also show improved lower bounds against subset spanners (where $P = S \times S$ for some node subset $S$), and lower bounds against $D$ threshold spanners (where $P$ is the set of node pairs at distance at least $D$).

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CP13
Better Distance Preservers and Additive Spanners

We make improvements to the upper bounds on several popular types of distance preserving graph sketches. The first part of our paper concerns pairwise distance preservers, which are sparse subgraphs that exactly preserve the pairwise distances for a set of given pairs of vertices. Our main result here is that all unweighted, undirected n-node graphs G and all pair sets P have distance preservers on |H| = O(n^{2/3}|P|^{2/3} + n|P|^{1/3}) edges. This improves the known bounds whenever |P| = ω(n^{3/4}). We then develop a new graph clustering technique, based on distance preservers, and we apply this technique to show new upper bounds for additive (standard) spanners, in which all pairwise distances must be preserved up to an additive error function, and for subset spanners, in which only distances within a given node subset must be preserved up to an error function. For both of these objects, we obtain the new best tradeoff between spanner sparsity and error allowance in the regime where the error is polynomial in the graph size. We leave open a conjecture that O(n^{2/3}|P|^{2/3} + n) pairwise distance preservers are possible for undirected unweighted graphs. Resolving this conjecture in the affirmative would improve and simplify our upper bounds for all the graph sketches mentioned above.

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CP13
Near-Optimal Light Spanners

In this work we consider spanners of light weight for general undirected graphs. For any n-vertex graph G and any positive integer k, we present a spanner with (2k - 1)·(1 + ε) stretch, O,ω(MST(G))n^{1/k} weight, and O,(n^{1+1/k}) edges. Up to a (1 + ε) factor in the stretch this matches the girth conjecture of Erdős. It improves over the recent work of Elkin, Neiman and Solomon [ICALP 14] by reducing the spanner weight by a factor of k/ log k.

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CP13
On Notions of Distortion and An Almost Minimum Spanning Tree with Constant Average Distortion

Minimum Spanning Trees of weighted graphs are fundamental objects in numerous applications. In particular in distributed networks, the minimum spanning tree of the network is often used to route messages between network nodes. Unfortunately, while being most efficient in the total cost of connecting all nodes, minimum spanning trees fail miserably in the desired property of approximately preserving distances between pairs. While known lower bounds exclude the possibility of the worst case distortion of a tree being small, it was shown in [ABN15] that there exists a spanning tree with constant average distortion. Yet, the weight of such a tree may be significantly larger than that of the MST. In this paper, we show that any weighted undirected graph admits a spanning tree whose weight is at most (1 + ρ) times that of the MST, providing constant average distortion O(1/ρ^2). The constant average distortion bound is implied by a stronger property of scaling distortion, i.e., improved distortion for smaller fractions of the pairs. The result is achieved by first showing the existence of a low weight spanner with small prioritized distortion, a property allowing to prioritize the nodes whose associated distortions will be improved. We show that prioritized distortion is essentially equivalent to coarse scaling distortion via a general transformation, which has further implications and may be of independent interest. In particular, we obtain an embedding for arbitrary metrics into Euclidean space with optimal prioritized distortion.

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CP13
Approximating Low-Stretch Spanners

Despite significant recent progress on approximating graph spanners (subgraphs which approximately preserve distances), there are still several large gaps in our understanding. We give new results for two of them: approximating basic k-spanner (particularly for small k), and the dependence on f when approximating f-fault tolerant spanners. We first design an O(n^{1/k})-approximation for 4-spanner (both basic and directed). This was the last value of k for which only an O(√n)-approximation was known for basic k-spanner, and thus implies that for any k the approximation ratio is at most O(n^{1/3}). For basic k-spanner, we also show an integrality gap for the natural flow-based LP (the main tool in almost all nontrivial spanner approximations) which nearly matches the trivial approximation of n^{(1+1/k)/2}. For f-fault tolerant spanners, we show that in the small-stretch setting (k ∈ {3, 4}) it is possible to entirely remove the dependence on f from the approximation ratio, at the cost of moving to bicriteria guarantees. The previous best dependence on f was either almost-linear (in the undirected setting) or exponential (in the directed setting for stretch 4).

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In the $k$-junta testing problem, a tester has to efficiently decide whether a given function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ is a $k$-junta (i.e., depends on at most $k$ of its input bits) or is $\varepsilon$-far from any $k$-junta. Our main result is a quantum algorithm for this problem with query complexity $\tilde{O}(\sqrt{k/\varepsilon})$ and time complexity $\tilde{O}(n\sqrt{k/\varepsilon})$. This quadratically improves over the query complexity of the previous best quantum junta tester, due to Atıcı and Servedio. Our tester is based on a new quantum algorithm for a gapped version of the combinatorial graph testing problem, with an up to quartic improvement over the query complexity of the best classical algorithm. For our upper bound on the time complexity we give a near-linear time implementation of a shallow variant of the quantum Fourier transform over the symmetric group, similar to the Schur-Weyl transform.
CP15
Multiscale Mapper: Topological Summarization Via Codomain Covers

Summarizing topological information from datasets and maps defined on them is a central theme in topological data analysis. Mapper, a tool for such summarization, takes as input both a possibly high dimensional dataset and a map defined on the data, and produces a summary of the data by using a cover of the codomain of the map. This cover, via a pullback operation to the domain, produces a simplicial complex connecting the data points. The resulting view of the data through a cover of the codomain offers flexibility in analyzing the data. However, it offers only a view at a fixed scale at which the cover is constructed. Inspired by the concept, we explore a notion of a tower of covers which induces a tower of simplicial complexes connected by simplicial maps, which we call multiscale mapper.

We study the resulting structure, and design practical algorithms to compute its persistence diagrams efficiently. We present a combinatorial version of the algorithm that acts on the 1-skeleton of the input complex. For general maps, we show that the algorithm approximates the exact persistence diagram only on vertex sets connected by simplicial maps, which we call multiscale mapper.

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CP15
Beyond the Richter-Thomassen Conjecture

If two closed Jordan curves in the plane have precisely one point in common, then it is called a touching point. All other intersection points are called crossing points. The main result of this paper is a Crossing Lemma for closed curves: In any family of \( n \) pairwise intersecting simple closed curves in the plane, no three of which pass through the same point, the number of crossing points exceeds the number of touching points by a factor of \( \Omega((\log \log n)^{1/8}) \).

As a corollary, we prove the following long-standing conjecture of Richter and Thomassen: The total number of intersection points between any \( n \) pairwise intersecting simple closed curves in the plane, no three of which pass through the same point, is at least \( (1 - o(1))n^2 \).

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CP15
Windrose Planarity: Embedding Graphs with Direction-Constrained Edges

Given a planar graph \( G(V,E) \) and an assignment of the neighbours of each vertex \( v \in V \) to the four quadrants around \( v \), problem Windrose Planarity asks whether \( G \) admits a windrose-planar drawing, i.e., a planar drawing in which edges are \( xy \)-monotone curves and the neighbours of each vertex \( v \) lie in the assigned quadrants. We give a polynomial-time algorithm for Windrose Planarity on embedded graphs, which in the positive case constructs a 1-bend windrose-planar drawing on an \( O(n) \times O(n) \) grid.

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CP16
An Improved Approximation Guarantee for the Maximum Budgeted Allocation Problem

We study the Maximum Budgeted Allocation problem, which is the problem of assigning indivisible items to players with budget constraints. The best approximation algorithms we know for the MBA problem achieve a 3/4-approximation ratio, and employ the natural Assignment-LP. In this paper, we present a 3/4 + \( c \)-approximation algorithm for MBA, for some constant \( c > 0 \). This algorithm rounds solutions to the Configuration-LP, therefore also showing that the Configuration-LP is stronger than the Assignment-LP for the MBA problem.

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CP16

On the Integrality Gap of Degree-4 Sum of Squares for Planted Clique

The problem of finding large cliques in random graphs and its “planted” variant, where one wants to recover a clique of size \( \omega \gg \log (n) \) added to an Erdős-Rényi graph \( G \sim G(n, \frac{1}{2}) \), have been intensely studied. Nevertheless, existing polynomial time algorithms can only recover planted cliques of size \( \omega = \Omega(\sqrt{n}) \). By contrast, information theoretically, one can recover planted cliques so long as \( \omega \gg \log (n) \). In this work, we continue the investigation of algorithms from the sum of squares hierarchy for solving the planted clique problem begun by Meka, Potechin, and Wigderson [STOC 2015] and Deshpande and Montanari [COLT 2015]. Our main results improve upon both these previous works by showing:

1. Degree four SoS does not recover the planted clique unless \( \omega \gg \sqrt{n}/\text{polylog} \; n \), improving upon the bound \( \omega \gg n^{1/4} \) due to Deshpande and Montanari.

2. For \( 2 < d = o(\sqrt{\log (n)}) \), degree \( 2d \) SoS does not recover the planted clique unless \( \omega \gg n^{1/(2d+1)} / (2d^{d}/\text{polylog} \; n) \), improving upon the bound due to Meka, Potechin and Wigderson.

Our proof for the second result is based on a fine spectral analysis of the certificate used in the prior works by decomposing it along an appropriately chosen basis. Along the way, we develop combinatorial tools to analyze the spectrum of random matrices with dependent entries and to understand the symmetries in the eigenspaces of the set symmetric matrices inspired by work of Grigoriev. An argument of Kelner shows that the first result cannot be proved using the same certificate. Rather, our proof involves constructing and analyzing a new certificate that yields the nearly tight lower bound by “correcting” the previous certificate.

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CP16

The Matching Problem Has No Small Symmetric Sdp

Yannakakis showed that the matching problem does not have a small symmetric linear program. Rothvöß recently proved that any, not necessarily symmetric, linear program also has exponential size. It is natural to ask whether the matching problem can be expressed compactly in a framework such as semidefinite programming (SDP) that is more powerful than linear programming but still allows efficient optimization. We answer this question negatively for symmetric SDPs: any symmetric SDP for the matching problem has exponential size. We also show that an \( O(k) \)-round Lasserre SDP relaxation for the metric traveling salesman problem yields at least as good an approximation as any symmetric SDP relaxation of size \( n^k \). The key technical ingredient underlying both these results is an upper bound on the degree needed to derive polynomial identities that hold over the space of matchings or traveling salesperson tours.

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cardinality of the ground set and the time for computing the saturation and exchange capacities. This is as fast as the best known algorithm for the polymatroid intersection problem.

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CP16
Online Contention Resolution Schemes

We introduce a new rounding technique designed for online optimization problems, called online contention resolution schemes (OCRSs). OCRSs are applicable to various online selection problems, including Bayesian online selection, oblivious posted pricing mechanisms, and stochastic probing models. Furthermore, they share many properties of offline contention resolution schemes. In particular, OCRSs for different constraint families can be combined to obtain an OCRS for their intersection. Moreover, we can approximately maximize submodular functions in various online settings.

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CP17
Lower Bounds for the Parameterized Complexity of Minimum Fill-In and Other Completion Problems

In this work, we focus on several completion problems for subclasses of chordal graphs: MINIMUM FILL-IN, INTERVAL COMPLETION, PROPER INTERVAL COMPLETION, THRESHOLD COMPLETION, and TRIVIALLY PERFECT COMPLETION. In these problems, the task is to add at most \( k \) edges to a given graph in order to obtain a chordal, interval, proper interval, threshold, or trivially perfect graph, respectively. We prove the following lower bounds for all these problems, as well as for the related CHAIN COMPLETION problem:

- Assuming the Exponential Time Hypothesis, none of these problems can be solved in time \( 2^{O(n^{1/2} / \log^c n)} \) or \( 2^{O(k^{1/4} / \log^c k)} \cdot n^{O(1)} \), for some integer \( c \).
- Assuming the non-existence of a subexponential-time approximation scheme for MIN BISECTION on \( d \)-regular graphs, for some constant \( d \), none of these problems can be solved in time \( 2^{o(n)} \) or \( 2^{o(\sqrt{k})} \cdot n^{O(1)} \).

The second result is an evidence, that a significant improvement of the best known algorithms for parameterized completion problems would lead to a surprising breakthrough in the design of approximation algorithms for MIN BISECTION.

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CP17
Linear Recognition of Almost Interval Graphs

Let interval + \( kv \), interval + \( ke \), and interval − \( ke \) denote the classes of graphs that can be obtained from some interval graph by adding \( k \) vertices, adding \( k \) edges, and deleting \( k \) edges, respectively. When \( k \) is small, these graph classes are called almost interval graphs. They are well motivated from computational biology, where the data ought to be represented by an interval graph while we can only expect an almost interval graph for the best. For any fixed \( k \), we give linear-time algorithms for recognizing all these classes, and in the case of membership, our algorithms provide also a specific interval graph as evidence. When \( k \) is part of the input, these problems are also known as graph modification problems, all NP-complete. Our results imply that they are fixed-parameter tractable parameterized by \( k \), thereby resolving the long-standing open problem on the parameterized complexity of recognizing interval + \( ke \), first asked by Bodlaender et al. [Bioinformatics, 11:49–57, 1995]. Moreover, our algorithms for recognizing interval + \( kv \) and interval − \( ke \) run in times \( O(6^k \cdot (n + m)) \) and \( O(8^k \cdot (n + m)) \), (where \( n \) and \( m \) stand for the numbers of vertices and edges respectively in the input graph,) significantly improving the \( O(k^{2e} \cdot n^m) \)-time algorithm of Hegedus et al. [STOC 2007; SICOMP 2009] and the \( O(10^k \cdot n^m) \)-time algorithm of Cao and Marx [SODA 2014; TALG 2015] respectively.

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CP17
Raising The Bar For Vertex Cover: Fixed-Parameter Tractability Above A Higher Guarantee

We investigate the following above-guarantee parameterization of the classical Vertex Cover problem: Given a graph \( G \) and an integer \( k \) as input, does \( G \) have a vertex cover of size at most \( (2LP - MM) + k \) ? Here \( MM \) is the size of a maximum matching of \( G \), \( LP \) is the value of an optimal solution to the relaxed (standard) LP for Vertex Cover on \( G \), and \( k \) is the parameter. It can be shown that \( (2LP - MM) \) is a lower bound on vertex cover size, and since \( (2LP - MM) \geq LP \geq MM \), this is a stronger pa-
rameterization than those—namely, above $MM$, and above $LP$—which have been studied so far. We prove that Vertex Cover is fixed-parameter tractable for this stronger parameter $k$: We derive an algorithm which solves Vertex Cover in time $O^*(3^k)$, and thus push the envelope further on the parameterized tractability of Vertex Cover.

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**CP17**

**Directed Multicut Is $W[1]$-Hard, Even for Four Terminal Pairs**

We prove that Multicut in directed graphs, parameterized by the size of the cutset, is $W[1]$-hard and hence unlikely to be fixed-parameter tractable even if restricted to instances with only four terminal pairs. This negative result almost completely resolves one of the central open problems in the area of parameterized complexity of graph separation problems, posted originally by Marx and Razgon [SIAM J. Comput. 43(2):355–388 (2014)], leaving only the case of three terminal pairs open. The case of two terminal pairs was shown to be FPT by Chitnis et al. [SIAM J. Comput. 42(4):1674–1696 (2013)]. Our gadget methodology also allows us to prove $W[1]$-hardness of the Steiner Orientation problem parameterized by the number of terminal pairs, resolving an open problem of Cygan, Kortsarz, and Nutov [SIAM J. Discrete Math. 27(3):1503-1513 (2013)].

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**CP17**  

**Subexponential Parameterized Algorithm for Interval Completion**

In the Interval Completion problem we are given an $n$-vertex graph $G$ and an integer $k$, and the task is to transform $G$ by making use of at most $k$ edge additions into an interval graph. This is a fundamental graph modification problem with applications in sparse matrix multiplication and molecular biology. The question about fixed-parameter tractability of Interval Completion was asked by Kaplan, Shamir and Tarjan [FOCS 1994; SIAM J. Comput. 1999] and was answered affirmatively more than a decade later by Villanger at el. [STOC 2007; SIAM J. Comput. 2009], who presented an algorithm with running time $O(k^{2k}n^3m)$. We give the first subexponential parameterized algorithm solving Interval Completion in time $k^{O(k)}n^{O(1)}$. This adds Interval Completion to a very small list of parameterized graph modification problems solvable in subexponential time.

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**CP18**  

**Improved Deterministic Algorithms for Linear Programming in Low Dimensions**

At SODA’93, Chazelle and Matoušek presented a derandomization of Clarkson’s sampling-based algorithm [FOCS’88] for solving linear programs with $n$ constraints and $d$ variables in $d^{(7+o(1))n}$ deterministic time. The time bound can be improved to $d^{(3+o(1))n}$ with subsequent work by Brönnimann, Chazelle, and Matoušek [FOCS’93]. We first point out a much simpler derandomization of Clarkson’s algorithm that avoids $\varepsilon$-approximations and runs in $d^{(3+o(1))n}$ time. We then describe a few additional ideas that eventually improve the deterministic time bound to $d^{(1/2+o(1))n}$.

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**CP18**  

**An Efficient Algorithm for Computing High-Quality Paths Amid Polygonal Obstacles**

We study a path-planning problem amid a set of obstacles in $\mathbb{R}^2$. Specifically, the problem asks for a path minimizing the reciprocal of the clearance integrated over the length of the path. We present the first polynomial-time approximation scheme for this problem. Let $n$ be the total number of obstacle vertices and let $\varepsilon \in (0,1]$. Our algorithm computes in time $O(\frac{n^2}{\varepsilon^2} \log \frac{1}{\varepsilon})$ a path of total cost at most $(1+\varepsilon)$ times the cost of the optimal path.

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**CP18**  

**Approximating the K-Level in Three-Dimensional Plane Arrangements**

For a set of $n$ non-vertical planes in three dimensions, and a parameter $r < n$, we give a simple alternative proof of the existence of $O(1/r)$-cutting of the first $n/r$ levels of $A(H)$ via a terrain consisting of $O(r/e^2)$ triangular faces. The proof avoids sampling, and exploits techniques based on planar separators and various structural properties of levels in three-dimensional arrangements and planar maps.

Sariel Har-Peled
Smolensky

We show how to solve all-pairs shortest paths on $n$ nodes in deterministic $n^3/2^{O((\sqrt{\log n})}$ time, and how to deterministically count the pairs of orthogonal vectors among $n$-1 vectors in $d = c \log n$ dimensions in $n^{d^2+1}/O(d^2 \log c)$ time. These running times essentially match the best known randomized algorithms of (Williams, STOC’14) and (Abboud, Williams, and Yu, SODA 2015) respectively, and the ability to count was open even for randomized algorithms. By reductions, these two results yield faster deterministic algorithms for many other problems. Our techniques can also be used to deterministically count $k$-SAT assignments on $n$ variable formulas in $2^{n - n/O(k)}$ time, roughly matching the best known running times for detecting satisfiability and resolving an open problem of Santhanam (2013). A key to our constructions is an efficient way to deterministically simulate certain probabilistic polynomials critical to the algorithms of prior work, carefully applying small-biased sets and modulus-amplifying polynomials.

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CP19
Higher Lower Bounds from the 3sum Conjecture

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CP19
Algorithmic Complexity of Power Law Networks

We define a deterministic condition for checking whether a graph has a power law degree distribution and experimentally validate it on real-world networks. This definition allows us to derive interesting properties of power law networks. We observe that for exponents of the degree distribution in the range \([1, 2]\) such networks exhibit double power law phenomenon that was observed for several real-world networks. Moreover, we give a novel theoretical explanation why many algorithms run faster on real-world data than what is predicted by algorithmic worst-case analysis. We show how to exploit the power law degree distribution to design faster algorithms for a number of classic problems including transitive closure, maximum matching, determinant, PageRank, matrix inverse, counting triangles and finding maximum clique. In contrast to previously done average-case analyses, we believe that this is the first “waterproof” argument that explains why many real-world networks are easier.

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CP19
Subtree Isomorphism Revisited

The Subtree Isomorphism problem asks whether a given tree is contained in another given tree. The problem is of fundamental importance and has been studied since the 1960s. For some variants, e.g., ordered trees, near-linear time algorithms are known, but for the general case truly subquadratic algorithms remain elusive. Our first result is a reduction from the Orthogonal Vectors problem to Subtree Isomorphism, showing that a truly subquadratic algorithm for the latter refutes the Strong Exponential Time Hypothesis (SETH). In light of this conditional lower bound, we focus on natural special cases for which no truly subquadratic algorithms are known. We classify these cases against the quadratic barrier, showing in particular that:

- Even for binary, rooted trees, a truly subquadratic algorithm refutes SETH.
- Even for rooted trees of depth \(O(\log \log n)\), where \(n\) is the total number of vertices, a truly subquadratic algorithm refutes SETH.
- For every constant \(d\), there is a constant \(\epsilon_d > 0\) and a randomized, truly subquadratic algorithm for degree-\(d\) rooted trees of depth at most \((1 + \epsilon_d) \log_d n\). In particular, there is an \(O(\min\{2.85^n, n^2\})\) algorithm for binary trees of depth \(h\).

Our reductions utilize new “tree gadgets” that are likely useful for future SETH-based lower bounds for problems on trees. Our upper bounds apply a folklore result from randomized decision tree complexity.

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CP20
Maximum Matchings in Dynamic Graph Streams and the Simultaneous Communication Model

We study the problem of finding an approximate maximum matching in two closely related computational models, namely, the dynamic graph streaming model and the simultaneous multi-party communication model. We resolve the space complexity of single-pass turnstile streaming algorithms for approximating matchings. Our results for dynamic graph streams also resolve the simultaneous communication complexity of approximating matchings in the edge partition model. We further design new random-
ized and deterministic protocols for the vertex partition model.

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CP20
Incidence Geometries and the Pass Complexity Of Semi-Streaming Set Cover

Set cover, over a universe of size $n$, may be modelled as a data-streaming problem, where the $m$ sets that comprise the instance are to be read one by one. A semi-streaming algorithm is allowed only $O(p \log n, \log m)$ space to process this stream. For each $p \geq 1$, we give a very simple deterministic algorithm that makes $p$ passes over the input stream and returns an appropriately certified $(p + 1)n^{1/(p+1)}$ approximation to the optimum set cover. More importantly, we proceed to show that this approximation factor is essentially tight, by showing that a factor better than $0.99 n^{1/(p+1)}/(p + 1)^2$ is unachievable for a $p$-pass semi-streaming algorithm, even allowing randomisation. In particular, this implies that achieving a $\Theta(\log n)$-approximation requires $\Omega(\log n/\log \log n)$ passes, which is tight up to the $\log \log n$ factor. These results extend to a relaxation of the set cover problem where we are allowed to leave an $\epsilon$ fraction of the universe uncovered: the tight bounds on the best approximation factor achievable in $p$ passes turn out to be $\Theta_p(\min\{n^{1/(p+1)}, \epsilon^{-1/p}\})$. Our lower bounds are based on a construction of a family of high-rank incidence geometries, which may be thought of as vast generalisations of affine planes. This construction, based on algebraic techniques, appears flexible enough to find other applications and is therefore interesting in its own right.

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CP20
Clustering Problems on Sliding Windows

We present the first polylogarithmic space $O(1)$-approximation to the metric $k$-median and metric $k$-means problems in the sliding window model. We show that using only polylogarithmic space we can maintain a summary of the current window from which we can construct an $O(1)$-approximate clustering solution. Additionally, we give the first algorithm that, given an insertion-only streaming coreset of space $s$ (using merge-and-reduce), maintains a coreset in the sliding window model using $O(s^2 \epsilon^{-2} \log W)$ space.

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CP20
Kernelization via Sampling with Applications to Finding Matchings and Related Problems in Dynamic Graph Streams

We present a simple but powerful subgraph sampling primitive that is applicable in a variety of computational models including dynamic graph streams (where the input graph is defined by a sequence of edge/hyperedge insertions and deletions) and distributed systems such as MapReduce. In the case of dynamic graph streams, we use this primitive to prove the following results:

- **Matching**: Our main result for matchings is that there exists an $\tilde{O}(k^2)$ space algorithm that returns the edges of a maximum matching on the assumption the cardinality is at most $k$. The best previous algorithm used $\tilde{O}(kn)$ space where $n$ is the number of vertices in the graph and we prove our result is optimal up to logarithmic factors. Our algorithm has $\tilde{O}(1)$ update time. We also show that there exists an $\tilde{O}(n^2/\alpha^3)$ space algorithm that returns an $\alpha$-approximation for matchings of arbitrary size. In independent work, Assadi et al. (SODA 2016) proved this approximation algorithm is optimal and provided an alternative algorithm. We generalize our exact and algorithms to weighted matching. For graphs with low arboricity such as planar graphs, the space required for constant approximation can be further reduced.

- **Vertex Cover and Hitting Set**: There exists an $\tilde{O}(d^k)$ space algorithm that solves the minimum hitting set problem where $d$ is the cardinality of the input sets and $k$ is an upper bound on the size of the minimum hitting set. We prove this is optimal up to logarithmic factors. Our algorithm has $\tilde{O}(1)$ update time.

Finally, we present fast, small-space dynamic graph stream algorithms for a larger family of parameterized problems and show lower bounds for natural problems outside this family.

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CP21

On the Complexity of Dynamic Mechanism Design

We introduce a dynamic mechanism design problem in which the designer wants to offer two items for sale to the same agent, one now and one in the future. We show that finding the revenue maximizing deterministic mechanism, subject to ex-post individual rationality and truthfulness, is NP-hard. We also prove several positive results and we show that, in an environment in which contracts cannot be enforced, the optimum mechanism requires multiple rounds of communication.

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CP21

Simple Pricing Schemes for Consumers With Evolving Values

We consider a pricing problem where a buyer is interested in purchasing/using a good, such as an app or music or software, repeatedly over time. The consumer discovers his value for the good only as he uses it, and the value evolves with each use as a martingale. We provide a simple pricing scheme and show that its revenue is a constant fraction of the buyers expected cumulative value.

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CP21

On the Economic Efficiency of the Combinatorial Clock Auction

Since the 1990s spectrum auctions have been implemented world-wide. This has provided for a practical examination of an assortment of auction mechanisms and, amongst these, two simultaneous ascending price auctions have proved to be extremely successful. These are the simultaneous multiround ascending auction (SMRA) and the combinatorial clock auction (CCA). It has long been known that, for certain classes of valuation functions, the SMRA provides good theoretical guarantees on social welfare. However, no such guarantees were known for the CCA. In this paper, we show that CCA does provide strong guarantees on social welfare provided the price increment and stopping rule are well-chosen. This is very surprising in that the choice of price increment has been used primarily to adjust auction duration and the stopping rule has attracted little attention. The main result is a polylogarithmic approximation guarantee for social welfare when the maximum number of items demanded $C$ by a bidder is fixed. Specifically, we show that either the revenue of the CCA is at least an $\Omega\left(\frac{\log n \log \log m}{\log n}\right)$-fraction of the optimal welfare or the welfare of the CCA is at least an $\Omega\left(\frac{\log n}{\log n}\right)$-fraction of the optimal welfare, where $n$ is the number of bidders and $m$ is the number of items. As a corollary, the welfare ratio – the worst case ratio between the social welfare of the optimum allocation and the social welfare of the CCA allocation – is at most $O(C^2 \cdot \log n \cdot \log^2 m)$. We emphasize that this latter result requires no assumption on bidders valuation functions. Finally, we prove that such a dependence on $C$ is necessary. In particular, we show that the welfare ratio of the CCA is at least $\Omega\left(\frac{C \cdot \log m}{\log \log m}\right)$.
Strong Budget Balance

This paper focuses on two-sided algorithmic mechanism design - a research area that has not been thoroughly studied as its one-sided counterpart. In the light of [Myerson and Satterwhite, 1983]’s impossibility result, which prevents any individually rational, dominant strategy incentive compatible, and strongly budget-balanced mechanism maximizing the social welfare, we devise a two-sided version of sequential posted price mechanisms to obtain a constant approximation to the social welfare while satisfying the above constraints.

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CP21
Interpolating Between Truthful and Non-Truthful Mechanisms for Combinatorial Auctions

We study the communication complexity of combinatorial auctions via interpolation mechanisms that interpolate between non-truthful and truthful protocols. Specifically, an interpolation mechanism has two phases. In the first phase, the bidders participate in some non-truthful protocol whose output is itself a truthful protocol. In the second phase, the bidders participate in the truthful protocol selected during phase one. Note that virtually all existing auctions have either a non-existent first phase (and are therefore truthful mechanisms), or a non-existent second phase (and are therefore just traditional protocols, analyzed via the Price of Anarchy/Stability). The goal of this paper is to understand the benefits of interpolation mechanisms versus truthful mechanisms or traditional protocols, and develop the necessary tools to formally study them. Interestingly, we exhibit settings where interpolation mechanisms greatly outperform the optimal traditional and truthful protocols. Yet, we also exhibit settings where interpolation mechanisms are provably no better than truthful ones. Finally, we apply our new machinery to prove that the recent single-bid mechanism of Devanur et. al. (the only pre-existing interpolation mechanism in the literature) achieves the optimal price of anarchy among a wide class of protocols, a claim that simply can’t be addressed by appealing just to machinery from communication complexity or the study of truthful mechanisms.

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ρN, the neighbor set \(\Gamma(S)\) contains at least \((1 - \delta)M\) distinct vertices. Our main results are strong integrality gaps in the Lasserre hierarchy and an approximation algorithm for dispersers.

1. For any \(\alpha > 0\), \(\delta > 0\), and a random bipartite graph \(G\) with left degree \(D = O(\log N)\), we prove that the Lasserre hierarchy cannot distinguish whether \(G\) is an \((N^\alpha, (1 - \delta)M)\)-dispenser or not an \((N^{1-\alpha}, \delta M)\)-dispenser.

2. For any \(\rho > 0\), we prove that there exist infinitely many constants \(d\) such that the Lasserre hierarchy cannot distinguish whether a random bipartite graph \(G\) with right degree \(d\) is a \((\rho N, (1 - (1 - \rho)^d)M)\)-dispenser or not a \((\rho N, (1 - \Omega(\frac{1}{d(1 - \rho)})M)\)-dispenser.

Our method gives an integrality gap in the Lasserre hierarchy for bipartite expanders with left degree \(D\). \(G\) on \([N],[M]\) is a \((\rho N, a)\)-expander if for any subset \(S \subseteq [N]\) of size \(\rho N\), the neighbor set \(\Gamma(S)\) contains at least \(a \cdot \rho N\) distinct vertices. We prove that for any constant \(\epsilon > 0\), there exist constants \(\epsilon' < \epsilon, \rho\), and \(D\) such that the Lasserre hierarchy cannot distinguish whether a bipartite graph on \([N],[M]\) with left degree \(D\) is a \((\rho N, (1 - \epsilon')D)\)-expander or not a \((\rho N, (1 - \epsilon)D)\)-expander.

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CP22
Packing Small Vectors

We study the online \(d\)-dimensional Vector Bin Packing problem for small vectors (relative to the size of a bin). We give a constant competitive ratio, improving the best known algorithm from \(O(\log d)\) to \(\approx \epsilon\), for arbitrary \(d\). For two dimensions, we present a First Fit variant with a competitive ratio \(\approx 1.48\). Additionally, we give an essentially tight algorithm (not via a First Fit variant) with a competitive ratio arbitrarily close to \(4/3\).

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CP22
Improved Approximation for Vector Bin Packing

We study the \(d\)-dimensional vector bin packing problem, a well-studied generalization of bin packing arising in resource allocation and scheduling problems. Here we are given a set of \(d\)-dimensional vectors \(v_1, \ldots, v_n\) in \([0,1]^d\), and the goal is to pack them into the least number of bins so that for each bin \(B\), the sum of the vectors in it is at
most 1 in every dimension, i.e., $\| \sum_{v_i \in B} v_i \|_\infty \leq 1$. For the 2-dimensional case we give an asymptotic approximation guarantee of $1 + \ln(1.5) + \epsilon \approx (1.405 + \epsilon)$, improving upon the previous bound of $1 + \ln 2 + \epsilon \approx (1.693 + \epsilon)$. We also give an almost tight $(1.5 + \epsilon)$ absolute approximation guarantee, improving upon the previous bound of 2. For the $d$-dimensional case, we get a $1.5 + \ln(\frac{d+1}{\epsilon}) + \epsilon \approx 0.807 + \ln(d+1) + \epsilon$ guarantee, improving upon the previous $(1 + \ln d + \epsilon)$ guarantee based on the Round & Approx framework. Here $(1 + \ln d)$ was a natural barrier as rounding-based algorithms can not achieve better than $d$ approximation. We get around this by exploiting various structural properties of (near)-optimal packings, and using multi-objective multi-budget matching based techniques and expanding the R & A framework to go beyond rounding-based algorithms. Along the way we also prove several results that could be of independent interest.

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CP22
Approximation Schemes for Machine Scheduling with Resource (in-)dependent Processing Times

We consider two related scheduling problems: resource constrained scheduling on identical parallel machines and a generalization with resource dependent processing times. In both problems, jobs require a certain amount of an additional resource and have to be scheduled minimizing the makespan, while at every point in time a given resource capacity is not exceeded. We present a method to obtain asymptotic fully polynomial approximation schemes (AFPTAS) for the problems.

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CP22
On Approximating Strip Packing with a Better Ratio Than 3/2

In the strip packing problem we are given a set of rectangular items that we want to place in a strip of given width such that we minimize the height of the obtained packing. It is a very classical two-dimensional packing problem that has received a lot of attention and it has applications in many settings such as stock-cutting and scheduling. A straight-forward reduction from Partition shows that the problem cannot be approximated with a better absolute factor than 3/2. However, this reduction requires the numeric values to be exponentially large. In this paper, we present a $(1.4 + \epsilon)$-approximation algorithm with pseudo-polynomial running time. This implies that for polynomially bounded input data the problem can be approximated with a strictly better ratio than for exponential input which is a very rare phenomenon in combinatorial optimization.

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CP23
Recovery and Rigidity in a Regular Stochastic Block Model

The stochastic block model is a natural model for studying community detection in random networks. Its clustering properties have been extensively studied in the statistics, physics and computer science literature. Recently this area has experienced major mathematical breakthroughs, particularly for the binary (two-community) version, see [Mossel, Neeman, Sly (2012, 2013)] and [Massoulie (2013)]. In this paper, we introduce a variant of the binary model which we call the regular stochastic block model (RSBM). We prove rigidity by showing that with high probability an exact recovery of the community structure is possible. Spectral methods exhibit a regime where this can be done efficiently. Moreover we also prove that, in this setting, any suitably good partial recovery can be bootstrapped to obtain a full recovery of the communities.

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CP23
Connectivity in Bridge-Addable Graph Classes: the McDiarmid-Steger-Welsh Conjecture

We prove a conjecture of McDiarmid, Steger and Welsh that says that if a class of graphs is bridge-addable (i.e. stable by adding edges between different components) then the random graph from this class is connected with proba-
ability at least $e^{-1/2} + o(1)$, when its number of vertices tends to infinity. This lower bound is tight since it is reached for forests. The best previously known constants where $e^{-1}$, $e^{-0.7963}$ and $e^{-2/3}$ proved respectively by McDiarmid, Steger and Welsh, by Balister, Bollobás and Gerke, and by Norin.

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CP23
The Power of Two Choices with Simple Tabulation

The power of two choices is a classic paradigm for load balancing when assigning $m$ balls to $n$ bins. When placing a ball, we pick two bins according to two hash functions $h_0$ and $h_1$, and place the ball in the least loaded bin. Assuming fully random hash functions, when $m = O(n)$, Azar et al. [STOC’94] proved that the maximum load is $\lg \lg n + O(1)$ with high probability. No such bound was known with a hash function implementable in constant time. In this paper, we investigate the power of two choices when the hash functions $h_0$ and $h_1$ are implemented with simple tabulation, which is a very efficient hash function evaluated in constant time. Following their analysis of Cuckoo hashing [J.ACM’12], Pătraşcu and Thorup claimed that the expected maximum load with simple tabulation is $O(\lg \lg n)$. This did not include any high probability guarantee, so the load balancing was not yet to be trusted. Here, we show that with simple tabulation, the maximum load is $O(\lg \lg n)$ with high probability, giving the first constant time hash function with this guarantee. We also give a concrete example where, unlike with fully random hashing, the maximum load is not bounded by $\lg \lg n + O(1)$, or even $(1 + o(1)) \lg \lg n$ with high probability. Finally, we show that the expected maximum load is $\lg \lg n + O(1)$, just like with fully random hashing.

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CP23
Distance in the Forest Fire Model How Far Are You from Eve?

Leskovec, Kleinberg and Faloutsos (2005) observed that many social networks exhibit properties such as shrinking (i.e. bounded) diameter, densification, and (power-law) heavy tail degree distributions. To explain these phenomena, they introduced a generative model, called the Forest Fire model, and using simulations showed that this model indeed exhibited these properties; however, proving this rigorously was left as an open problem. In this paper, we analyse one of these properties, shrinking diameter. We define a restricted version of their model that incorporates the main features that seem to contribute towards this property, and prove that the graphs generated by this model exhibit shrinking distance to the seed graph. We prove that an even simpler model, the random walk model, already exhibits this phenomenon.

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CP23
Species Trees from Gene Trees Despite a High Rate of Lateral Genetic Transfer: A Tight Bound

Reconstructing the tree of life from molecular sequences is a fundamental problem in computational biology. Modern data sets often contain a large number of genes which can complicate the reconstruction problem due to the fact that different genes may undergo different evolutionary histories. This is the case in particular in the presence of lateral genetic transfer (LGT), whereby a gene is inherited from a distant species rather than an immediate ancestor. Such an event produces a gene tree which is distinct from (but related to) the species phylogeny. In previous work, a stochastic model of LGT was introduced and it was shown that the species phylogeny can be reconstructed from gene trees despite surprisingly high rates of LGT. Both lower and upper bounds on this rate were obtained, but a large gap remained. Here we close this gap, up to a constant.

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CP24
Tight Conditional Lower Bounds for Counting Perfect Matchings on Graphs of Bounded Treewidth, Cliquewidth, and Genus

Assuming the Strong Exponential-Time Hypothesis, we show that the problem of counting perfect matchings has no $(2-\epsilon)^k n^{O(1)}$ time algorithm for any $\epsilon > 0$ on graphs of treewidth $k$, whereas it is known to be solvable in time $2^n O(1)$ if a tree decomposition of width $k$ is given. We also show that the problem has no $O(n^{(1-\epsilon)k})$ time algorithm for any $\epsilon > 0$ on graphs of cliquewidth $k$, but it can be solved in time $O(n^{k+1})$ if a $k$-expression is given. Furthermore, assuming the Exponential-Time Hypothesis, we show that there is no $2^\omega(k) n^{O(1)}$ time algorithm for the problem on graphs of genus $k$.

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CP24
Tight Bounds for Graph Homomorphism and Sub-
**graph Isomorphism**

We prove that unless Exponential Time Hypothesis (ETH) fails, deciding if there is a homomorphism from graph $G$ to graph $H$ cannot be done in time $|V(H)|^{|V(G)|}$. We also show an exponential-time reduction from graph homomorphism to subgraph isomorphism. This rules out (subject to ETH) a possibility of $|V(H)|^{|V(H)|}$-time algorithm deciding if graph $G$ is a subgraph of $H$. For both problems our lower bounds asymptotically match the running time of brute-force algorithms trying all possible mappings of one graph into another. Thus, our work closes the gap in the known complexity of these fundamental problems.

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**CP24**

**Constructive Algorithm for Path-Width of Matroids**

We present a fixed-parameter tractable algorithm to construct a linear layout $V_1, V_2, \ldots, V_n$ of the subspaces such that $\dim((V_1 + V_2 + \cdots + V_i) \cap (V_{i+1} + \cdots + V_n)) \leq k$ for all $i$, if it exists, for input subspaces of a finite-dimensional vector space over $F$. When restricted to 1-dimensional subspaces, this problem is equivalent to computing the path-width of an $F$-represented matroid in matroid theory and computing the trellis-width of a linear code in coding theory.

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**CP24**

**Blocking Optimal $k$-Arborescences**

Given a digraph $D = (V, A)$ and a positive integer $k$, an arc set $F \subseteq A$ is called a $k$-arborescence if it is the disjoint union of $k$ spanning arborescences. The problem of finding a minimum cost $k$-arborescence is known to be polynomial-time solvable using matroid intersection. In this paper we study the following problem: find a minimum cardinality subset of arcs that contains at least one arc from every minimum cost $k$-arborescence. For $k = 1$, the problem was solved in [A. Bernáth, G. Pap, Blocking optimal arborescences, IPCO 2013]. In this paper we give an algorithm for general $k$ that has polynomial running time if $k$ is fixed.

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**CP24**

**Discovering Archipelagos of Tractability for Constraint Satisfaction and Counting**

This paper addresses the general limit of standard tractability results for CSP and #CSP, that they only apply to instances where all constraints belong to a single tractable language. For this purpose we utilize the notion of a strong backdoor of a CSP instance, as introduced by Williams et al. (IJCAI 2003). We consider strong backdoors into scattered classes, consisting of CSP instances where each connected component belongs entirely to some class from a list of tractable classes and show that one can solve a given CSP instance in time fixed-parameter tractable in the size of the strong backdoor.

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**CP25**

**Nearly Optimal NP-Hardness of Unique Coverage**

The Unique Coverage problem, given a universe $V$ of elements and a collection $E$ of subsets of $V$, asks to find $S \subseteq V$ to maximize the number of $e \in E$ that intersects $S$ in exactly one element. When each $e \in E$ has cardinality at most $k$, it is also known as 1-in-$k$ Hitting Set, and admits a simple $\Omega(\frac{1}{\log k})$-approximation algorithm. For constant $k$, we prove that 1-in-$k$ Hitting Set is NP-hard to approximate within a factor $O(\frac{1}{\log k})$. This improves the result of Guruswami and Zhou [SODA’11, ToC’12], who proved the
same result assuming the Unique Games Conjecture. For Unique Coverage, we prove that it is hard to approximate within a factor \(O\left(\frac{\log \frac{1}{\epsilon}}{n}\right)\) for any \(\epsilon > 0\), unless NP admits quasipolynomial time algorithms. This improves the results of Demaine et al. [SODA’06, SICOMP’08], including their \(\approx 1/\log^{1/3} n\) inapproximability factor which was proven under the Random 3SAT Hypothesis. Our simple proof combines ideas from two classical inapproximability results for Set Cover and Constraint Satisfaction Problem, made efficient by various derandomization methods based on bounded independence.

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CP25
Algorithms and Adaptivity Gaps for Stochastic Probing

A stochastic probing problem consists of a set of elements whose values are independent random variables. The algorithm knows the distributions of these variables, but not the actual outcomes. The only way to learn the actual outcomes is to probe these elements. However, there are constraints on which set of elements may be probed. (E.g., we may have to travel in some metric to probe elements but have limited time.) These constraints are called outer constraints. We want to develop a strategy that allows us to maximize the (expected) value, subject to the picked subset of elements satisfying some other set of constraints, called the inner constraints. In the past, probing problems were studied for the case when both inner and outer constraints were intersections of matroids; these modeled kidney matching and Bayesian auctions applications. One limitation of past work was their reliance on linear-programming-like techniques, which made going beyond matroid-like structures difficult. In this work, we give a very general adaptivity gap result that holds for all prefix-closed outer constraints, as long as the inner constraints are intersections of matroids. The adaptivity gap is \(O(\log n)\) for any constant number of inner matroid constraints. The prefix-closedness captures most ‘reasonable’ outer constraints, like orienteering, connectivity, and precedence. Based on this we obtain the first approximation algorithms for a number of stochastic probing problems, which have applications, e.g., to path-planning and precedence-constrained scheduling.

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CP25
Discrete Gaussian Sampling Reduces to CVP and SVP

The discrete Gaussian \(D_{L-t,x}\) is the distribution that assigns to each vector \(x\) in a shifted lattice \(L - t\) probability proportional to \(e^{-\|x\|^2/s^2}\). It has long been an important tool in the study of lattices. More recently, algorithms for discrete Gaussian sampling (DGS) have found many applications in computer science. In particular, polynomial-time algorithms for DGS with very high parameters \(s\) have found many uses in cryptography and in reductions between lattice problems. And, in the past year, Aggarwal, Dadush, Regev, and Stephens-Davidowitz showed a polynomial-time algorithms for DGS with a much wider range of parameters and used them to obtain the current fastest known algorithms for the two most important lattice problems, the Shortest Vector Problem (SVP) and the Closest Vector Problem (CVP). Motivated by its increasing importance, we investigate the complexity of DGS itself and its relationship to CVP and SVP. Our first result is a polynomial-time dimension-preserving reduction from DGS to CVP. There is a simple reduction from CVP to DGS, so this shows that DGS is equivalent to CVP. Our second result, which we find to be more surprising, is a polynomial-time dimension-preserving reduction from centered DGS (the important special case when \(t = 0\)) to SVP. In the other direction, there is a simple reduction from \(\gamma\)-approximate SVP for any \(\gamma = \Omega(\sqrt{n/\log n})\), and we present some (relatively weak) evidence to suggest that this might be the best achievable approximation factor.

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CP25
Approximation of Nonboolean 2CSP

We develop a polynomial time \(\Omega(\frac{1}{\epsilon} \log R)\) approximate algorithm for Max 2CSP\(-R\), the problem where we are given a collection of constraints, each involving two variables, where each variable ranges over a set of size \(R\), and we want to find an assignment to the variables that maximizes the number of satisfied constraints. Assuming the Unique Games Conjecture, this is the best possible approximation up to constant factors.

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CP25
An FPTAS for Minimizing Indefinite Quadratic Forms over Integers in Polyhedra

We present a generic approach that allows us to develop a fully polynomial-time approximation scheme (FPTAS) for minimizing nonlinear functions over the integer points in a rational polyhedron in fixed dimension. The approach combines the subdivision strategy of Papadimitriou and Yannakakis (2000) with ideas similar to those commonly used to derive real algebraic certificates of positivity for polynomials. Our general approach is widely applicable. We apply it, for instance, to the Motzkin polynomial and to indefinite quadratic forms \(x^T Q x\) in a fixed number of variables, where \(Q\) has at most one positive, or at most
one negative eigenvalue. In dimension three, this leads to an FPTAS for general $Q$.

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CP26
Finding Perfect Matchings in Bipartite Hypergraphs

Haxell’s condition is a natural hypergraph analog of Hall’s condition, which is a well-known necessary and sufficient condition for a bipartite graph to admit a perfect matching. That is, when Haxell’s condition holds it forces the existence of a perfect matching in the bipartite hypergraph. Unlike in graphs, however, there is no known polynomial time algorithm to find the hypergraph perfect matching that is guaranteed to exist when Haxell’s condition is satisfied. We prove the existence of an efficient algorithm to find perfect matchings in bipartite hypergraphs whenever a stronger version of Haxell’s condition holds. Our algorithm can be seen as a generalization of the classical Hungarian algorithm for finding perfect matchings in bipartite graphs. The techniques we use to achieve this result could be of use more generally in other combinatorial problems on hypergraphs where disjointness structure is crucial, e.g. Set Packing.

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CP26
Obstructions for Three-Coloring Graphs with One Forbidden Induced Subgraph

The complexity of coloring graphs without long induced paths is a notorious problem in algorithm graph theory, an especially intriguing case being that of 3-colorability. So far, not much was known about certification in this context. We prove that there are only finitely many 4-critical $P_6$-free graphs, and give the complete list that consists of 24 graphs. In particular, we obtain a certifying algorithm for 3-coloring $P_6$-free graphs, which solves an open problem posed by Golovach et al. Here, $P_6$ denotes the induced path on six vertices. Our result leads to the following dichotomy theorem: if $H$ is a connected graph, then there are finitely many 4-critical $H$-free graphs if and only if $H$ is a subgraph of $P_6$. This answers a question of Seymour. The proof of our main result involves two distinct automatic proofs, and an extensive structural analysis by hand.

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CP26
An Algorithmic Hypergraph Regularity Lemma

Szemerédi’s Regularity Lemma is a powerful tool in graph theory. It asserts that all large graphs $G$ admit a bounded partition of $E(G)$, most classes of which are bipartite subgraphs with uniformly distributed edges. The original proof of this result was non-constructive. A constructive proof was given by Alon, Duke, Lefmann, Rödl and Yuster, which allows one to efficiently construct a regular partition for any large graph. Szemerédi’s Regularity Lemma was extended to hypergraphs by various authors. Frankl and Rödl gave one such extension to 3-uniform hypergraphs, and Rödl and Skokan extended this result to $k$-uniform hypergraphs. W.T. Gowers gave another such extension. Similarly to the graph case, all of these proofs are non-constructive. We present an efficient algorithmic version of the Hypergraph Regularity Lemma for $k$-uniform hypergraphs.

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CP26
Independence and Efficient Domination on $P_k$-Free Graphs

In the Max Independent Set problem, the input is a graph $G$, every vertex has a non-negative integer weight, and the task is to find a set $S$ of pairwise non-adjacent vertices, maximizing the total weight of the vertices in $S$. We give an $n^{O(k^2 \log k)}$ time algorithm for this problem on graphs excluding the path $P_k$ on 6 vertices as an induced subgraph. Currently, there is no constant $k$ known for which Max Independent Set on $P_k$-free graphs becomes NP-complete, and our result implies that if such a $k$ exists, then $k > 6$ unless all problems in NP can be decided in (quasi)polynomial time. Using the combinatorial tools that we develop for the above algorithm, we also give a polynomial-time algorithm for Efficient Domination on $P_k$-free graphs. In this problem, the input is a graph $G$, every vertex has an integer weight, and the objective is to find a set $S$ of maximum weight such that every vertex in $G$ has exactly one vertex in $S$ in its closed neighborhood, or to determine that no such set exists. Prior to our work, the class of $P_6$-free graphs was the only class of graphs defined by a single forbidden induced subgraph on which the computational
complexity of Efficient Domination was unknown.

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CP26
Sparsity and Dimension

We prove that posets of bounded height whose cover graphs belong to a fixed class with bounded expansion have bounded dimension. Bounded expansion, introduced by Nešetřil and Ossona de Mendez as a model for sparsity in graphs, is a property that is naturally satisfied by a wide range of graph classes, from graph structure theory (graphs excluding a minor or a topological minor) to graph drawing (e.g. graphs with constant book thickness). Therefore, our theorem generalizes a number of results including the most recent one for posets of bounded height with cover graphs excluding a fixed graph as a topological minor (Walczak, SODA 2015). We also show that the result is in a sense best possible, as it does not extend to nowhere dense classes; in fact, it already fails for cover graphs with locally bounded treewidth.

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CP27
Using Optimization to Obtain a Width-Independent, Parallel, Simpler, and Faster Positive Sdp Solver

We study the design of polylogarithmic depth algorithms for approximately solving packing and covering semidefinite programs (or positive SDPs for short). This is a natural SDP generalization of the well-studied positive LP problem. Although positive LPs can be solved in polylogarithmic depth while using only \( \log^3 n/cps^3 \) parallelizable iterations [3], the best known positive SDP solvers due to Jain and Yao [18] require \( \log^{14}(n)/cps^{13} \) parallelizable iterations. Several alternative solvers have been proposed to reduce the exponents in the number of iterations [19, 29]. However, the correctness of the convergence analyses in these works has been called into question [29], as they both rely on algebraic monotonicity properties that do not generalize to matrix algebra. In this paper, we propose a very simple algorithm based on the optimization framework proposed in [3] for LP solvers. Our algorithm only needs \( \log^2 n/cps^8 \) iterations, matching that of the best LP solver. To surmount the obstacles encountered by previous approaches, our analysis requires a new matrix inequality that extends Lieb-Thirring’s inequality, and a sign-consistent, randomized variant of the gradient truncation technique proposed in [2, 3].

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CP27
Improved Cheeger’s Inequality Using Vertex Expansion and Expansion Profile

We prove two generalizations of the Cheeger’s inequality. The first generalization relates the second eigenvalue to the edge expansion and the vertex expansion of the graph \( G \),

\[
\lambda_2 = \Omega(\phi^V(G) \cdot \phi(G)),
\]

where \( \phi^V(G) \) denotes the robust vertex expansion of \( G \) and \( \phi(G) \) denotes the edge expansion of \( G \). The second generalization relates the second eigenvalue to the edge expansion and the expansion profile of \( G \), for all \( k \geq 2 \),

\[
\lambda_2 = \Omega(\frac{1}{k} \cdot \phi_k(G) \cdot \phi(G)),
\]

where \( \phi_k(G) \) denotes the \( k \)-way expansion of \( G \). These show that the spectral partitioning algorithm has better performance guarantees when \( \phi^V(G) \) is large (e.g. planted random instances) or \( \phi_k(G) \) is large (instances with few disjoint non-expanding sets). Both bounds are tight up to a constant factor.

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CP27
Approximate Undirected Maximum Flows in \( O(m \text{ Polylog}(n)) \) Time

We give the first \( O(m \text{ Polylog}(n)) \) time algorithms for approximating maximum flows in undirected graphs and constructing polylog(n)-quality cut-approximating hierarchical tree decompositions. Our algorithm invokes existing algorithms for these two problems recursively while gradually incorporating size reductions. These size reductions are in turn obtained via ultra-sparsifiers, which are key tools in solvers for symmetric diagonally dominant (SDD) linear systems.

Richard Peng
CP27
How to Round Subspaces: A New Spectral Clustering Algorithm

We propose a new spectral clustering algorithm, which can recover a $k$-partition whose span is $O(\sqrt{OPT})$ close to the input in spectral norm. This implies an algorithm for approximating $k$-expansion in a graph where each cluster have expansion $O(\lambda_{k+1})$ ($\lambda_{k+1}$ being $(k+1)^{th}$ eigenvalue of Laplacian). This significantly improves upon the previous algorithms, which required $O(\lambda_{k+1}/k)$.

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CP27
Natural Algorithms for Flow Problems

In the last few years, there has been a significant interest in the computational abilities of Physarum polycephalum (a slime mold). This arose from a remarkable experiment which showed that this organism can compute shortest paths in a maze [NTY '00]. Subsequently, the workings of Physarum were mathematically modeled as a dynamical system and algorithms inspired by this model were proposed to solve several graph problems: shortest paths, flows, and linear programs to name a few. Indeed, computer scientists have initiated a rigorous study of these dynamics and a first step towards this was taken by BMV '12 and BBDKM '13 who proved that the Physarum dynamics for the shortest path problem are efficient (when edge-lengths are polynomially bounded). In this paper, we take this further: we prove that the discrete time Physarum-dynamics can also efficiently solve the uncapacitated min-cost flow problems on undirected and directed graphs; problems that are non-trivial generalizations of shortest path. This raises the tantalizing possibility that nature, via evolution, developed algorithms that efficiently solve some of the most complex computational problems, about a billion years before we did.

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CP28
Robust Positioning Patterns

In this paper, we construct large sequences and matrices with the property that the contents of any small window determine the location of the window, robustly. Such objects have found many applications in practical settings, from positioning of wireless devices to smart pens, and have recently gained some theoretical interest. In this context, we give the first explicit constructions of sequences and matrices with high rate and constant relative distance. Accompanying these efficient constructions, we also give efficient decoding algorithms, which can determine the position of the window given its contents, even if a constant fraction of the contents have been corrupted.

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CP28
Towards Optimal Deterministic Coding for Interactive Communication

We study efficient, deterministic interactive coding schemes that simulate any interactive protocol both under random and adversarial errors, and can achieve a constant communication rate independent of the protocol length. For channels that flip bits independently with probability $\epsilon < 1/2$, our coding scheme achieves a communication rate of $1 - O(\sqrt{H(\epsilon)})$ and a failure probability of $\exp(-n/\log n)$ in length $n$ protocols. Prior to our work, all nontrivial deterministic schemes (either efficient or not) had a rate bounded away from 1. Furthermore, the best failure probability achievable by an efficient deterministic coding scheme with constant rate was only quasi-polynomial, i.e., of the form $\exp(-\log^{O(1)} n)$ (Braverman, ITCS 2012).

For channels in which an adversary controls the noise pattern our coding scheme can tolerate $\Omega(1/\log n)$ fraction of errors with rate approaching 1. Once more, all previously known nontrivial deterministic schemes (either efficient or not) in the adversarial setting had a rate bounded away from 1, and no nontrivial efficient deterministic coding schemes were known with any constant rate. Essential to both results is an explicit, efficiently encodable and decodable systematic tree code of length $n$ that has relative distance $\Omega(1/\log n)$ and rate approaching 1, defined over an $O(\log n)$-bit alphabet. No nontrivial tree code (either efficient or not) was known to approach rate 1, and no nontrivial distance bound was known for any efficient constant rate tree code.

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CP28
An Improved Bound on the Fraction of Correctable Deletions

We consider codes over fixed alphabets against worst-case symbol deletions. For any fixed $k \geq 2$, we construct a family of codes over alphabet of size $k$ with positive rate, which allow efficient recovery from a worst-case deletion fraction approaching $1 - \frac{\epsilon}{k}$. In particular, for binary codes, we are able to recover a fraction of deletions approaching...
1/3. Previously, even non-constructively the largest deletion fraction known to be correctable with positive rate was $1 - \Theta(1/\sqrt{k})$, and around 0.17 for the binary case. Our result pins down the largest fraction of correctable deletions for $k$-ary codes as $1 - \Theta(1/k)$, since $1 - 1/k$ is an upper bound even for the simpler model of erasures where the locations of the missing symbols are known. Closing the gap between 1/3 and 1/2 for the limit of worst-case deletions correctable by binary codes remains a tantalizing open question.

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CP28
Efficient Low-Redundancy Codes for Correcting Multiple Deletions

We consider the problem of constructing binary codes to recover from $k$-bit deletions with efficient encoding/decoding, for a fixed $k$. The single deletion case is well understood, with the Varshamov-Tenengolts-Levenshtein code from 1965 giving an asymptotically optimal construction with $2^{n/2}$ codewords of length $n$, i.e., at most $\log n$ bits of redundancy. However, even for the case of two deletions, there was no known explicit construction with redundancy less than $n^{O(1)}$. For any fixed $k$, we construct a binary code with $c_k \log n$ redundancy that can be decoded from $k$ deletions in $O_k(n \log^2 n)$ time. The coefficient $c_k$ can be taken to be $O(k^2 \log k)$, which is only quadratically worse than the optimal, non-constructive bound of $O(k)$. We also indicate how to modify this code to allow for a combination of up to $k$ insertions and deletions. We also note that among linear codes capable of correcting $k$ deletions, the $(k+1)$-fold repetition code is essentially the best possible.

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CP28
Communication Complexity of Permutation-Invariant Functions

Motivated by the quest for a broader understanding of upper bounds in communication complexity, at least for simple functions, we introduce the class of "permutation-invariant" functions. A partial function $f : \{0,1\}^n \times \{0,1\}^n \to \{0,1,\}^*$ is permutation-invariant if for every bijection $\pi : \{1,\ldots,n\} \to \{1,\ldots,n\}$ and every $x,y \in \{0,1\}^n$, it is the case that $f(x,y) = f(x^{\pi(1)},y^{\pi(2)})$. Most of the commonly studied functions in communication complexity are permutation-invariant. For such functions, we present a simple complexity measure (computable in time polynomial in $n$ given an implicit description of $f$) that describes their communication complexity up to polynomial factors and up to an additive error that is logarithmic in the input size. This gives a coarse taxonomy of the communication complexity of simple functions. Our work highlights the role of the well-known lower bounds of functions such as ‘Set-Disjointness’ and ‘Indexing’, while complementing them with the relatively lesser-known upper bounds for ‘Gap-Inner-Product’ (from the sketching literature) and ‘Sparse-Gap-Inner-Product’ (from the recent work of Canonne et al. [ITCS 2015]). We also present consequences to the study of communication complexity with imperfectly shared randomness where we show that for total permutation-invariant functions, imperfectly shared randomness results in only a polynomial blow-up in communication complexity after an additive $O(\log \log n)$ overhead.

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CP29
Sampling on Lattices with Free Boundary Conditions Using Randomized Extensions

Many statistical physics models are defined on an infinite lattice by taking appropriate limits of finite lattice regions, where a key consideration is how the boundaries are defined. For several models on planar lattices, such as 3-colorings and lozenge tilings, efficient sampling algorithms are known for regions with fixed boundary conditions, where the colors or tiles around the boundary are pre-specified (Luby et al., 2002), but much less is known about how to sample when these regions have free boundaries, where we want to include all configurations one could see within a finite window. We introduce a method using randomized extensions of a lattice region to relate sampling problems on regions with free boundaries to a constant number of sampling problems on larger regions with fixed boundaries. We demonstrate this principled approach to sample 3-colorings of regions of $Z^2$ and lozenge tilings of regions of the triangular lattice, building on arguments for the fixed boundary cases due to Luby et al. Our approach also yields an efficient algorithm for sampling 3-colorings with free boundary conditions on regions with one reflex corner, the first such result for a nonconvex region. Our approach can also be generalized to a broad class of mixed boundary conditions. Sampling for these families of regions is significant because it allows us to establish self-reducibility, giving the first algorithm to approximately count the total number of 3-colorings of rectangular lattice regions.

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CP29
On the Switch Markov Chain for Perfect Matchings

We study a simple Markov chain, the switch chain, on the set of all perfect matchings in a bipartite graph. This Markov chain was proposed by Diaconis, Graham and Holmes as a possible approach to a sampling problem arising in Statistics. They considered several classes of graphs, and conjectured that the switch chain would mix rapidly for graphs in these classes. Here we settle their conjecture almost completely. We ask: for which graph classes is the Markov chain ergodic and for which is it rapidly mixing? We provide a precise answer to the ergodicity question and close bounds on the mixing question. We show for the first time that the mixing time of the switch chain is polynomial...
in the class of monotone graphs. This class was identified by Diaconis, Graham and Holmes as being of particular interest in the statistical setting.

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CP29
Algorithmic and Enumerative Aspects of the Moser-Tardos Distribution

Moser & Tardos have developed a powerful algorithmic approach (henceforth “MT”) to the Lovasz Local Lemma (LLL); the basic operation done in MT and its variants is a search for “bad” events in a current configuration. In the initial stage of MT, the variables are set independently. We examine the distributions on these variables which arise during intermediate stages of MT. We show that these configurations have a more or less “random” form, building further on the “MT-distribution” concept of Haeupler et al., in understanding the (intermediate and) output distribution of MT. This has a variety of algorithmic applications; the most important is that bad events can be found relatively quickly, improving upon MT across the complexity spectrum: it makes some polynomial-time algorithms sublinear (e.g., for Latin transversals, which are of basic combinatorial interest), gives lower-degree polynomial run-times in some settings, transforms certain super-polynomial-time algorithms into polynomial-time ones, and leads to Las Vegas algorithms for some coloring problems for which only Monte Carlo algorithms were known. We show that in certain conditions when the LLL condition is violated, a variant of the MT algorithm can still produce a distribution which avoids most of the bad events. We show in some cases this MT variant can run faster than the original MT algorithm itself, and develop the first-known criterion for the case of the asymmetric LLL. This can be used to find partial Latin transversals – improving upon earlier bounds of Stein (1975) – among other applications. We furthermore give applications in enumeration, showing that most applications (where we aim for all or most of the bad events to be avoided) have many more solutions than known before by proving that the MT-distribution has “large” Renyi entropy and hence that its support-size is large.

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CP30
Weak Duality for Packing Edge-Disjoint Odd $(u,v)$-Trails

There is no packing-covering duality for odd edge-disjoint $(u,v)$-paths: a graph with no two edge-disjoint $(u,v)$-paths may need an arbitrarily large number of edges to cover all such paths. In contrast, we show that the relaxed problem of packing odd trails does have an approximate duality. The proof leads to a polynomial-time algorithm to find, for any given $k$, either $k$ edge-disjoint odd $(u,v)$-trails or a...
set of fewer than 8k edges intersecting all odd $(u, v)$-trails.

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CP30
The $k$-Mismatch Problem Revisited
We revisit the complexity of one of the most basic problems in pattern matching. In the $k$-mismatch problem we must compute the Hamming distance between a pattern and every pattern-length substring of a text, as long as that Hamming distance is at most $k$. Whenever the Hamming distance is greater than $k$, we simply output “No”. We give improved algorithms for this problem in both the standard offline setting and also in the streaming model.

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CP30
Communication with Contextual Uncertainty
We introduce a simple model illustrating the role of context in communication and the challenge posed by uncertainty of knowledge of context. We consider a variant of distributional communication complexity where Alice gets some information $x$ and Bob gets $y$, where $(x, y)$ is drawn from a known distribution, and Bob wishes to compute some function $g(x, y)$ (with high probability over $(x, y)$). In our variant, Alice does not know $y$, but only knows some function $f(x)$ which is an approximation of $g$. Thus, the function being computed forms the context for the communication, and knowing it imperfectly models (mild) uncertainty in this context. A naive solution would be for Alice and Bob to first agree on some common function $h$ that is close to both $f$ and $g$ and then use a protocol for $h$ to compute $h(x, y)$. We show that any such agreement leads to a large overhead in communication ruling out such a universal solution. In contrast, we show that if $g$ has a one-way communication protocol with complexity $k$ in the standard setting, then it has a communication protocol with complexity $O(k \cdot (1 + I))$ in the uncertain setting, where $I$ denotes the mutual information between $x$ and $y$. In the particular case where the input distribution is a product distribution, the protocol in the uncertain setting only incurs a constant factor blow-up in communication and error. Furthermore, we show that the dependence on the mutual information $I$ is required. Namely, we construct a class of functions along with a non-product distribution over $(x, y)$ for which the communication complexity is a single bit in the standard setting but at least $\Omega(\sqrt{\sigma})$ bits in the uncertain setting.

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CP30
On the Maximum Quartet Distance Between Phylogenetic Trees
A conjecture of Bandelt and Dress states that the maximum quartet distance between any two phylogenetic trees on $n$ leaves is at most $(\frac{\sigma}{2} + o(1))^{(1)}$. Using the machinery of flag algebras we improve the currently known bounds regarding this conjecture, in particular we show that the maximum is at most $(0.69 + o(1))^{(1)}$. We also give further evidence that the conjecture is true by proving that the maximum distance between caterpillar trees is at most $(\frac{\sigma}{2} + o(1))^{(1)}$.

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CP30
Range Predecessor and Lempel-Ziv Parsing
The Lempel-Ziv (LZ77) parsing of a string is one of the most important algorithmic tools in data compression and string processing. We show that the Lempel-Ziv parsing of a string of length $n$ on an alphabet of size $\sigma$ can be computed in $O(n \log \log \sigma)$ time using $O(n \log \sigma)$ bits of working space. The previous fastest algorithm using $O(n \log \sigma)$ space takes $O(n \log \sigma + \log \log n)$ time. We also consider the important rightmost variant of the problem, where the goal is to associate with each phrase of the parsing its most recent occurrence in the input string. We solve this problem in $O(n(1 + \log \sigma/\log n))$ time, using the same working space as above. The previous best solution for rightmost parsing uses $O(n(1 + \log \sigma/\log n))$ time and $O(n \log n)$ space. As a bonus, we provide a faster construction method for efficient 2D orthogonal range reporting, which is of independent interest.

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Wulff-Nilsen, Christian, CP6, 2:25 Sun

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Yang, Jiyan, CP9, 5:20 Sun
Yokoi, Yu, CP16, 4:55 Mon
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