Final Program

Third SIAM Conference on Parallel Processing for Scientific Computing

DECEMBER 1-4, 1987

The Westin Bonaventure Hotel Los Angeles, CA

CONFERENCE THEMES

- **■** High-Speed Computer Architectures
- **■** Computer Performance Evaluation
- **■** High-Speed Computer Environments
- Numerical Domain Decomposition Methods
- **■** Scientific Programming Languages
- Numerical Linear Algebra
- **■** Distributed Computing
- **■** Computational Complexity
- Adaptive Numerical Methods
- Numerical Particle Methods

Sponsored by the SIAM Activity Group on Supercomputing

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ORGANIZING COMMITTEE

Garry Rodrigue, Chair Lawrence Livermore National Laboratory and University of California, Davis

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FUNDING AGENCIES

SIAM is conducting this conference with the partial support of the Department of Energy and the National Science Foundation.

MEETING HIGHLIGHTS

INVITED PRESENTATIONS

Tuesday, December 1, 8:30 AM Invited Presentation 1

Parallel Computing Via Multidomain Techniques for Spectral Methods

Multidomain techniques have become a popular way for the application of spectral methods to simulating flow on complex geometries. It turns out that these techniques suggest a natural way of doing the computations in parallel, namely each subdomain is treated in a different processor. The way of transferring the information from one processor to another becomes crucial, especially for hyperbolic problems that arise from simulating compressible flows.

The author will review spectral multidomain techniques in the context of parallel computations. Ways of improving boundary conditions were explored that will speed up convergence. In particular, ways of composing matching conditions were tested for hyperbolic equations — in particular the Euler equations of gas dynamics.

The plan is to use the Flex machine at ICASE to test the validity and efficiency of the multidomain technique. This research attempts to give an answer to the question whether it is preferable to parallelize the algorithm without physically subdividing the domain of computations, or whether it is more gainful to divide the domain and to treat the problem in different subdomain.

In both cases, the connection between boundary conditions and transfer of information between processor will be discussed and clarified.

David Gottlieb Brown University

Tuesday, December 1, 9:15 AM Invited Presentation 2

Adaptive Mesh Refinement for Parallel Processors

Adaptive algorithms can be difficult to parallelize efficiently. This is because: (1) adaptive methods use more complicated data structures than arrays, making parallel access to these data structures difficult; (2) adaptive methods concentrate the work in localized regions, making load balancing on a multiprocessor difficult.

In this work, an adaptive mesh refinement algorithm for solving the Euler equations is implemented on a Cray XMP with four processors. This method of adaptive refinement is based on the use of logically rectangular fine grid patches embedded in a coarse, underlying grid. This method of nesting multiple grids allows the bulk of the computational work to be done on regular grids, although the complete grid hierarchy is irregular. We discuss both a fine grained and coarse grained approach to parallelizing the algorithm. A binary decomposition strategy is used to load balance the computation. The binary decomposition algorithm is easy to implement, but does not take into account the communication overhead that would occur on a non-shared memory machine. We show results, and discuss the overhead of

this method. Surprisingly, the localization of work, combined with the adaptive data structures, make implementation of the binary decomposition an easy task.

Marsha J. Berger Courant Institute of Mathematical Sciences New York University

Tuesday, December 1, 1:30 PM Invited Presentation 3 Neural Computing

The questions to be addressed in this talk are two-fold: what is the appropriate set of tools for the study of the highly parallel computing networks in the human brain; and what are the strategies for building highly parallel computers suitable for programming with "intelligence." The thrust of this paper is that there overall architectural principles that unite both sides of this study: namely, that a computer no longer be thought of as a unitary system, but rather as a network of more specialized devices: and that many of these devices be structured as highly parallel arrays of interacting neuron like components. We shall illustrate this with a brief discussion of the following topics: the high level architecture of the frog's brain as revealed in studies of the mechanisms visual motor coordination; the use of supercomputing in modelling the interacting levels of the retina: stability analysis in the study of neural networks; the implications of associative memory techniques for intelligence computer design.

Michael A. Arbib University of Southern California, Los Angeles

Tuesday, December 1, 2:15 PM Invited Presentation 4 Optical Computing Systems

There are several computing problems that require much more computing power than now is available. Optics can offer potentially greater speed and parallelism. In addition, in several cases, the data are in optical form. Examples of the latter are optical fiber data systems where the electronic control now slows operation and image processing systems where the data are naturally in optical form. Several optical computing systems being built within the Center for Optoelectronic Computing Systems and elsewhere will be discussed.

Examples are a bit-serial machine to explore architectures and the speed limits; an optical cellular array for signal processing; a connection intensive system for pattern recognition and processing; optical symbolic computing for optical artificial intelligence machines; and optical associative memory that uses four-wave mixing and holography.

W. Thomas Cathey University of Colorado, Boulder

INVITED PRESENTATIONS

Wednesday, December 2, 8:30 AM Invited Presentation 5

Automatic Decomposition of Fortran Programs for Execution on Multiprocessors

It is clear that future generations of scientific supercomputers will employ multiple independent processors. What form of programming support software should be provided with such machines? Existing Fortran programs, written for sequential machines, are not well suited to parallel execution. If these programs are to run efficiently on a multiprocessor system, they must be decomposed into subproblems that can be executed in parallel.

The usual approach to decomposition is to provide language primitives or system calls that permit concurrent programming in Fortran. The programmer is thus responsible for handling all the synchronization. Unfortunately, concurrent programming is unnatural for many scientific programmers. It would be better to permit the programmer to design parallelism into his programs at an abstract level and have the programming environment automatically construct synchronization primitives.

In his talk, the author explores automatic techniques for uncovering parallelism in Fortran programs and for translating them to run efficiently on multiprocessor supercomputers. These techniques rely upon a sophisticated theory of dependence in programs developed by Kuck at Illinois and extended for our previous work on automatic vectorization. In addition, the methods rely extensively on intraprocedural and interprocedural data flow analysis.

In order to achieve enough parallelism to make the translation profitable, a number of very ambitious transformations are required, including loop distribution, loop interchange, loop alignment, and loop fusion. The problem of generating optimal code using these transformations has been shown to be intractable, but heuristic techniques have proved effective in practice. The talk will conclude with a discussion of our experience with a prototype translation system, called PFC Plus, developed at Rice.

Ken Kennedy Rice University

Wednesday, December 2, 9:15 AM Invited Presentation 6

Schwarz Splittings and Template Operators

Schwarz' alternating method (SAM) is an old mathematical technique dating from 1860. The method was commonly employed as a tool for theoretical analysis, but it is still not yet widely used for large scale scientific computations. The earlier experiences showed that SAM converged slowly. Our analyses show that some new methods based on SAM are competitive with other solution techniques. Some generalizations of SAM, namely Schwarz splittings (SS), are presented here. Several acceleration schemes are also discussed in this talk. In particular, when these techniques are applied to the solution of the model problem, an optimal complexity can be achieved. For many important applications, such as performing parallel computations in a non-shared memory environment, using composite grids and also

applying fast solvers in an irregular region, SS's are found to be useful techniques.

In order to identify the types of problems for which these new techniques are most suitable, a new structure for the linear operators called Template operators has been developed. Some decay results for the elements of the inverses of sparse operators are given. These results provide a theoretical basis for determining when these SS techniques can be used successfully.

Wei Pai Tang Stanford University

Wednesday, December 2, 1:30 PM Invited Presentation 7

Parallel Systems: Performance Modeling and Numerical Algorithms

Historically, there have been two major techniques for modeling parallel systems: discrete event simulation and mathematical analysis. Although simulation models can mimic a real-world system as closely as understanding permits and needs require, highly detailed simulation models can be computationally taxing. In contrast, analytic techniques often can quickly provide mathematical insight into the behavior of systems over a broad range of parameter values. Their major limitation is the number of restrictive assumptions that often must be satisfied to insure tractability and accuracy. For many parallel systems, evaluation of realistic analytic models is as complex or even more complex than equivalent simulations.

Transient behavior, complex access patterns, and large state spaces all make analysis difficult. Eliding unnecessary detail while retaining important features is the essence of the modeler's art. The author presents a set of techniques for analyzing the performance of parallel architectures. These techniques permit the comparative evaluation of interconnection networks for message-based processors (e.g., hypercubes) as a function of communication pattern, including communication locality. In addition, the author shows how simple analytic techniques can capture the salient details of network traffic and memory conflicts in shared memory processors.

Daniel A. Reed University of Illinois, Urbana

Wednesday, December 2, 2:15 PM Invited Presentation 8

Asymptotics and Domain Decomposition in Parallel Computing

Mathematical models for scientific phenomena usually involve multiple scales, with the consequence of rapid temporal or spatial variations in the solution. While the details of these rapid variations may not be of interest, the numerical computations must somehow incorporate any effects on those aspects of the solution one is interested to determine.

Where the local behavior of the solution is known, analytic patches may be used, while in other regions special local grids will have to be employed in order to resolve the special features of the solution.

The sub-domains together with rules governing their interactions form a domain

decomposition method. Domain decomposition derived from asymptotic analysis can take advantage of parallel vector architecture and represents, indeed, a truly "new" class of parallel algorithms.

In his talk, the author surveys the asymptotic methods employed to identify rapid varying regions and to analyze their local behaviors. He will discuss the incorporation of the local asymptotic solutions and the partition of subdomains into a numerical scheme via a domain decomposition algorithm. An analysis of the domain decomposition algorithm within the frame-work of asymptotic preconditioning and deferred correction will be presented and the mapping of an asymptotics-induced domain decomposition method onto two parallel processing machines will be discussed.

Raymond C. Y. Chin Lawrence Livermore National Laboratory

Thursday, December 3, 8:30 AM Invited Presentation 9

A Fast Multipole Algorithm for Large-Scale Particle Simulations

The evaluation of Coulombic or gravitational interactions in large-scale ensembles of particles is an integral part of the numerical simulation of a large number of physical processes. Examples include celestial mechanics, plasma physics, the vortex method in fluid dynamics, molecular dynamics, and the solution of the Laplace equation via potential theory. In a typical application, a numerical model follows the trajectories of a number of particles moving in accordance with Newton's second law of motion in a field generated by the whole ensemble. In many situations, in order to be of physical interest, the simulation has to involve thousands of particles (or more), and the fields have to be evaluated for a large number of configurations. Unfortunately, an amount of work of the order O(N2) has traditionally been required to evaluate all pairwise interactions in a system of N particles. unless some approximation or truncation method is used. As a result, large-scale simulations requiring high accuracy (as encountered in highly correlated systems) have been extremely expensive in some cases, and prohibitive in others.

The author presents an algorithm for the rapid evaluation of the potential and force fields in large-scale systems of particles. In order to evaluate all pairwise Coulombic interactions of N particles to within round-off error, the algorithm requires an amount of work proportional to N, and this estimate does not depend on the statistics of the distribution. Both two and three dimensional versions of the algorithm have been constructed, as we will discuss their applications to several problems in physics, chemistry, biology, and numerical complex analysis.

The author concludes with a discussion of the implementation of the algorithm on a variety of parallel architectures and observes that almost all communications required by the procedure are local, making it well-suited for such machines.

Leslie Greengard Yale University

INVITED PRESENTATIONS

Thursday, December 3, 9:15 AM Invited Presentation 10 Lattice Gas Methods for Solving Partial Differential Equations

Lattice gas methods are logical, discrete techniques for efficiently solving partial differential equations. These methods have been shown to solve 2D compressible Navier Stokes, reaction-diffusion, 2D MHD, and Burger's equations. They can exactly conserve mass, energy, and momentum; they are fast (about one billion cell updates per second on a CRAY XMP48); they are memory efficient (one CRAY word defines 10.6 cells). The codes are robust, efficient, and short (approximately 50 FORTRAN lines); complicated boundaries are easily incorporated. They are fully parallel (N processors run N times faster). Estimates indicate that a (512)3 processor computer could cost \$300,000. Small simulators (256 imes 514 cells) have been built for around \$1000. Present lattice gas methods are constrained by Mach numbers that cannot exceed 1.4. density-dependent viscosity, velocity-dependent equations of state, and cell-averaging requirements.

The author will review recent developments, presenting details of 2D and 3D Navier Stokes solvers along with movies of complicated flows and concluding with a discussion of future challenges.

Gary Doolen Los Alamos National Laboratory

Thursday, December 3, 1:30 PM Invited Presentation 11 Strategies for the Mapping and Scheduling of Scientific Computations

Work underway at Yale and at ICASE on automated problem mapping, runtime scheduling of computations and load balancing will be described. Through the analysis of a variety of problems in scientific computing, the following tradeoffs are explored: (1) communication costs and load imbalance in multiprocessors with preferential access to local memories, and (2) synchronization costs and load imbalance in machines where the access time to all memory is approximately equal.

Parameterized mapping schemes are developed that allow one to adjust the granularity of parallelism in several different ways. The chosen mapping influences the balance of load as well as the observed communication and synchronization overheads. Both analytic modeling and empirical studies on the Intel iPSC hypercube and the Encore Multimax are utilized in this work. The prospects for incorporating what has been learned in these studies into the Crystal/ACRE parallel programming environment will be discussed.

Joel Saltz Yale University Thursday, December 3, 2:15 PM Invited Presentation 12

Graph Theory plus Group Therory equals Answers to Communication Problems

A primary obstacle in obtaining significant speedup of highly parallel computer programs on massively parallel distributed processors is global communication. For this reason, we have been investigating theories which enable us to exhibit provably optimal global communication schemes on a wide class of interconnection networks.

One such class for which a theoretical treatment is successful are those networks which can be represented as the directed graph of a group. The elements of the group represent processors (the nodes of the graph) while the generators correspond to communication lines (the arcs of the graph). The author has shown that if the group processes a certain sequencing property P then various global communications tasks (e.g., universal broadcast, accumulation or exchange) can be performed in minimal time.

The hypercube is a prime example of a network which can be represented as the graph of a group (an elementary abelian 2-group) and the author has been able to show that this group does in fact have property P. At this time no groups are known to the author that do not have property P.

Plans are to extend these notions to networks which cannot be represented as the graph of the group, but may still possess optimal global communications schemes. Ultimately, one should expect to be able to measure the communication performance of an interconnection network in these terms, in order to identify those networks that are in some sense "optimal" in terms of global communication.

Vance Faber Los Alamos National Laboratory

Friday, December 4, 8:30 AM Invited Presentation 13 Parallel Processing Research at IBM

Two highly parallel processor prototyping efforts are currently underway at the IBM T.J. Watson Research Center: GF11, a Single Instruction Multiple Data Stream (SIMD) array of processors which can achieve close to 10 billion floating point operations per second on certain computations. The Research Parallel Processor Prototype (RP3), which is a Multiple Instruction Stream Multiple Data Stream (MIMD) system consisting of conventional microprocessors. A high bandwidth, low latency switch together with a flexible memory structure makes this machine potentially suitable for a variety of applications. The talk will briefly describe these two projects and report on their current status. Motivations, goals, and technical plans will be discussed to provide an overall perspective for these efforts. Tilak Agerwala IBM T.J. Watson Research Center

Friday, December 4, 9:15 AM Invited Presentation 14 Programming Environments for Multiprocessor Architectures

Programming multiprocessor architectures is a complex task which requires the user to have thorough knowledge of the machine being targeted. Several different approaches are being explored so as to make this task easier. These range from efforts aimed at automatically transforming sequential programs for parallel execution to designing new languages which will allow the explicit control of the parallelism. The primary goal of all these efforts is to provide an environment which allows the programmer to express the algorithm in a manner which not only leads to efficient execution but also provides some measure of portability across architectures. In his talk the author will survey the different approaches that are currently being explored.

Piyush Mehrotra Purdue University

SPECIAL FUNCTIONS

Welcoming Reception
Monday, November 30, 8:00 PM - 10 PM
Pasadena Exhibit Hall
Westin Bonaventure
Cash Bar

Beef and Beer Party
Tuesday, December 1, 7:30 PM-9:30 PM
Pasadena Exhibit Hall
Westin Bonaventure
\$15.00

Meeting of the SIAM Activity Group on Supercomputing Wednesday, December 2, 7:00 PM Pasadena Exhibit Hall Cash Bar

Monday, November 30/PM

4:00 PM/California Foyer **Registration Opens**

8:00 PM/California Foyer **Registration Closes**

Tuesday, December 1/AM

7:00 AM/California Foyer **Registration Opens**

8:15 AM/San Diego Room **Welcoming Remarks**

8:30 AM/San Diego Room Invited Presentations 1 and 2 Chair: Garry Rodrique

Lawrence Livermore National Laboratory

Parallel Computing Via Multidomain **Techniques for Special Methods**

David Gottlieb Brown University

Adaptive Mesh Refinement for Parallel Processors

Marsha J. Berger

Courant Institute of Mathematical Sciences New York University

10:00 AM/Pasadena Exhibit Hall Coffee

10:40 AM **CONCURRENT SESSIONS**

Tuesday, December 1, 10:40 AM-12:00 Noon Contributed Presentations 1/San Diego Room

Chair: Robert G. Babb, II, Oregon Graduate Center

10:40/46/A1

The Design and Implementation of Parallel Algorithms in Ada

Richard F. Sincovec

University of Colorado, Colorado Springs

11:00/123/A1

DINO — A New Language for Numerical Computation on Local Memory Multiprocessors

Matthew Rosing and Robert B. Schnabel University of Colorado, Boulder

11:20/160/A1

Scientific Parallel Processing with LGDF2 Robert G. Babb, II and David C. DiNucci Oregon Graduate Center

11:40/157/A2

The BF (Boundary-Fitted) Coordinate Transformation Technique of DEQSOL (Differential EQuation SOlver Language) Chisato Konno, Michiru Yamabe, Miyuki Saji and Yukio Umetani

Hitachi Ltd.

Tuesday, December 1, 10:40 AM-12:00 Noon Contributed Presentations 2/Sacramento Room

Chair: Kai Hwang, University of Southern California

10:40/105/A2

Communication Complexity of Matrix Multiplication

David C. Fisher Harvey Mudd College

11:00/19/A2

Orderings for Parallel Sparse Symmetric Factorization

Charles E. Leiserson, Massachusetts Institute of Technology and Thinking Machines Corporation and John G. Lewis, Boeing Computer Services and Massachusetts Institute of Technology Center for Advanced Engineering Study

11:20/162/A3

An Efficient Fixed Size Array for Solving Large Scale Toeplitz Systems

Afshin Daghi, V. K. Prasanna Kumar and Ali Safavi

University of Southern California

11:40/188/A3

Architecture and Operation of a Systolic Sparse Matrix Engine

Robert J. Dunki-Jacobs, Steven W. Hammond, Robert M. Hardy and Terry M. Topka, GE Corporate Research and Development Center

Tuesday, December 1, 10:40 AM-12:00 Noon Contributed Presentations 3/San Gabriel A-C

Chair: Ahmed Sameh, University of Illinois, U-C

10:40/170/A3

Accelerating with Rank-One Updates Timo Eirola and Olavi Nevanlinna

Helsinki University of Technology, Finland

A Parallel, Hybrid Algorithm for the Generalized Eigenproblem

Shing Ma, Merrell Patrick and Daniel Szyld Duke University

11:20/143/A4

Multiprocessor Algorithm for the Symmetric Generalized Eigenvalue Problem

Bill Harrod, Ahmed Sameh and Mark Schaefer University of Illinois, Urbana-Champaign

11:40/182/A4

Parallel QR Factorization on a Hypercube Multiprocessor

Eleanor Chu, University of Waterloo, Alan George, University of Tennessee and Oak Ridge National Laboratory

Tuesday, December 1, 10:40 AM-12:00 Noon Contributed Presentations 4/Santa Barbara A-C

Chair: G.W. Hedstrom, Lawrence Livermore National Laboratory

10:40/14/A4

Shock Calculations on Multi Processors Bjorn Sjogreen

Uppsala University, Sweden

11:00/126/A5

Parallel Processing of a Domain Decomposition Method

R. C. Y. Chin, G. W. Hedstrom, Lawrence Livermore National Laboratory, J. S. Scroggs, University of Illinois, Urbana-Champaign and D. C. Sorensen, Argonne National Laboratory

Moving Point and Particle Methods for **Convection-Diffusion Problems** Michael Rees

Numerical Analysis Group, United Kingdom

11:40/177/A5
Random-Walk Simulation of Diffusion-Reaction-Convection

D. J. Hebert

University of Pittsburgh

12:00 Noon Lunch

Tuesday, December 1/PM

1:30 PM/San Diego Room Invited Presentations 3 and 4 Chair: Burton Smith Supercomputing Research Center

1:30

Neural Computing Michael A. Arbib

University of Southern California, Los Angeles

Optical Computing Systems W. Thomas Cathey University of Colorado, Boulder

3:00 PM/Pasadena Exhibit Hall Coffee

Tuesday, December 1, 3:40-5:00 PM Contributed Presentations 5/San Diego Room

Chair: Edward K. Blum, University of Southern California

3:40/3/A5

Parallel Elliptic Solvers

E. Gallopoulos and Y. Saad

University of Illinois, Urbana-Champaign

4:00/17/A6

The Parallel Waveform Relaxation Multigrid Method

Vandewalle Stefan

Katholieke Universiteit Leuven, Belgium

A Pipelined Block QR Algorithm for a Ring of **Vector Processors**

Christian H. Bischof Cornell University

4:40/32/A6

Conjugate Gradient Preconditioners for Parallel Computers

Gerard Meurant

Departement de Mathematiques Appliques, France

Tuesday, December 1/PM Continued

Tuesday, December 1, 5:00-7:00 PM

Poster Session 1/Pasadena Exhibit Hall/Cash Bar

1/A22

Scheduling of Precedence-Constrained Tasks on Multiprocessors

C. C. Price, Stephen F. Austin State University and M. A. Salama, California Institute of Technology

2/A22

Iterative Suboptimal Strategies for Improving the Transients in Adaptive Control Systems M. De La Sen and M. J. Gonzalez-Gomez, Universidad Del Pais Vasco, Spain

5/A22

A Comparison of Hydrodynamics Multitasking Algorithms

David A. Mandell, Los Alamos National Laboratory

6/A23

Efficient Parallel Algorithms for Controlability and Eigenvalue Assignment Problems

B. N. Datta and Karabi Datta, Northern Illinois University

7/A23

Use of Matrix Determinant of Kronecker Product In Parallel Matrix Computations Karabi Datta, Northern Illinois University

11/A23

Task Assignment in a Multiprocessor System C. Siva Ram Murthy and V. Rajaraman, Indian Institute of Science, India

12/A23

Algorithm and Performance Notes Concerning Block LU Factorization

Jim Armstrong, CONVEX Computer Corporation

13/A24

Use of Mathematical Switches to Solve Differential Equation Problems

Yi-ling F. Chiang, New Jersey Institute of Technology

16/A24

Block Truncated-Newton Methods for Unconstrained Minimization

Stephen G. Nash, George Mason University

22 /A24

Calculating Good Initial Guesses for Iterative Solutions of Parabolic Partial Differential Equations in Parallel

A. Louise Perkins, University of California at Davis in Livermore

23/A25

Systematic Method of Condition/Event Nets Design

Boleslaw Mikolajczak, Southeastern Massachusetts University

26/A25

Neural Nets and Sparse Matrices

Horst D. Simon, Boeing Computer Services

27/A25

Large Memory Applications on the SCS-40 R. E. Anderson, R. G. Grimes, A. C. Mong and

H. D. Simon, Boeing Computer Services

29/A25

A Parallel Matrix Inversion Algorithm for Triangular Toeplitz Systems

You Zhaoyong and Li Lei, Xian Jiaotong University, China

30/A26

Two Supernodal Implementations of General Sparse Factorization for Vector Computers C. Cleve Ashcraft, John G. Lewis and Barry W. Peyton

Boeing Computer Services

33/A26

Recursive Doubling Algorithm for Solution of Tridiagonal Systems on Hypercube Multiprocessors

Omer Egecioglu, Alan J. Laub and Cetin K. Koc, University of California, Santa Barbara

37/A26

Matrix Processing Computers — The Logical Step into the Future

David E. T. Foulser, Saxpy Computer Corporation

39/A26

Parallel Implementation of a Neumann-Dirichlet Preconditioner in Domain Decomposition

Wlodek Proskurowski and Majid Haghoo, University of Southern California

40/A27

The Network Emulation Problem

James Abello, University of California, Santa Barbara, Donna J. Brown, University of Illinois, Michael R. Fellows, University of New Mexico and Michael A. Langston, Washington State University

44/A2

Performance Evaluation for Stochastic Dynamic Programming on Vector Multiprocessor

Floyd B. Hanson, Argonne National Laboratory and University of Illinois at Chicago

45/A27

Parallel Runge-Kutta Methods for Ordinary Differential Equations

Syvert P. Norsett and <u>Ivar Lie</u>, University of Trondheim, Norway

47/A27

Scientific Computation in FORTRAN 8X Alan Wilson, Active Memory Technology

50/A28

Simulation of Larger Dimensions on the Intel iPSC Hypercube

Johnny Petersen and Jan Lindheim, Chr. Michelsen Institute, Norway

51/A28

Parallel implementation of a New Linear Matrix Solver

Lisette de Pillis, Johnny Petersen, Chr. Michelsen Institute and John de Pillis, Bergen Scientific Centre, Norway

53/A28

Parallel Implementation of Fourier Pseudospectral Methods on a Loosely Coupled Array of Processors

Zaphiris Christidis and Vijay Sonnad, IBM Corporation

54/A28

Optimal Algorithms for Parallel Matrix Multiplication on Vector Computer

You Zhaoyong and Li Lei, Xian Jiaotong University, China 55/A29

Exploiting Multilevel Parallelism in PDE Software on the CRAY X-MP

Michael Bieterman, Boeing Computer Services

57/420

An Extension of Ryckaer's Algorithm Sheng-Bai Zhu, Texas Tech University

58/A29

Solving Very Large Dense Systems of Linear Equations

Michael Pernice, National Center for Atmospheric Research

59/A29

Low-Accuracy, High Speed Methods in the Numerical Solution of Ordinary Differential Equations

Chris R. Coray, Utah State University and Jean Michel Favre, Tektronix Inc.

60/A29

Application of Supercomputers in MD Study of Liquid Phase Isomerizations

Sheng-Bai Zhu, Texas Tech University

61/A30

Fortran 9y: Extensions to FORTRAN for Supporting Network Connected Computers

D. E. Stevenson, Clemson University

62/A30

Parallel Algorithms for Market Equilibrium Problems

Ana Nagurney, University of Massachusetts

63/A30

Concurrent I/O System for the Hypercube Mulliprocessor

Andrew Witkowski, California Institute of Technology

193/A30

A Pseudo-Spectral Method Implemented on a SIMD Machine

J. V. W. Reynders, Naval Research Laboratory, and Princeton University

1 -1 /431

Parallel Multigrid Algorithms in Computational Fluid Dynamics

Eric Barszcz, NASA Ames Research Center; Tony F. Chan, University of California, Los Angeles, and RIACS NASA Ames Research Center; Dennis Jespersen, NASA Ames Research Center; and Ray S. Tuminaro, Stanford University and RIACS NASA Ames Research Center

L-2/A31

Matrix Decomposition on Message Passing Architectures

M. Cosnard and B. Tourancheau, Institut National Polytechnique de Grenoble, France

L-3/A3

A Linear Processor Network for the Knapsack Problem

M. Cosnard and A. G. Ferreira, Institut National Polytechnique de Grenoble, France

7:30 PM/California Foyer Registration Closes

7:30 PM/Pasadena Exhibit Hall

Beef and Beer Party \$15.00

Wednesday, December 2/AM

8:30 AM/San Diego Room Invited Presentations 5 and 6 Chair: Joseph Oliger Stanford University

8:30

Automatic Decomposition of Fortran Programs for Execution on Multiprocessors Ken Kennedy

Rice University

9:15

Schwarz Splittings and Template Operators Wei Pai Tang

Stanford University

10:00 AM/Pasadena Exhibit Hall Coffee

10:40 AM CONCURRENT SESSIONS

Wednesday, December 2, 10:40 AM-12:00 Noon Contributed Presentations 6/Santa Barbara A-C

Chair: A. Louise Perkins, University of California, Davis

10:40/129/A6

Concurrent Management of Priority Queues for Adaptive Algorithms

John Kapenga and Elise deDoncker-Kapenga Western Michigan University

11:00/83/A7

Programming Abstractions for Run-Time Partitioning of Scientific Continuum Calculations Running on Multiprocessors Scott B. Baden

University of California, Berkeley

11:20/64/A7

Implementation of the Acceptance-Rejection Method on Parallel Processors: A Case Study in Scheduling

William Celmaster Bolt Beranek and Newman

11:40/74/A7

On the Placement of Parallel Processes Michael R. Leuze and Stephen R. Schach

Vanderbilt University

Wednesday, December 2, 10:40 AM-12:00 Noon Contributed Presentations 7/San Gabriel A-C

Chair: Gregory A. Darmohray, Lawrence Livermore National Laboratory

10:40/104/A7

Mapping Large Scale Computational Problems on a Highly Parallel SIMD Computer

H. M. Liddell, Queen Mary College D. Parkinson, University of London, United Kingdom

11:00/24/A8

Multiprogramming and the Performance of Parallel Programs

Muhammad S. Benten and Harry F. Jordan University of Colorado, Boulder

11:20/122/A8

Using Mathematical Modeling to Aid in Parallel Algorithm Development

Elizabeth Eskow and Robert B. Schnabel University of Colorado, Boulder

11:40/4/A8

The Cerberus Multiprocessor Simulator
Eugene D. Brooks, III, Timothy S. Axelrod and
Gregory A. Darmohray

Lawrence Livermore National Laboratory

Wednesday, December 2, 10:40 AM-12:00 Noon Contributed Presentations 8/San Diego Room

Chair: Youcef Saad, University of Illinois, U-C

10:40/109/A8

A New Parallel Algorithm for Linear Triangular Systems

Avi Lin and He Zhang Temple University

11:00/142/A9

A Parallel Algorithm for the Singular Value Decomposition of Rectangular Matrices William D. Shoaff and David C. Chan

Florida Institute of Technology

11:20/125/A9

A Parallel Algorithm for the Singular Value Decomposition

E. R. Jessup, Yale University and D. C. Sorensen, Argonne National Laboratory

11:40/89/A9

A Multiprocessor Scheme for Singular Value Decomposition

Michael Berry and Ahmed Sameh University of Illinois, Urbana-Champaign

Wednesday, December 2, 10:40 AM-12:00 Noon Contributed Presentations 9/Sacramento Room

Chair: Jeffrey Scroggs, University of Illinois, Urbana

10:40/35/A9

A Parallel Homotopy Method for Solving a System of Polynomials

Lionel M. Ni, Shui-Nee Chow and Yun-Qiu Shen Michigan State University

11:00/94/A10

The Granularity of Homotopy Algorithms for Polynomial Systems of Equations

L. T. Watson, J. Harimoto and J. P. Bixler Virginia Polytechnic Institute and State University

11:20/167/A10

A Parallel Nonlinear Integer Programming Algorithm Based on Branch and Bound and Simulated Annealing

Ken S. Bosworth and G. S. Stiles Utah State University

11:40/131/A10

Recursive Binary Partitions

George Cybenko, Tufts University and Tom Allen, Alphatech, Inc.

12:00 Noon Lunch

Wednesday, December 2/PM

1:30 PM/San Diego Room Invited Presentations 7 and 8 Chair: Jack Dongarra Argonne National Laboratory

1:30

Parallel Systems: Performance Modeling and Numerical Algorithms

Daniel A. Reed

University of Illinois, Urbana

2:15

Asymptotics and Domain Decomposition in Parallel Computing

Raymond C. Y. Chin

Lawrence Livermore National Laboratory, and National Science Foundation

3:00 PM/Pasadena Exhibit Hall Coffee

Wednesday, December 2, 3:40 – 5:00 PM Contributed Presentations 10/San Diego Room

Chair: Peter Linz, University of California, Davis

3:40/49/A11

Parallel Architectures for Optical Computing

Kai Hwang and Ahmed Louri
University of Southern California

4:00/76/A11

Synchronization of Nonhomogeneous Parallel Computations

Dan C. Marinescu and John R. Rice Purdue University

4:20/120/A11

Intrinsically Parallel Algorithms

Oliver A. McBryan

University of Colorado, Boulder

4:40/124/A11

A Software Tool for Building Supercomputer Applications

Dennis Gannon, Daya Atapattu, Mann Ho Lee and Bruce Shei

Indiana University, Bloomington

Special Notice to Contributed Presentation Authors and Chairmen of Contributed Presentation Sessions:

Twenty minutes are allowed for each contributed presentation. Presenters are requested to spend a maximum of 17 minutes for their presentation, and 3 minutes for questions and answers.

Please note:

For presentations with more than one author, an underlineation is used to denote the author who will present the paper.

Wednesday, December 2/PM

Wednesday, December 2, 5:00-7:00 PM
Poster Session 2/Pasadena Exhibit Hall/Cash Bar

65/A32

Fast Fourier Transforms on the BBN Butterfly Parallel Processor

William Celmaster, Bolt Baranek and Newman

66/A32

Scientific Computer Performance Evaluation: A Guide for the Perplexed

M. Edward Borasky, Floating Point Systems, Inc.

67/A32

Fast Banded Matrix Multiplication on Parallel Computers

Richard Bertram and Jerry F. Magnan, Florida State University

68/A32

Parallel and Vector Processing of the Incremental Dynamic Programming Algorithm

Daan Sandee, Control Data Corp. and Dennis Morrow, and Mohan Ramamurthy, Supercomputer Computations Research Institute

69/A33

Simulation of Transonic Flow on a Concurrent Computer

Mark E. Bassett and Christopher J. Catherasoo, AMETEK Computer Research Division

71/A33

Distributed Orthogonal Factorization: Givens and Householder Algorithms

Alex Pothen and Padma Raghavan, The Pennsylvania State University

72/A33

A Constraint Algorithm for Maintaining Rigid Bonds in Molecular Dynamics Simulations of Large Molecules

S. G. Lambrakos, E. S. Oran, J. P. Boris, and M. Nagumo

Naval Research Laboratory

75/A33

Pattern Recognition by Neural Networks on Hypercube

Alex W. Ho and Wojtek Furmanski, California Institute of Technology

77/A34

A Parallel Tridiagonal Equations Solver for the Butterfly Multiprocessor

Swarn P. Kumar, Boeing Computer Services

78/A34

Electromagnetic Scattering Analysis on the Hypercube

Jean E. Patterson, Ruel H. Calalo, Farzin Manshadi, Paulett C. Liewer, William A. Imbriale and James R. Lyons Jet Propulsion Laboratory

79/A34

A Performance Comparison of Three Mini Supercomputers: The Alliant FX/8, The SCS-40, and The Convex C-1

Harvey J. Wasserman, Margaret L. Simmons, and Olaf M. Lubeck, Los Alamos National Laboratory

80/A34

Multitasking a Sparse Matrix Code on the Alliant FX/8

lain S. Duff, Harwell Laboratory, England

81/A35

Parallel Triangular Matrix Inversion

David C. Fisher, Harvey Mudd College and Linda Sattler, The Claremont Graduate School

82/A35

Parallel Processing Bibliography (in Machine-Readable Form)

E. N. Miya, NASA Ames Research Center

85/A35

Analysis of Vectorisable Preconditioners for 3D Problems

Victor Eijkhout, University of Nijmegen, The Netherlands

86/A35

Decomposition and Solution of Transient Finite Element Problems on Hypercubes

Gregory A. Lyzenga, Arthur Raefsky, California Institute of Technology and Bahram Nour-Omid, Lockheed Palo Alto Research Laboratory

87/A36

Serial-to-Parallel Conversion of Real Time Continuous System Simulations

Alan M. Baum and Donald J. McMillan, General Motors Research Laboratories

88/A36

Design of a Real Time, Interactive Simulation Computer with a Parallel Processing Organization

Yukoh Kobayashi, California Research and Development

92/A36

Parallel Subspace Iteration for Solving the Generalized Eigenvalue Problem

<u>Umesh Mahajan</u> and Merrell Patrick, Duke University

93/A36

Performance of the NCUBE Hypercube Geoffrey C. Fox, <u>David W. Walker</u>, California Institute of Technology and Gary R. Montry, Sandia National Laboratories

95/A36

Convergence of Certain Parallel Source Iteration Schemes in Particle Transport

Paul Nelson, Air Force Weapons Laboratory, and Texas A&M University, C. P. Katti, Jawaharlal Nehru University, India and Beny Neta, Naval Postgraduate School

96/A37

Study of Program Termination Using Z-Transform Theory

P. A. Venkatachalam, Asian Institute of Technology, Thailand and S. Arumugam, Computer Centre, GCT, India

97/A37

Parallel Distributed Processing Models for Natural Language Inference

Peter G. Tripodes, Los Angeles City College

100/A37

A Methodology for Dynamic Scheduling of Computation Tasks in Parallel Processing Systems

Ed P. Andert Jr., California State University, Fullerton

101/A37

Analysis of Fortran and C Performance on the FPS T Series Parallel Supercomputer Thomas E. Bauer and Ian A. Taylor, Floating Point Systems, Inc. 102/A37

Application of Book Thickness of Graphs to Parallel Computing

Paul C. Kainen, The Analytic Sciences Corporation

Parallel Optimization of Large-Scale Linear and Nonlinear Networks

Robert R. Meyer, University of Wisconsin-Madison 110/A38

Pipeline SOR: A Parallel Iterative Method for Shared Memory Machines

John P. Bonomo and Wayne R. Dyksen, Purdue University

111/A38

Optimization Strategies for Pattern Selection Problems and Their Implementation on the BBN Butterfly Parallel Processor

William Jeffrey, Robert Rosner, Robert Simon, Harvard-Smithsonian Center for Astrophysics, William Celmaster and Eric Tenenbaum, Bolt Beranek and Newman ACI

113/43

Efficient Implementation of Parallel 3-D Seismic Algorithms on the CRAY X-MP

Mickey Edwards, Christopher C. Hsiung, Cray Research Inc., Dan D. Kosloff, Tel Aviv University, Israel and Moshe Reshef, Cray Research, Inc.

114/A39

Sparse Matrix Solution for Circuit Simulation on Multiprocessors

P. Sadayappan, Ohio State University and V. Visvanathan, AT&T Bell Laboratories

115/A39

Reliable Parallel Eliptic PDE Solution
Bruce M. McMillin and Lionel M. Ni, Michigan
State University

116/A39

Efficient Matrix Multiplication on Hypercube Multiprocessors

<u>Chung-Ta King</u> and Lionel M. Ni, Michigan State University

117/A39

Parallel Simulated Annealing Optimization of Distributed Database Networks

G. S. Stiles, Allison Lee, Utah State University and Philippe Julien-Laferriere, Xerox Corporation, France

118/A40

Parallel Branch-and-Bound Algorithms on Hypercube Multiprocessors

Tarek S. Abdel-Rahman and <u>Trevor N. Mudge</u>, University of Michigan

119/A40

Parallel Programming Without Tears Richard Foster, Chris Thomson and Dan Wilson, Myrias Research Corporation

-4/A4(

Multiplying Arbitrarily Shaped Matrices on a Data Parallel Computer

S. Lennart Johnsson, Alan Ruttenberg, Alex Vasilevsky, and Jill Mesirov, Thinking Machines Corporation

L-5/A40

Computing Fast Fourier Transforms on a Data Parallel Computer

S. Lennart Johnsson, Alan Ruttenberg, Alex Vasilevsky, and Guy Blelloch, Thinking Machines Corporation L-7/A41

A Product Integration Method for Vortex **Dynamics**

Michael Shelley and Moshe Israeli, Princeton University

L-10/A41

A Parallel Algorithm for Matrix Inversion Jerry F. Magnan, Florida State University

7:00 PM/Pasadena Exhibit Hall Meeting of the SIAM Activity Group on Supercomputing/Cash Bar

Thursday, December 3/AM

8:30 AM/San Diego Room Invited Presentations 9 and 10 Chair: Robert G. Voigt **ICASE**

A Fast Multipole Algorithm for Large-Scale **Particle Simulations**

Leslie Greengard Yale University

Lattice Gas Methods for Solving Partial Differential Equations

Gary Doolen

Los Alamos National Laboratory

10:00 AM/Pasadena Exhibit Hall Coffee

10:40 AM **CONCURRENT SESSIONS**

Thursday, December 3, 10:40 AM-12:00 Noon Contributed Presentations 11/Santa Barbara

Chair: Milo Dorr, Lawrence Livermore National Laboratory

10:40/15/A12

Large Scale FE Parallel Nonlinear Computations on High Performance Architectures Charbel Farhat and Luis Crivelli

University of Colorado, Boulder

11:00/34/A12

A Family of Concurrent Algorithms for the **Solution of Transient Finite Element Equations**

B. Nour-Omid, Lockheed Palo Alto Research Laboratory

M. Ortiz, Brown University

A Parallel Algorithm for Rapid Computation of Transient Fields

Robert J. Krueger Iowa State University

11:40/169/A13

Analysis of a Parallelized Ellliptic Solver for **Reacting Flows**

David E. Keyes and Mitchell D. Smooke Yale University

Thursday, December 3, 10:40 AM-12:00 Noon Contributed Presentations 12/San Diego Room

Chair: Eugene Brooks III, Lawrence Livermore National Laboratory

10:40/196/A13

Performance Analyses of the Nested-Dissection Algorithms for MIMD Multiprocessor **Systems**

Vijay K. Naik and Merrell L. Patrick

Institute for Computer Applications in Science and Engineering, NASA Langley Research Center and

11:00/190/A13

A Set of Level 3 Basic Linear Algebra Subprograms

Jack Dongarra, Argonne National Laboratory Jeremy Du Croz, Numerical Algorithms Group Ltd., lain Duff, Harwell Laboratory, Sven Hammarling, Numerical Algorithms Group Ltd.

11:20/56/A14

Gaussian Techniques on Shared Memory Multiprocessor Computers

Gregory A. Darmohray and Eugene D. Brooks, Ill Lawrence Livermore National Laboratory

LU Factorization on Distributed-Memory Multiprocessors

George A. Geist and Charles H. Romine Oak Ridge National Laboratory

Thursday, December 3, 10:40 AM-12:00 Noon Contributed Presentations 13/San Gabriel A-C

Chair: Raymond C. Y. Chin, Lawrence Livermore National Laboratory, and National Science Foundation

10:40/36/A14

Semi-Analytical Shape Functions in Parallel **Computing Environment**

J. A. Puckett and R. J. Schmidt University of Wyoming

11:00/145/A14

The P-S Shape Functions in a Parallel **Computing Environment**

J. A. Puckett and R. J. Schmidt University of Wyoming

Invariant imbedding and the Method of Lines for Multiprocessor Computers

Richard C. Allen, Jr., Lorraine S. Baca and David E. Womble

Sandia National Laboratories

11:40/128/A15

A Portable Parallel Algorithm for Multivariate **Numerical Integration and Its Performance Analysis**

E. deDoncker and John Kapenga Western Michigan University

Thursday, December 3, 10:40 AM-12:00 Noon Contributed Presentations 14/Sacramento Room

Chair: Richard Hickman, Lawrence Livermore **National Laboratory**

Preconditioned Conjugate Gradient Methods for General Sparse Matrices on Shared **Memory Machines**

Edward Anderson and Youcef Saad University of Illinois, Urbana-Champaign 11:00/52/A15

Parallel Multi-Level Finite Element Method with Hierarchical Basis Functions

G. Brussino, R. Herbin and V. Sonnad **IBM Corporation**

11:20/48/A16

A Double Augmented Tree Architecture for **Parallel Execution of Multigrid Algorithms**

H. C. Wang and Kai Hwang University of Southern California

11:40/38/A16

Implementation of a Parallel Multigrid Method on a Loosely Coupled Array of **Processors**

Raphaele Herbin, IBM Corporation Stephane Gerbi, E.N.T.P.E., France Vijay Sonnad, IBM Corporation

12:00 Noon Lunch

Thursday, December 3/PM

1:30 PM/San Diego Room Invited Presentations 11 and 12 Chair: Michael R. Raugh NASA Ames Research Center

Strategies for the Mapping and Scheduling of Scientific Computations

Joel Saltz Yale University

Graph Theory Plus Group Theory Equals Answers to Communication Problems Vance Faber

Los Alamos National Laboratory

3:00 PM/Pasadena Exhibit Hall Coffee

Thursday, December 3, 3:40-5:00 PM Contributed Presentations 15/San Diego Room

Chair: Włodek Proskurowski, University of Southern California

3:40/127/A16

SCHEDULE: An Environment for Developing Transportable Explicitly Parallel Codes in Fortran

J. J. Dongarra and D. C. Sorensen Argonne National Laboratory

4:00/140/A17

Domain Decomposition Methods for Three-Dimensional Elliptic Problems O. Axelsson

University of Nijmegen, The Netherlands

4:20/146/A17

Fast Poisson Solver on Irregular Regions by Boundary Integral-based Domain Decomposition

Daeshik Lee

University of Illinois, Urbana-Champaign

4:40/178/A17

A High-Performance FFT Algorithm for **Vector Supercomputers**

David H. Bailey NASA Ames Research Center

Thursday, December 3/PM

Thursday, December 3, 5:00-7:00 PM
Poster Session 3/Pasadena Exhibit Hall/Cash Bar

121/A41

Obtaining Nonuniform Systolic Designs Using Modified Mapping Techniques

Keith R. Allen and Roy P. Pargas, Clemson University

130/A42

A Parallel Processing Pre-Processor for the Fortran Language

Bruce W. Suter, Robert M. Hyatt and Joydeep Roy, University of Alabama at Birmingham

132/A42

Inhomogenous Parallel SupercomputersAnthony E. Terrano, Rutgers University

133/A42

Fault-Tolerant Mapping of Large Scale Resource Allocation Algorithms onto Concurrent Processors

Rong T. Lee, Krishna R. Pattipati and Peter B. Luh, University of Connecticut

135/A42

Research Issues in 1000-Processor Linear Algebra Algorithms in Finite Element Analysis Robert E. Benner and Gary R. Montry, Sandia

National Laboratories

137/A43

Implementation of Iterative and Direct Solvers on The UNIX CRAY X-MP24 Michael A. Driscoll and Elissa C. Pariser, AT&T Bell Laboratories

138/A43

A Multiprocessor Model for a System Evaluation Tool

Kathleen M. Nichols, AT&T Bell Laboratories

139/A43

A Case Study in Scientific Parallel Processing: Prefix-Based Computation
Binay Sugla, AT&T Bell Laboratories

141/A43

The Effects of Parallel Processing on Matrix Computations

Laurie Melany Darmody, Laurel, MD

144/A44

The Behavior of Conjugate Gradient Methods on a Multi-Vector Processor with a Hierarchical Memory

William Jalby, <u>Ulrike Meier</u> and Ahmed Sameh, University of Illinois-Urbana-Champaign

147/A44

The Impact of Hierarchical Memory Systems on Linear Algebra Algorithm Design

Kyle Gallivan, William Jalby, Ulrike Meier and Ahmed Sameh, University of Illinois, Urbana-Champaign

148/A44

An Analysis of a Multivector Processor Implementation of Cyclic Reduction

Kyle Gallivan, Youcef Saad, University of Illinois at Urbana-Champaign and William Jolly, INRIA, France

149/A44

An Out-Of-Core Algorithm for Solving Dense Sets of Linear Equations on the FPS-164 Mark S. Woodyard, Martin Marietta Orlando Aerospace and Hassan N. Srinidhi, University of Central Florida

152/A45

Linear Permutations on a Multistage Interconnection Network

Anujan Varma, IBM Thomas J. Watson Research Center

156/A45

Solution of Elliptic Systems on Highly Parallel Computers

Paul O. Frederickson, Los Alamos National Laboratory

158/A45

Minimum Spanning Tree on The Hmesh Architecture

R. V. Boppana and C. S. Raghavendra, University of Southern California

159/A45

Comparison of Methods of Solving Symmetric Positive Definite Linear Systems on a Hypercube Concurrent Computer Lois Mansfield, University of Virginia

161/A45

A Simulation Tool for Optimizing Production

Steve Howard, Scientific Computer Systems Corporation

163/A48

Task Migration in Partitionable Parallel Processing Systems

Thomas Schwederski, Thomas L. Casavant, Purdue University and <u>Howard Jay Siegel</u>, Supercomputing Research Center

164/A46

Early Experience With the PASM System Prototype

Samuel A. Fineberg, Thomas L. Casavant and Thomas Schwederski, Purdue University

165/A46

Iterative Algorithms in a Data Driven Environment

P. Evripidou and J. L. Gaudiot, University of Southern California

166/A46

Computational and Reliability Measures of Homogeneous Multiprocessor Systems

Walid Najjar and Jean-Luc Gaudiot, University of Southern California

168/A46

Adaptive Grids of Multivariate, NTH-Order Calculi

Peter Nwoye O. Mbaeyi, West Germany

171/A47

Compiler-Drive Cache Management
Chi-Hung Chi and Henry Dietz, Purdue University

172/A47

CRegs: A Hardware Solution to the Alias Problem

Henry Dietz and Chi-Hung Chi, Purdue University

173/A47

Barrier MIMD: Beyond VLIW

Henry Dietz and Thomas Schwederski, Purdue University

174/A47

Efficient Iterative Solution of Elliptic PDE's Fabian Waleffe, Massachusetts Institute of Technology and Thinking Machines Corporation

176/A47

A MIMD Eigenvalue Algorithm for a Parallel Transputer-Microcomputer System Ina Vincent and Peter W. Aitchison, University of

179/A48

Manitoba, Canada

Parallel Conjugate Gradients on a 1024-Node Hypercube

John L. Gustafson and Robert E. Benner, Sandia National Laboratories

183/A48

A Parallel Algorithm for Computing the Generalized Singular Value Decomposition (GSVD)

Zhaojun Bai, University of Maryland

185/449

Finding Public Domain, Mathematical Software

Greg Astfalk, AT&T, Jack Dongarra, Argonne National Laboratory and Eric Grosse, AT&T Bell Laboratories

186/A49

The Lawrence Livermore Loops with a Vectorizing C Compiler

Greg Astfalk, AT&T and <u>Dale Lancaster</u>, Convex Computer Corp.

194/A49

A Graphical Programming Language for the Navier-Stokes Computer

Sherryl Tomboulian, Thomas W. Crockett and David Middleton, NASA Langley Research Center

195/A50

A Comparison of Different Parallel Programming Languages Using Choleski's Method of Multi-Processor Computer

Mark Jones, Duke University and Robert Voigt, NASA Langley Research Center

191/A51

Generalized SOR — New Spectra from Old John de Pillis, Bergen Science Center, Norway

-8/450

Performance of Some CFD Codes on the Alliant FX/8

Walid Abu-Sufah and Amjad Daoud, Virginia Polytechnic Institute and State University

1-9/450

Experiences with the Use of the Cray X-MP/48 Performance Monitoring Tools Walid Abu-Sufah and Amjad Daoud, Virginia Polytechnic Institute and State University

L-11/A50

Combined Vector and Parallel Processing on the IBM 3090

John Tesch and Ken Jacobs, IBM, Dallas

L-12/A51

Applications of Level 2 BLAS in the NAG Library

Richard Brankin, Jeremy Du Croz, and Peter Mayes, The Numerical Algorithms Group Limited, Oxford, England

Friday, December 4/AM

8:30 AM/San Diego Room Invited Presentations 13 and 14 Chair: John Van Rosendale Argonne National Laboratory

8:30

Parallel Processing Research at IBM Tilak Agerwala

IBM T. J. Watson Research Center

9:15

Programming Environments for Multiprocessor Architectures

Piyush Mehrotra Purdue University

10:00 AM/Pasadena Exhibit Hall Coffee

10:40 AM CONCURRENT SESSIONS

Friday, December 4, 10:40 AM – 12:00 Noon Contributed Presentations 16/San Gabriel A – C

Chair: Merrell Patrick, Duke University

10:40/153/A17

Parallel Vision Algorithms: An Approach Sharat Chandran and Larry S. Davis University of Maryland

11:00/154/A18

Performance of an Ocean Circulation Model on LCAP

<u>Hsiao-Ming Hsu</u>, Woods Hole Oceanographic Institution

Jih-Kwon Peir, IBM T. J. Watson Research Center, **Dale B. Haidvogel**, The Johns Hopkins University

11:20/28/A18

Ising Spin on a Shared Memory Machine: Computational Experience

James L. Blue and Francis Sullivan National Bureau of Standards

11:40/189/A18

Parallel Iterative Algorithms To Solve The Discrete LQR Optimal Control Problem With Hard Control Bounds

Gerard G. L Meyer, The John Hopkins University Louis J. Podrazik, Allied Signal, Inc.

Friday, December 4, 10:40 AM-12:00 Noon Contributed Presentations 17/Sacramento Room

Chair: Mark K. Seager, Lawrence Livermore National Laboratory

10:40/175/A19

A Comparison of the Performance of LINPACK and Vectorized ICCG in Solving a Poisson Equation for a Problem of Two-Phase Flow in a Centrifuge

Eugene L. Poole, NASA Langley Research Center, Jeffery W. Frederick, University of Virginia

11:00/187/A19

On the Solution of Block Tridiagonal Systems of Equations on Multi-Vector Machines — A Comparative Study

Ken Dowers, S. Lakshmivarahan, Sudarshan K. Dhall and Julio Diaz, University of Oklahoma

11:20/112/A19

Linear Algebra at 30 MFLOPS on a 9 MFLOP Machine

Stewart F. Reddaway, ICL, England

11:40/134/A19

Performance of Blocked Gaussian Elimination on Multiprocessors

Elizabeth Jean O'Neil, BBN Advanced Computers, Inc., Henno Allik, BBN Laboratories Inc.

Friday, December 4, 10:40 AM-12:00 Noon Contributed Presentations 18/Santa Barbara A-C

Chair: Scott Baden, University of California Berkeley

10:40/98/A20

Loop Blocking for Parallel Memory Optimiza-

Michael Wolfe, Kuck and Associates, Inc.

11:00/150/A20

Parallel FÓRTRAN: Why You Can't. How You Can.

Clifford Arnold, ETA Systems, Inc.

11:20/180/A20

Performance Based Distributed Programming Phillip Q. Hwang, Naval Surface Weapons Center

11:40/20/A20

A Programming Aide for Message-passing Systems

Min-You Wu and Daniel D. Gajski University of California, Irvine Friday, December 4, 10:40 AM – 12:00 Noon Contributed Presentations 19/San Diego Room

Chair: William D. Gropp, Yale University

10:40/10/A21

A Gray-Code Scheme for Local Uniform Mesh Refinement on Hypercubes William D. Gropp and Ilse C. F. Ipsen Yale University

11:00/43/A21

Iteration with Mesh Refinement for Systems of ODE's

Olavi Nevanlinna

Helsinki University of Technology, Finland

11:20/108/A21

Parallelization of Adaptive Grid Domain Mappings

Calvin J. Ribbens, Virginia Polytechnic Institute and State University

11:40/25/A21

Automated Mesh Decomposition and Concurrent Finite Element Analysis for Hypercube Multiprocessor Computers James G. Malone

General Motors Research Laboratories

1:00 PM Conference Adjourns TUESDAY, DECEMBER 1 10:40 AM - 12:00 Noon San Diego Room

Contributed Presentations 1

#46/10:40 AM
The Design and Implementation of Parallel
Algorithms in Ada

The features of Ada that support parallel computing in mathematical software are covered and their advantages emphasized. The design and implementation of reusable parallel mathematical software components based on Ada's package, generic, data abstraction, and information hiding features are discussed. We also describe how to create an Ada shell that permits the use of existing Fortran software while maintaining the advantages of Ada. Mathematical software examples, primarily linear algebra, are presented that illustrate these Ada features.

Richard F. Sincovec Computer Science Department University of Colorado Colorado Springs, CO 80933-7150

#123/11:00 AM
DINO -- A New Language
for Numerical Computation
on Local Memory Multiprocessors

We describe a new language for parallel processing that is specifically oriented towards writing numerical programs for local memory multiprocessors. DINO consists of extensions to standard C. The key new feature of DINO is the declaration by the user of a problem topology, which can be thought of as a virtual processor topology related to the natural physical or graphical topology of the problem. The topology is used in two additional DINO language extensions, composite procedure calls and distributed variables, that allow the execution of parallel procedures and communication between them without explicit com-

munication statements. Experience using a prototype of DINO written in C++ will be described.

Matthew Rosing Robert B. Schnabel Department of Computer Science Campus Box 430 University of Colorado Boulder, Colorado 80309

#160/ll:20 AM Scientific Parallel Processing with LGDF2

Large-Grain Data Flow was developed for writing portable, parallel scientific programs. The method uses a graphical language for specifying parallelism, combining dataflow-like programming constructs with blocks of ordinary sequential code. Macro expansion automatically generates source code containing efficient parallel scheduling mechanisms for a particular (parallel) processor. original LGDF model was used successfully to implement several medium-sized programs on the Denelcor HEP. An improved version of the programming model (LGDF2) allows more efficient scheduling. Performance results for the Sequent Balance 21000 are given. Also, software tools are described that allow reliable construction, tuning, and debugging of parallel programs.

Robert G. Babb II David C. DiNucci Dept. of Computer Science and Engineering Oregon Graduate Center 19600 NW Von Neumann Dr Beaverton, OR 97006

#157/11:40 AM

"The BF (Boundary-Fitted) Coordinate Transformation Technique of DEQSOL (Differential EQuation SOlver Language) "

DEQSOL is a high-level programming language specially designed to describe PDE problems in a way quite natural for numerical analyses. The DEQSOL translator automatically generates highly vectorizable simulation codes from high-level descriptions.

The BF facility has been developed in addition to the existing FDM and FEM facility. This facility automatically transforms any PDEs and regions defined in physical space into equivalent PDEs and rectangular regions in transformed space. This facility has made DEQSOL applicable to the development of simulators on curved and moving boundary regions. such as the impurity diffusion simulator in semiconductor devices.

Despite this extension, programming productivity is still an order of magnitude higher than FORTRAN programming, and generated codes achieve extremely high vectorization ratios(over 90%) on the Hitachi S-810 vector processor.

Chisato Konno, Michiru Yamabe, Miyuki Saji, Yukio Umetani 8th Dept., Central Research Laboratory Hitachi Ltd. Kokubunji, Tokyo 185, Japan

TUESDAY, DECEMBER 1 10:40 AM - 12:00 Noon Sacramento Room

Contributed Presentations 2

#105/10:40 AM

Communication Complexity of Matrix Multiplication

Suppose a parallel processing machine multiplies two $n \times n$ matrices so only one copy of each input exists at the start of computations. We show that no matter what algorithm is used, it is possible to halve the machine so the amount of data exchanged between the halves is at least $(4n^4 - L^2)/16n^2$ where L is the maximum number of inputs entering one port. This supercedes a similar bound of Savage (Journal of Computer and System Science, 22, 230-242 (1981)). This bound is used to find lower bounds on the time needed to do matrix multiplication on various parallel architectures.

David C. Fisher Department of Mathematics Harvey Mudd College Claremont, CA 91711

#19/11:00 AM

ORDERINGS FOR PARALLEL SPARSE SYMMETRIC-FACTORIZATION

We have developed a new heuristic for obtaining a nested dissection ordering, based on graph bisection algorithms. We use the Kernighan and Lin algorithm, as modified by Fiduccia and Matheyses, to partition the relevant section graphs. We present a new graph model so that the grid partitionings provide nodal separators, rather than the usual edge separators. We have refined the standard algorithm to further reduce the size of the separator set.

Our heuristic produces consistently and sometimes significantly better parallel orderings than the usual modifications of scalar orderings. In particular the depth of the elimination tree is reduced. This is the appropriate measure of time for our target machine, the Connection Machine, and for other parallel machines. The orderings are also comparable in space requirements to the best scalar orderings.

Charles E. Leiserson*
Laboratory for Computer Science
NE43-321
Massachusetts Institute of Technology
Cambridge MA 02139

3.

John G. Lewis^{*}
Boeing Computer Services Company
Mail Stop 7L-21
PO Box 24346
Seattle WA 98214

- * and Thinking Machines Corporation
- * and MIT Center for Advanced Engineering Study

#162/11:20 AM
An Efficient Fixed Size Array for Solving
Large Scale Toeplitz Systems

In this paper systolic implementation and partitioning of Bareiss algorithm for solving an (n+1) by (n+1) Toeplitz system of equations is studied. While systolic solutions to Toeplitz systems has been well studied, the main problem with the known results is that they cannot solve any arbitrary Toeplitz system of equations on a fixed size array of processors. Our systolic implementation solves the Bareiss algorithm on a fixed size linear array. Our

algorithm achieves $O(\frac{n^2}{p})$ computing time and $O(n^2)$ storage on p (p may be fixed) processors. Even without the partitioning, our implementation has superior time performance compared to known results within a constant factor. Furthermore, our implementation achieves a processor utilization of almost 100%. Our solution allows the real time solution of large Toeplitz matrices which occur in most Toeplitz applications.

Afshin Daghi

Systolic Solutions Inc.

V.K. Prasanna Kumar

Department of Electrical Engineering-Systems
University of Southern California
LA, California 90089-0781
Ali Safavi

USC Information Sciences Institute and Computer Science Department

#188/11:40 AM Architecture and Operation of a Systolic Sparse Matrix Engine

This paper presents the architecture and operation of the FEM, a systolic array tuned to sparse matrix operations. It enables high-performance computation of the preconditioned conjugate gradient method for solving the large, sparse, linear systems of equations common to applications such as finite element analysis. Matrix operations are typically very simple and regular and thus ideally suited for systolic computation. The desire was to design a robust system that could handle sparse matrices with a general sparsity structure. Dense matrices are treated as special cases of sparse matrices instead of the other way around which is typical of vector processors. The objective of this work is to provide an efficient systolic architecture for computing the operations of iterative techniques.

Robert J. Dunki-Jacobs, Steven W. Hammond, Robert M. Hardy, and Terry M. Topka GE Corporate Research and Development Center P.O. Box 8 Schenectady, NY 12301 USA

TUESDAY, DECEMBER 1 10:40 AM - 12:00 Noon San Gabriel Rooms A-C

Contributed Presentations 3

#170/10:40 AM Accelerating with Rank-One Updates

Consider iterations for solving Ax = b (An by n nonsingular) which are based on splitting A=M-N. We discuss rank-one updates to improve inv(M) as an approximation to inv(A) during the iteration. The update kills and reduces singular values of Ninv(M) and thus speeds up the convergence. The basic algorithm terminates after at most n sweeps, and if full n sweeps are needed, then inv(A) has been computed. Suitability for parallel processing and the effect of restricted memory length for the updates are also discussed.

Timo Eirola and Olavi Nevanlinna Institute of Mathematics Helsinki University of Technology SF-02150 Espoo, Finland #91/11:00 AM A Parallel, Hybrid Algorithm for the Generalized Eigenproblem

We present a parallel algorithm for computing all eigenvalues, and their corresponding eigenvectors, in a specified interval for the generalized eigenproblem, $Ax = \lambda Bx$, where A and B are real, symmetric and B is positive definite. Eigenvalues are isolated, in parallel, using the Sturm sequence property of leading principal minors of A - μB . Concurrently, eigenvalues and eigenvectors are computed accurately using a combination of inverse and Rayleigh quotient iterations.

Results obtained from implementation of this algorithm on a shared memory MIMD architecture are presented. Factors which affect the efficiency of the algorithm are discussed.

Shing Ma Merrell Patrick Daniel Szyld

Department of Computer Science Duke University Durham, NC 27706

#143/11:20 AM

Multiprocessor Algorithm for the Symmetric Generalized Eigenvalue Problem

A trace minimization algorithm for computing a few of the smallest (or largest) eigenvalues and associated eigenvectors of the generalized eigenvalue problem $Ax = \lambda Bx$ is presented. Here we assume that the matrices A and B are symmetric of order n, with B being positive definite, and that both A and B are so large and sparse that a factorization of either matrix is impractical. In each iteration of this algorithm, first investigated by A. Sameh and J. Wisniewski, we are simultaneously approximating the p desired eigenpairs p << n by minimizing the trace of a $p \times p$ matrix subject to quadratic constraints. This talk presents an improved computational scheme which incorporates a variety of acceleration strategies, as well as demonstrates the suitability of the algorithm on a multiprocessor with two levels of parallelism, such as the Alliant FX/8

Bill Harrod, Ahmed Sameh and Mark Schaefer University of Illinois at Urbana-Champaign Center for Supercomputing R & D 305 Talbot Laboratory 104 South Wright Street Urbana, Illinois 61801-2932 USA

#182/11:40 AM

Parallel OR Factorization on a Hypercube Multiprocessor

We describe a new algorithm for computing an orthogonal decomposition of a rectangular $m \times n$ matrix A on a hypercube multiprocessor. The algorithm uses Givens rotations, and requires the embedding of a two-dimensional grid on the hypercube network. We design a global communication scheme which uses redundant computation to maintain the data proximity, and we employ a mapping strategy for data allocation so that the processor idle time remains constant for a fixed number of processors regardless of the size of a square matrix. We describe how the algorithm is easily generalized to include the case when m > n. The proposed global communication scheme and the data mapping strategy result in reduced computation and communication cost compared to other known results for the same problem. Complexity results and numerical experiments on an Intel iPSC hypercube will be presented.

Eleanor Chu
Department of Computer Science
University of Waterloo
Waterloo, Ontario, Canada

Alan George
Department of Computer Science
University of Tennessee
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and
Oak Ridge National Laboratory
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TUESDAY, DECEMBER 1 10:40 AM - 12:00 Noon Santa Barbara Rooms A-C

Contributed Presentations 4

#14/10:40 AM Shock Calculations on Multi Processors

We present results from computing solutions to the compressible two dimensional Euler equations on a hypercube multiprocessor. We use modern high resolution finite difference schemes to get sharp shock profiles, mainly the Roe scheme and its generalizations to higher order accuracy. We discuss how to decompose the computational domain into different processors and we analyze if we can

increase the computational speed by not doing communication on every time level. Some results from solving the same problems on the Connection Machine will also be presented.

Bjorn Sjogreen
Department of ScientificComputing
Uppsala University
Sturegatan 4B
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SWEDEN

#126/11:00 AM
Parallel Processing of a
Domain Decomposition Method

We present a parallel algorithm for the efficient solution of a time dependent singularly perturbed convection diffusion equation. The method is based upon domain decomposition that is dictated by singular perturbation analysis. A transformation is made to a coordinate system induced by the characteristics of the reduced hyperbolic equation. Asymptotic analysis is used to determine regions where certain reduced equations may be solved in place of the full equation. This reduces the number and size of the domains where the full equation must be solved. Treatment of the internal layer leads naturally to a shock detection technique appropriate for the iterative solution of nonlinear problems. Parallelism is evident at 3 levels. Domain decomposition provides parallelism at the highest level, and within each subdomain we exploit parallel-vector processing. Important features of the method include independent solution of the characteristic curves at the lowest level and low communication requirements between processes devoted to solving in the various domains. Tightly coupled processes are only required in domains where the full equation must be solved. We shall discuss the implementation and some aspects of the performance of this algorithm on existing parallel computers.

R.C.Y. Chin and G.W. Hedstrom Lawrence Livermore National Laboratory P.O. Box 808 Livermore, CA 94550

J.S. Scroggs Center for Supercomputer Research and Development University of Illinois at Urbana-Champaign Urbana, IL 61801

D.C. Sorensen Mathematics and Computer Science Division Argonne National Laboratory 9700 South Cass Avenue Argonne, IL 60439-4844

#18/11:20 AM Moving Point and Particle Methods for Convection-Diffusion Problems

We present several recently proposed numerical methods for the solution of

multi-dimensional convection-diffusion problems. The common feature of all methods is the presence of a set of points which move according to the local fluid velocity field. However, the type of information that is associated with each of these points may be different. In addition, there may be a fixed mesh present, providing even greater variety of numerical method. In particular, we consider a moving point method employing a fixed grid and a grid-free particle method. The suitability of these algorithms to vector processors will be discussed.

Michael Rees Numerical Analysis Group Computing Laboratory 8-11 Keble Road Oxford, United Kingdom

#177/11:40 AM

Random-Walk Simulation of Diffusion-Reaction-Convection

A simple stochastic model for populations of diffusing, interacting, and drifting particles is used to derive simulation algorithms which are closely related to random-walk methods and particle methods for the solution of partial differential equations of reaction-diffusion and fluid flow. These algorithms are then directly realized as simple vector or parallel programs. Examples are given which compare the resulting simulations with some standard numerical solutions of diffusion-reaction and diffusion-convection systems of partial differential equations.

D. J. Hebert Department of Mathematics and Statistics University of Pittsburgh Pittsburgh, Pa. 15260

TUESDAY, DECEMBER 1 3:40 - 5:00 PM San Diego Room

Contributed Presentations 5

#3/3:40 PM Parallel Elliptic Solvers

We consider parallel versions of the standard fast elliptic solvers and discuss their merits and disadvantages. In particular, we present an adaptation of the Block Cyclic Reduction (BCR) algorithm to vector and parallel machines. The main bottleneck of BCR lies in the solution of linear systems whose coefficient matrix is the product of tridiagonal matrices. This bottleneck is handled by expressing the rational function corresponding to the inverse of this product as a sum of elementary fractions. This

leads to parallel solutions of tridiagonal systems. FFT based algorithms are also discussed. We discuss implementations of these algorithms on shared memory and distributed memory machines.

E. Gallopoulos Y. Saad

University of Illinois at Urbana-Champaign Center for Supercomputing R & D 305 Talbot Laboratory 104 South Wright Street Urbana, Illinois 61801-2932 USA

#17/4:00 PM

The Parallel Waveform Relaxation Multigrid Method

A great deal of attention has recently been given to the development of the highly parallel Waveform Relaxation method (WR) for solving very large systems of ordinary differential equations. Attempts to use WR to solve the equations arising from the numerical method of lines have not been very successful due to the slow convergence of the normally used Jacobi, Gauss-Seidel and SOR methods.

In this paper we present a new algorithm that combines the parallel nature of WR with the fast convergence of multigrid to solve time dependent partial differential equations.

We also discuss its implementation on the Caltech Hypercube.

Vandewalle Stefan Katholieke Universiteit Leuven dept computerscience Celestijnenlaan 200A B - 3030 Heverlee, Belgium

#31/4:20 PM

A Pipelined Block QR Algorithm for a Ring of Vector Processors

We present a parallel algorithm for computing the QR factorization of a dense matrix on a ring of vector processors. Regarding the matrix A as a matrix of block columns, we use the WY representation (C. Bischof and C. Van Loan, "The WY Representation for Products of Householder Matrices" SIAM J. Scientific and Statistical Computing, Vol. 8, No. 1, 1987) to bundle Householder updates. This strategy results in an algorithm that is rich

in matrix-matrix multiplication, the operation of choice for pipelined vector units. Here we show how this technique can be used in the coarse-grained parallel setting. We distribute block columns across the ring of processors and stagger the block column updates so that the computation can proceed in a pipelined fashion. We have performed experiments on the LCAP system at IBM Kingston which show that this method is indeed well suited for coarse-grained distributed systems.

Christian H. Bischof Dept. of Computer Science 5133 Upson Hall Cornell University Ithaca, NY 14853

#32/4:40 PM

CONJUGATE GRADIENT PRECONDITIONERS FOR PARALLEL COMPUTERS

In this paper, we show how to derive efficient preconditionners for the conjugate gradient method on parallel computers with a large number of processors. These preconditioners use domain decomposition techniques, the domain being divided either in strips or in boxes. Numerical experiments will show that for a given problem dimension, the increase in the number of iterations as a function of the number of processors, is very small, making these preconditioners very attractive for future parallel computers.

Gérard MEURANT Département de Mathématiques Appliquées CEL-V - B.P. N° 27 94190 - VILLENEUVE St GEORGES FRANCE

WEDNESDAY, DECEMBER 2

10:40 AM - 12:00 Noon Santa Barbara Rooms A-C

Contributed Presentations 6

#129/10:40 AM Concurrent Management of Priority Queues for Adaptive Algorithms

A study of the demands of several types of parallel numerical algorithms, which are synchronized by a task pool, is presented. This is used to evaluate different methods of concurrent priority queue management on various shared memory architectures. Both an analytic model and empirical results from the Denelecor HEP, Sequent Balance, Alliant FX-8 and Encore Multimax are considered.

This work is part of our effort to develop a set of macros, layered over the Argonne macro package, which provide a means of producing portable programs for adaptive partitioning problems. Portable efficient concurrent priority queue management is critical to this endeavor.

John Kapenga Elise deDoncker-Kapenga Department of Computer Science Western Michigan University Kalamazoo, MI 49008-5021

#83/11:00 AM

Programming Abstractions for Run-Time Partitioning of Scientific Continuum Calculations Running on Multiprocessors

I will discuss a set of software abstractions for dynamically partitioning various math-physics calculations on a team of processors. I tried out the abstractions on Anderson's Method of Local Corrections, a type of vortex method for computational fluid dynamics. I ran experiments on 32 processors of the Intel iPSC -- a message-passing hypercube architecture -- and on 4 processors of a Cray X-MP -- a shared-memory vector architecture -- and achieved good parallel speedups of 24 and 3.6, respectively. The abstractions should apply to diverse applications, including finite difference methods, and to diverse architectures without requiring that the application be reprogrammed extensively for each new architecture.

Scott B. Baden
Mathematics Department
Lawrence Berkeley Laboratory
University of California
Berkeley, CA 94720

#64/11:20 AM

Implementation of the Acceptance-Rejection Method on Parallel Processors:

A Case Study in Scheduling

The Acceptance-Rejection method (AR) is often used in connection with the generation of nonuniform random distributions. We examine the problem of generating, via AR, vectors of nonuniform numbers on various kinds of vector and parallel processors. Vector and SIMD implementations are generally quite inefficient

compared to MIMD implementations. Various scheduling models are presented and the theoretical predictions are verified by tests done on a BBN Butterfly Parallel Processor.

William Celmaster

Bolt Beranek and Newman Advanced Computers Inc. 10 Fawcett Street, Cambridge, MA 02238

#74/11:40 AM

On the Placement of Parallel Processes

An important problem of parallel computation is to determine the best placement in a multiprocessor system of each of a set of related processes. An earlier approach to this problem attempted to maximize the number of communicating processes which lie on adjacent processors. We demonstrate that this approach is not suitable for large problems through a comparison of mappings generated by this approach with random mappings and with mappings generated by a heuristic which attempts to minimize the distance from each process to all processes with which it communicates. The various approaches are experimentally compared with respect to a specific problem.

Michael R. Leuze and Stephen R. Schach Computer Science Department Vanderbilt University Box 77 Station B Nashville, TN 37235

WEDNESDAY, DECEMBER 2 10:40 AM - 12:00 Noon San Gabriel Rooms A-C

Contributed Presentations 7

#104/10:40 AM

Mapping Large Scale Computational Problems on a Highly Parallel SIMD Computer

The mapping strategies employed to solve various size computational problems on an SIMD computer with p = 1000 or more processors are described. For 'small' problems, whose size n << p, many problems can be solved simultaneously, and when n is approximately equal to p, straightforward 'direct' mappings can be used. The most interesting, and difficult, situation arises for large problems, n >> p where various techniques have been applied to a wide variety of computational problems. These include sheet mapping, crinkled mapping, linear array, multi-serial, parallel

data transforms, permutations or a combination of mappings. The problem of changing mappings is also considered. The application of these strategies to a number of problems and their implementation on the various DAP systems will be described

Heather M. Liddell
Dennis Parkinson
Department of Computer Science & Statistics
Queen Mary College
University of London
Mile End Road
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El 4NS

#24/11:00 AM

Multiprogramming and the Performance of Parallel Programs

Tight synchronization in parallel programs executed on multiprogrammed multiprocessors may result in catastrophic performance losses as a result of the absence of swapped out processes. Our work introduces a programming methodology that utilizes computational synchronization and avoids tight synchronization in control flow parallel programs. In this methodology, each phase of the computation is assigned a status that can be ready, blocked or completed, and tasks in each computational phase are selfscheduled to ensure computational progress by the available executing processes. Results obtained indicate that this methodology avoids the catastrophic performance losses resulting from the swapping of processes in multiprogrammed multiprocessors.

Muhammad S. Benten Harry F. Jordan Electrical and Computer Engineering Department University of Colorado Boulder, CO 80309-0425

#122/11:20 AM
Using Mathematical Modeling to Aid in
Parallel Algorithm Development

The intent of this presentation is to demonstrate the advantages of using mathematical modeling in conjunction with parallel computer implementation in designing and understanding parallel algorithms. We describe the formation of a model that accurately matched execution times of a rather complex parallel global optimization method on an Intel hypercube. We then discuss the application of this model to detect weakness in the algorithm and to analyze several possible improvements to it. We believe that this example helps to indicate that the use of mathematical

modeling together with computer implementation is useful in making parallel algorithm development more efficient in human and computer time.

Elizabeth Eskow Robert B. Schnabel Department of Computer Science Campus Box 430 University of Colorado Boulder, Colorado 80309

#4/11:40 AM

The Cerberus Multiprocessor Simulator

We describe a simulation facility for scalable shared memory multiprocessors. The processors are a RISC architecture with a minimum of additional instructions added to support synchronization. The functional units of the processors, including access to the shared memory, are fully pipelined. The multiprocessor is simulated at a very detailed level, with pipeline delays and conflicts in the packet switched shared memory server being accurately accounted for. Applications are written in C. The simulator can be used to examine algorithm performance as the number of processors in the multiprocessor is scaled, and accurately measures the effects of pipeline delays and conflicts in the shared memory subsystem.

Eugene D. Brooks, III, Timothy S. Axelrod, Gregory A. Darmohray

Computational Physics Division University of California, Lawrence Livermore National Laboratory P. O. Box 808 Livermore, California 94550

WEDNESDAY, DECEMBER 2

10:40 AM - 12:00 Noon San Diego Room

Contributed Presentations 8

#109/10:40 AM A New Parallel Algorithm for Linear Triangular Systems

A parallel algorithm for solving linear triangular system is presented and analyzed. This algorithm has a small start up phase which is different from the rest of the parallel phases so that the first processor will be able to accumulate enough results before shipping them to the all other processors. In the rest of the parallel phases, the first processor makes a substitution of the results of the rest of the processors from the previous phase and then makes an elemination of more variables. In the same time the rest of the processors substitute the results of the first processor from the previous

phase. The most important features of this algorithm are: (1) It works well for any number of processors. (2) It achieves the optimal speed up, optimal efficiency and very low communication complexity. (3) It can be used in both distributed parallel computing environment and tightly coupled parallel computing system. (4) This algorithm can be mapped easily onto any parallel architecture without any programming difficulties.

Avi Lin He Zhang Department of Mathematics Temple University Philadelphia, PA 19122

#142/11:00 AM A Parallel Algorithm for the Singular Value Decomposition of Rectangular Matrices

The classical sequential algorithm for computing the singular value decomposition (SVD) of a matrix has time complexity $O(n^3)$ and therefore can only be used on small matrices. Parallel algorithms for computing the SVD have recently been proposed. The best of these use Jacobi rotations to diagonalize the matrix. However, Jacobi-SVD algorithms require that the matrix be square, and, when this is not the case, the matrix must be preprocessed. A parallel algorithm is presented that computes the SVD of a rectangular matrix without preprocessing. The running time and numerical stability of this algorithm is compared against a selection of other parallel SVD algorithms.

William D. Shoaff and David C. Chan Mathematical and Computer Sciences Florida Institute of Technology 150 W. University Boulevard Melbourne, FL 32901-6988

#125/11:20 AM
A Parallel Algorithms for the Singular Value Decomposition

A parallel algorithm for computing the singular value decomposition will be presented. The basic step of the algorithm is a divide and conquer step based upon a rank one modification of a bidiagonal matrix. Numerical difficulties associated with forming the product of a matrix with its transpose are avoided and numerically stable formulae for obtaining left singular vectors after computing updated right singular vectors are derived. A deflation technique is described which together with a robust root finding method will assure computation of the singular values to full accuracy in the residuals and will also assure the orthogonality of the singular vectors.

E.R. Jessup Department of Computer Science Yale University New Haven, CT 06520

D.C. Sorensen
Mathematics and Computer Science Division
Argonne National Laboratory
9700 South Cass Avenue
Argonne, Illinois 60439-4844

#89/11:40 AM
A Multiprocessor Scheme for Singular Value Decomposition

We present a multiprocessor scheme for determining the singular value decomposition of rectangular matrices in which the number of rows is substantially larger or smaller than the number of columns. In this scheme, we perform an initial QR factorization on the tall matrix (either A or A) using a multiprocessor block Householder algorithm. We then apply a one-sided Jacobi multiprocessor method on the resulting upper triangular R to yield $RV = U\Sigma$, from which the desired singular value decomposition is obtained. Preliminary experiments on an Alliant FX/8 computer system with 8 processors indicate speedups near 5 for our scheme over an optimized implementation of the Linpack/Eispack routines which perform the classical bidiagonalization technique. Our scheme performs exceptionally well for rank deficient matrices as well as for those rectangular matrices having clustered or multiple singular values, and may be well suited for applications such as real-time signal processing. We present performance results on the Alliant FX/8 and Cray X-MP/48 computer systems with particular emphasis on speedups obtained for our scheme over classical SVD algorithms.

Michael Berry and Ahmed Sameh

Center for Supercomputing Research and Development University of Illinois at Urbana-Champaign 305 Talbot Lab 104 S. Wright St. Urbana, IL 61801-2932

WEDNESDAY, DECEMBER 2 10:40 AM - 12:00 Noon Sacramento Room

Contributed Presentations 9

#35/10:40 AM A Parallel Homotopy Method for Solving a System of Polynomials

The homotopy method is theoretically able to find all the roots of a system of polynomials. In its numerical implementation, each root is obtained by following a distinct homotopy curve. Tracking the curve very closely may guarantee the finding of all roots, but computationally expensive. A loose following of the curve can expedite the computation, but some roots will be missing due to the merge of curves and the undesirable divergence may

occur. Our proposed method can expedite the computation process and can automatically detect and correct curve merging. Potential parallelism of the proposed method is fully exploited to allow the implementation in parallel processors.

Lionel M. Ni Department of Computer Science

Shui-Nee Chow and Yun-Qiu Shen Department of Mathematics Michigan State University East Lansing, Michigan 48824

#94/11:00 AM
The Granularity of Homotopy Algorithms for Polynomial Systems of Equations

Polynomial systems consist of n polynomial functions in n variables, with real or complex coefficients. Finding zeros of such systems is challenging because there may be a large number of solutions, and Newton-type methods can rarely be guaranteed to find the complete solution list. There are homotopy algorithms for polynomial systems of equations that are globally convergent from an arbitrary starting point with probability one, are guaranteed to find all the solutions, and are robust, accurate, and reasonably efficient. There is inherent parallelism at several levels in these algorithms. Several parallel homotopy algorithms with different granularities are studied on several different parallel machines, using actual industrial problems from chemical engineering and solid modelling.

L. T. Watson

J. Harimoto

J. P. Bixler

Department of Computer Science Virginia Polytechnic Institute & State University Blacksburg, VA 24061

#167/11:20 AM

A Parallel Nonlinear Integer Programming Algorithm Based on Branch and Bound and Simulated Annealing

In this talk, we shall present a model problem which we wish to solve, a sketch of standard branch and bound techniques applicable to the solution of the problem, a short outline of the idea behind simulated annealing, and lastly an upper level description of an algorithm which uses a parallel branch and bound technique simultaneously with a modified simulated annealing process. The two techniques, operating in parallel, and "weakly" connected, complement each other in efficiency and determinacy. The parallel algorithm is being implemented on a multiprocessor

transputer network system operating under an OCCAM shell. We will present comparisons made with standard sequential and parallel branch and bound algorithms, and with sequential and parallel simulated annealing algorithms.

Ken W. Bosworth Mathematics Department Utah State University Logan, Utah 84322-3900

G. S. Stiles Electrical Engineering Department Utah State University Logan, Utah 84322-4120

#131/11:40 AM Recursive Binary Partitions

A Recursive Binary Partition (RBP) is a partition of space or data that is obtained from a single recursive partitioning scheme. The generic scheme makes recursive calls to a splitting function. By varying the definition of the splitting function and the depth of the recursion, one can obtain many of the currently popular partitioning strategies that are being used for vortex methods, domain decomposition and multiobject tracking. The recursive nature of the scheme allows all RBP's to be easily computed and maintained on parallel architectures with binary recursive structures such as hypercubes, the Connection Machine, RP3 and the Butterfly.

George Cybenko Department of Computer Science Tufts University Medford, MA 02155

> Tom Allen Alphatech, Inc. 111 Middlesex Turnpike Burlington, MA 01803

WEDNESDAY, DECEMBER 2

3:40 - 5:00 PM San Diego Room

Contributed Presentations 10

#49/3:40 PM Parallel Architectures for Optical Computing

Recent advances in optical devices, conjunction with the inherent parallelism in optics have created the hope of designing a computer that can achieve tremendous speed improvement over existing electronic computers. In this paper, we present a massively parallel architecture for optical computing. The system is very well suited for applications that are 2-D in nature such as signal and image processing. Next we show how parallel algorithms can be mapped very efficiently onto the system. Then we present the optical implementation of different units comprising the Then we discuss its programmability, and predict its performance, as compared with conventional computers. The key message of this article is that optical computers can be potentially 10 s to 10s faster than today's computers.

> Kai Hwang and Ahmed Louri Computer Research Institute University of Southern California University Park, SAL 324B Los Angeles, CA 90089-0781

#76/4:00 PM Synchronization of Nonhomogeneous Parallel Computations

We investigate the effect of non deterministic execution times upon the "speed-up" factor in the parallel execution of algorithms. This investigation has been triggered by the observations in the modeling of domain splitting algorithm for the numerical solution of partial differential equations. While the modeling of an unsynchronized execution leads to a very high processor utilization, requiring synchronization conditions results in a considerably lower level of performance. In this paper we show that this effect is quite general for nonhomogeneous parallel computations. It is an algorithmic effect independent of hardware and we give analytic results for a variety of nonhomogeneous computations.

Dan C. Marinescu
John R. Rice
Department of Computer Sciences
Purdue University
West Lafayette, IN 47907

#120/4:20 PM Intrinsically Parallel Algorithms

Most algorithms implemented on parallel computers have been optimal serial algorithms, slightly modified or parallelized. An exciting possibility is the search for intrinsically parallel algorithms. These are algorithms which do not have a sensible serial equivalent - any serial equivalent is so inefficient as to of little use.

We describe a multiscale algorithm for the solution of PDE systems that is designed specifically for massively parallel supercomputers. Unlike conventional multigrid algorithms, the new algorithm utilizes the same number of processors at all times. Convergence rates are much faster than for standard multigrid methods - the solution error decreases by up to three digits per iteration.

We present results of the implementation of the algorithm on the 65,536 processor Conection Machine, and we compare the results with more standard algorithms implemented on that computer.

Oliver A. McBryan Dept. of Computer Science, University of Colorado, Boulder, CO 80309-0430.

#124/4:40 PM

A Software Tool For Building Supercomputer
Applications

In this paper we describe a programming tool designed to help users of parallel supercomputers retarget and optimize application codes. The system is an interactive program editing and transformation system that helps the user with this task. Each program that enters the system is completely parsed and all data dependences are recorded. The user then works with the system to restructure his code to a form suitable for a given target architecture. At any time the user can ask the system to make estimates of program performance based on potential memory hierarchy conflicts or vector instruction utilization. The system will generate code for the BBN butterfly or Alliant FX/8.

Dennis Gannon Daya Atapattu Mann Ho Lee Bruce Shei

Department of Computer Science Indiana University Bloomington, Indiana 47405

THURSDAY, DECEMBER 3

10:40 AM - 12:00 Noon Santa Barbara Rooms A-C

Contributed Presentations 11

#15/10:40 AM Large Scale FE Parallel Nonlinear Computations on High Performance Architectures

Here we revisit Finite Element algorithms based on homotopy equations, for implementation on shared memory and local memory multiprocessors, which represent both extremes of today's high performance architectures. To achieve this goal, two nonnumerical algorithms for automatic domain decomposition are developed first. Then, a computational strategy that ties these with a numerical nonlinear algorithm based on homotopy equations is presented and discussed. Its implementation on parallel architectures features two levels of parallelism, namely, concurrency and vectorization. It requires little storage and minimizes synchronization and/or interprocessor communication. Numerical experiments conducted on a variety of multiprocessors (CRAY XMP-4, Alliant FX/8, Encore Multimax, Sequent Balance-8000, iPSC-d7) validate the computational strategy and assess its performance.

Charbel Farhat

Department of Aerospace Engineering Sciences and Center for Space Structures and Controls University of Colorado at Boulder - Campus Box 429 Boulder, CO 80309-0429

Luis Crivelli
Department of Mechanical Engineering
University of Colorado at Boulder - Campus Box 427
Boulder, CO 80309-0427

#34/11:00 AM

A Family of Concurrent Algorithms for the Solution of Transient Finite Element Equations

In this paper, we describe a solution method for problems in structural dynamics. This method makes effective use of the architecture of concurrent computers. The algorithm is based on partitioning the structure into substructures. Each substructure is then processed over a time step independently of the others. Thus, high levels of concurrency can be achieved at this phase of the analysis. The solution of the complete system is constructed by gluing the solution for the substructures. This involves averaging the solution at the interfaces between the subsystems. This averaging scheme is derived by

enforcing the consistency condition on the underlying algorithm. The resulting two parameter algorithm is shown to be unconditionally stable.

B. Nour-Omid

Mechanics and Materials Engineering Laboratory,

Lockheed Palo Alto Research Laboratory,

3251 Hanover Street, Palo Alto, CA 94304.

and

M. Ortiz

Division of Engineering,

Brown University,

Providence, RI 02912.

#103/11:20 AM A Parallel Algorithm for Rapid Computation of Transient Fields

The problem considered here is that of determining the internal electromagnetic field produced by an arbitrary incident pulse impinging on a onedimensional medium with spatially varying permittivity and conductivity profiles. This is done by introducing a special set of Green's functions for the problem. These functions have the property that they can be computed for all time (at a fixed spatial location) via an algorithm which uses as input data the values of the functions for a fixed initial period of time. Although this technique is more computationally intensive than standard methods for computing Green's functions, it is ideally suited for a parallel environment because there is no sequential code involved. Examples comparing this method versus standard techniques are presented.

Robert J. Krueger Applied Mathematical Sciences Ames Laboratory Iowa State University Ames, Iowa 50011

#169/11:40 AM
Analysis of a Parallelized Elliptic Solver for Reacting Flows

A parallelized finite difference code for systems of nonlinear elliptic boundary value problems in two dimensions, based on Newton's method, is analyzed in terms of computational complexity and parallel efficiency. An approximate cost function depending on 15 dimensionless parameters (including discrete problem dimensions, convergence parameters, and machine characteristics) is derived for algorithms based on stripwise and boxwise decompositions of the domain and a one-to-one assignment of the strip or box subdomains to processors. The sensitivity of the cost function to the parameters is explored

throughout a region of parameter space corresponding to a coupled system of nineteen equations with very expensive function evaluations (a detailed-kinetics diffusion flame). The algorithm has been implemented on parallel hardware and some experimental results for a flamesheet calculation with stripwise decompositions are presented and compared with the theory.

David E. Keyes & Mitchell D. Smooke Department of Mechanical Engineering Yale University P.O. Box 2159 Yale Station New Haven, CT 06520

THURSDAY, DECEMBER 3

10:40 AM - 12:00 Noon San Diego Room

Contributed Presentations 12

#196/10:40 AM

Performance Analyses of the Nested-Dissection Algorithms for MIMD Multiprocessor Systems

Sequential algorithms based on the idea of nested dissection orderings are known to be effective in solving sparse positive definite linear systems which usually arise from grid problems. Recently, several attempts have been made in developing efficient parallel algorithms for solving such sparse linear systems. Most of these parallel algorithms assume mesh-connected, systolic type architectures or massively parallel systems. In this paper we consider various performance issues related to implementing the nested-dissection algorithms on some of the currently available MIMD architectures. Analytic models are developed and are verified by implementing the methods on a 64-node BBN Butterfly multiprocessor system.

Vijay K. Naik Merrell L. Patrick

Institute for Computer Applications in Science and Engineering NASA Langley Research Center Hampton, VA 23665

and

Computer Science Department Duke University Durham, NC 27706 #190/11:00 AM
A Proposal for a Set of Level 3 Basic
Linear Algebra Subprograms

This Presentation describes a proposal for a set of Level 3 Basic Linear Algebra Subprograms (Level 3 BLAS), which are targeted at matrix-matrix operations, with the aim of providing more efficient, portable implementations of algorithms on high-performance computers, especially those with hierarchical memory and parallel processing capability.

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Jeremy Du Croz; Sven Hammarling Numerical Algorithms Group Ltd NAG Central Office, Mayfield House 256 Banbury Road Oxford OX2 7DE U.K.

Iain Duff
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#56/11:20 AM
Gaussian Techniques on Shared Memory
Multiprocessor Computers

Implementations of Gaussian elimination and Gauss-Jordan algorithms for a shared memory multiprocessor simulator are presented. Gaussian elimination algorithms require processor synchronization due to data dependencies. We show the necessity for a hardware implementation of processor synchronization. The efficiency of Gauss-Jordan versus Gaussian elimination with software and hardware processor barriers is also given. Even though Gauss-Jordan requires more operations than Gaussian elimination for a given size problem, Gauss-Jordan is faster for certain matrix sizes and numbers of processors. The cross-over points are shown for both hardware and software barriers. The algorithms are written in C and are run on the Cerberus Multiprocessor Simulator, a simulation facility for scalable shared memory multiprocessors with pipelined functional units.

Gregory A. Darmohray, Eugene D. Brooks III

Computational Physics Division University of California Lawrence Livermore National Laboratory P.O. Box 808, L-298 Livermore, CA 94550 #41/11:40 AM

LU Factorization on Distributed-Memory Multiprocessors

We discuss and compare two methods of improving the efficiency of LU factorization with pivoting on distributed-memory multiprocessors. The first method uses a dynamic load balancing scheme to avoid work imbalances caused by pivoting. While this method increases the amount of communication required, it significantly reduces the execution time. The second method uses pipelining to mask the cost of pivoting. While pipelining complicates the algorithm, the execution time is slightly less than in the first method. Both algorithms have achieved efficiences of 85% to 95% running on a 64 node iPSC hypercube.

George A. Geist Charles H. Romine

Mathematical Sciences Section Engineering Physics and Mathematics Division Oak Ridge National Laboratory P. O. Box Y, Bldg. 9207-A, MS-3 Oak Ridge, Tennessee 37831

THURSDAY, DECEMBER 3

10:40 AM - 12:00 Noon San Gabriel Rooms A-C

Contributed Presentations 13

#36/10:40 AM

Semi-Analytical Shape Functions in a Parallel Computing Environment

Semi-analytical shape functions have been used successfully to model many physical systems in continuum mechanics. The orthogonality of appropriately selected shape functions allows a continuous system to be effectively transformed into a series of smaller discrete subsystems that can be solved independently. uncoupling of the subsystems produces a highly parallel algorithm that is well suited to either a shared or a distributed memory environment. The semi-analytical shape functions are discussed with regard to the algorithmic implications. Examples are presented to illustrate the methodology.

J. A. Puckett
R. J. Schmidt
Institute for Scientific Computation
University of Wyoming
Laramie, Wyoming 82071

#145/11:00 AM
The P-S Shape Functions

Parallel Computing Environment

A shape function which combines polynomial and sine series functions (P-S Shape Functions) is used to develop an efficient finite element algorithm for the solution of the elliptic equations that model many phenomena in continuum mechanics. This nontraditional approach results in an algorithm that is massively parallel, adaptive, and easily implemented using standard numerical techniques. The key to these favorable features is the orthogonality of the sine series which uncouples element matrices associated with the series terms. Because of the special structure of the system equation, condensation methods are used to form a series of matrix equations that are much smaller than the original system equation. These equations can be triangularized independently and thus the algorithm is highly parallel.

J. A. Puckett
R. J. Schmidt
Department of Civil Engineering
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#90/11:20 AM

Invariant Imbedding and the Method of Lines for Multiprocessor Computers

Recursive relations have been used to allow the solution of invariant imbedding equations with singularities. We demonstrate that these same relations can be used in an efficient implementation of invariant imbedding for multiprocessor computers and investigate the maximum attainable speedup.

The parallel implementation of imbedding can be used in conjunction with the method of lines to solve partial differential equations. We consider the problem of assigning lines to processors to minimize communication delays. We also consider the question of convergence of the resulting asynchronous algorithm. Timing and speedup data are presented for algorithms implemented on the NCUBE hypercube.

Richard C. Allen, Jr. Lorraine S. Baca David E. Womble

Sandia National Laboratories Albuquerque, NM 87185 #128/11:40 AM
A Portable Parallel Algorithm for Multivariate Numerical Integration and Its Performance Analysis

An adaptive algorithm for numerical integration over an N-dimensional cube will be discussed, together with its implementation in a portable manner on a range of MIMD shared memory architectures.

Results from different machines will be given, showing good speedups especially in higher dimensions, where a large granularity is in effect.

We shall investigate the influence of various elements of the algorithm on the performance. A technique for a-priori performance analysis, on the basis of the serial code, will be applied and justified at the hand of performance results on the Sequent Balance, Alliant FX-8 and Encore Multimax.

E. deDoncker John Kapenga Department of Computer Science Western Michigan University Kalamazoo, MI 49008-5021

THURSDAY, DECEMBER 3

10:40 AM - 12:00 Noon Sacramento Room

Contributed Presentations 14

#84/10:40 AM

Preconditioned Conjugate Gradient Methods for General Sparse Matrices on Shared Memory Machines

For very large general sparse linear systems such as those arising from three-dimensional models, preconditioned Krylov subspce methods represent a good alternative to using direct methods. When implementing these methods on parallel machines one faces the problem that the standard preconditioners are very sequential and become a serious bottleneck. The first remedy is to use a polynomial preconditioning which consists of approximating A by a polynomial in A. This preconditioning technique offers a high degree of parallelism and its performance can be enhanced by a simple diagonal or block-diagonal scaling. An alternative is to parallelize the standard Incomplete LU (ILU) preconditioners by performing the incomplete factorization in a rank-one modification manner instead of computing one row of L and U at a time as is usually done. This implies manipulating the row and column structures of A but the resulting overhead is outweighed by the benefits of parallelism. Numerical experiments will be presented comparing these approaches.

Edward Anderson Youcef Saad University of Illinois at Urbana-Champaign Center for Supercomputing R & D 305 Talbot Laboratory 104 South Wright Street Urbana, Illinois 61801-2932 USA #52/11:00 AM

PARALLEL MULTI-LEVEL FINITE ELEMENT METHOD WITH HIERARCHICAL BASIS FUNCTIONS

The p-version of the finite element method using a hierarchy of basis functions is employed in a multi-level computational scheme closely resembling the multi-grid approach. The advantage is that it is possible to preserve the speed of multi-grid techniques while avoiding the very difficult problem of generating a sequence of nested grids on complex geometries. The parallel implementation of this technique on a Loosely Coupled Array of Processors at IBM Kingston is described, and results are presented for the Poisson equation on a square with Dirichlet boundary conditions, using a standard smoothing technique.

G. Brussino IBM Corporation Department 48B Kingston, NY 12401, USA.

R. Herbin IBM Corporation Department 48B Kingston, NY 12401, USA.

V. Sonnad IBM Corporation Department 48B Kingston, NY 12401, USA. Present address: Alliant Computer Systems Corp. Littleton Massachusetts 01460 USA.

Present address: Ecole Polytechnique de Lausanne GASOV / ASTRID 1015 Lausanne, Switzerland.

#48/11:20 AM A Double Augmented Tree Architecture for Parallel Execution of Multigrid Algorithms

Abstract: We propose a multiprocessor architecture for the efficient execution of multigrid PDE algorithms. Using a binary tree as the skeleton, the processors of the same tree level are connected into a ring to form an augmented tree. Two augmented trees are stacked together with their roots coalesced to form an X-shaped tree structure. Leaf nodes of the two augmented trees are also used for I/O purpose. Feedback lines are provided from output nodes to input nodes. The architecture is shown to be well suited to parallel processing of multigrid algorithms via a close match of hardware links with the communication patterns of the algorithms. Both conventional and concurrent multigrid algorithms are considered. Performance of this architecture is compared with the hypercube architecture.

H.C. Wang and Kai Hwang

Department of Electrical Engineering-Systems
University of Southern California
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Los Angeles, CA 90089

#38/11:40 AM

Implementation of a Parallel Multigrid Method on a Loosely Coupled Array of Processors

A parallel method with multicolour ordering has been implemented on the loosely Coupled Array of Processors at IBM Kingston. We consider various smoothers, restrictions and prolongations. The parallelization is obtained by an a-priori decomposition of the finest grid; each step of the multigrid method is then vectorized and parallelized, up to the direct solve on the coarse grid, which is done sequentially; data communication via the shared memory occurs at the boundaries of the subregions thus defined. We show that because of the relative parallel unefficiency of the smoother on coarser grids, it is more interesting to do the direct solve on a grid of level superior to 1. We present the numerical results that we obtained for the different algorithms, along with the speed-ups according to the number of processors that we use.

Raphaele Herbin

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Kingston, NY, USA. 1015 Lausanne, Switzerland.

Stephane Gerbi E.N.T.P.E. Lyon, France.

Present address: Yale University,

Computer Science Department,

New Haven, CT, USA.

Vijay Sonnad IBM Corporation, Department 48B Kingston, New York, USA.

THURSDAY, DECEMBER 3

3:40 - 5:00 PM San Diego Room

Contributed Presentations 15

#127/3:40 PM

SCHEDULE: An Environment for Developing Transportable Explicitly Parallel Codes in Fortran

SCHEDULE is a software package designed to aid in programming explicitly parallel algorithms for numerical calculations. The design goal of SCHEDULE is to aid a programmer with implementation of explicitly parallel algorithms in a style of Fortran programming that will lend itself to transporting the resulting programs across a wide variety of parallel machines. The approach relies upon the user adopting a particular style of expressing a parallel program. Once this has been done the subroutines and data structure provided by SCHEDULE will allow implementation of the parallel program without dependence on specific machine intrinsics. The basic philosophy taken here is that Fortran programs are naturally broken into subroutines which identify units of computation that are self-contained and which operate on shared data structures.

A graphics post processor has been developed to analyze the execution pattern and structure of a parallel program. A graphics pre-processor is under development that will aid in constructing a parallel program directly from the large grain data dependency graph.

J.J. Dongarra and D.C. Sorensen Mathematics and Computer Science Division Argonne National Laboratory 9700 South Cass Avenue Argonne, IL 60439-4844

#140/4:00 PM

Domain decomposition methods for three-dimensional elliptic problems.

Decomposing a domain in R 3 by parallel vertical planes leads to a block structure of the stiffness matrix where it is natural to form the reduced Schur complement matrix corresponding to the dividing planes. Formation of the matrix is in general impossible so the Schur complement system must be solved by iteration. This requires only the computation of the action of the matrix but this action requires the solution of the same problem on the subdomains. We describe how an effective (optimal order) preconditioner can be constructed and how a recursive subdomain method can be used for the solution of the subdomain problems and of the preconditioner. Pre- and postsmoothing of the residuals on the fine mesh implies that the correction can be computed on the coarse mesh which can decrease the cost by a factor close to 1/8. We describe also how one can implement the method on a cluster of loosely coupled parallel processors where each cluster consists of more tightly coupled processors.

O. Axelsson
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#146/4:20 PM Fast Poisson Solver On Irregular Regions by Boundary Integral-based Domain Decomposition

A method, called the Boundary Integralbased Domain Decomposition (BIDD) method, is proposed for a fast solution of the Poisson equation on irregular regions. The idea of the method is to adopt the domain decomposition approach by embedding a number of regular subdomains into the irregular problem domain, but to use the standard integral equation formulation of the potential theory to compute the interface values. The method provides a particularly efficient algorithm for parallel processing. Some results of numerical experiments on the Alliant multiprocessor will be reported.

Daeshik Lee, E. Gallopoulos Center for Supercomputing Research and Development University of Illinois 305 Talbot Lab 104 S. Wright St. Urbana, IL 61801-2932

#178/4:40 PM

A High-Performance FFT Algorithm for Vector Supercomputers

Many traditional algorithms for computing the fast Fourier transform (FFT) on conventional computers are unacceptable for advanced vector and parallel systems because they employ nonunit, power-of-two memory strides. This paper presents a technique for computing the fast Fourier transform that avoids such strides and appears to be near-optimal for a variety of current vector and parallel computers. Performance results of a program based on this technique are presented. Notable among these results is that a Fortran implementation of this algorithm on the Cray-2 runs up to 75% faster than Cray's assembly-coded library routine.

David H. Bailey NASA Ames Research Center Mail Stop 258-5 Moffett Field, CA 94035

FRIDAY, DECEMBER 4 10:40 AM - 12:00 Noon San Gabriel Rooms A-C

Contributed Presentations 16

#153/10:40 AM Parallel Vision Algorithms: An Approach

In this work, we address the issue of implementing some common image processing applications on parallel processors. Our work involved two non Von-Neumann machines, viz. the Butterfly and the NCUBE. Our approach is fundamentally algorithmic in nature.

One basic technique used in computer vision is the Hough technique of feature detection. An application of this is the detection of collinear or nearly collinear feature points. These are evidences of edges, and thus shapes of objects in images.

We define an approach to proper processor management depending on the two quantities as parameters: the image size and the number of processors. Our approach led to a linear speedup.

Sharat Chandran Larry S. Davis Center for Automation Research University of Maryland College Park, MD 20742-3411

#154/11:00 AM
Performance of a Ocean Circulation Model on LCAP

A numerical ocean model based on the primitive equations has been parallelized and run on the ICAP system - an experimental parallel machine with an IBM host and ten FPS-x64 attached processors. Using coarse-grain parallelism approach, the computational domain is partitioned into a number of sub-domains according to one of the horizontal directions. Each attached processor computes one sub-domain. Inter-processor communication is accomplished through a set of shared bulk memories. Five synchronization (barrier) points are inserted into the program to guarantee a correct sequence of execution.

A set of ocean model benchmarks have been carried out. The best result shows an impressive speed-up of 6.9 on a 10-processors scale can be achieved.

Hsiao-Ming Hsu Department of Physical Oceanography Woods Hole Oceanographic Institution Woods Hole, Massachusetts 02543

Jih-Kwon Peir Computer Science Department IBM Thomas J. Watson Research Center Yorktown Heights, New York 10598

Dale B. Haidvogel The Chesapeake Bay Institute The John Hopkins University Baltimore, Maryland 21211

#28/11:20 AM

Ising spin on a shared memory machine: computational experience.

Ising spin simulations occur in areas such as studies of phase transitions, and alloy solidification. For a large system, hundreds of thousands of updates spin sites must be performed in order to approach a single equilibrium configuration, and averages over many configurations are required.

The kernel process appears to be "iteration on regular grids." This suggests parallel algorithms related to those used for

solving elliptic pde's by finite difference approximations. We have devised algorithms which have similarities with the techniques of red-black ordering and chaotic relaxation. Preliminary tests indicate that, for a five processor machine, one can expect a speedup factor of 3.5 to 4.5.

James L. Blue and Francis Sullivan Center for Applied Mathematics National Bureau of Standards Building 101, Room A438 Gaithersburg, MD 20899

#189/11:40 AM
Parallel Iterative Algorithms To Solve The Discrete LQR
Optimal Control Problem With Hard Control Bounds

In this paper we present two gradient based parallel iterative algorithms to solve the linear quadratic regulator (LQR) optimal control problem with hard control bounds. The first algorithm is a parallel implementation of a parametrized gradient projection method and the second algorithm is a parallel implementation of a combination of the first algorithm with a linearly constrained version of the Fletcher-Reeves conjugate gradient method. We show that at each iteration parallel evaluation of the step length and projected gradient of the quadratic cost function can be efficiently performed as a function of the number of processors. We then embed our parallel step length and gradient projection procedures to produce two parallel algorithms which are suitable for real-time online implementation on a SIMD machine.

Gerard G. L. Meyer Electrical and Computer Engineering Department The Johns Hopkins University Baltimore, Maryland 21218

Louis J. Podrazik Bendix Environmental Systems Division Allied-Signal, Inc. Baltimore, Maryland 21284-9840

FRIDAY, DECEMBER 4

10:40 AM - 12:00 Noon Sacramento Room

Contributed Presentations 17

#175/10:40 AM

A Comparison of the Performance of LINPACK and Vectorized ICCG in Solving a Poisson Equation for a Problem of Two-Phase Flow in a Centrifuge

A problem in fluid dynamics, two phase fluid flow in a centrifuge, is used to compare the direct solution of linear systems using LINPACK routines with the iterative solution of linear systems using Incomplete Cholesky Conjugate Gradient (ICCG). We compare two different approaches to vectorize ICCG, using the diagonal ordering, which yields shorter vector lengths, O(n), for an $n \times n$ grid, and using multicolor orderings which yield long vector lengths, $O(n^2)$. The key consideration is the greater degree of vectorization for the multi-color approach versus the somewhat better convergence results for the diagonal ordering. Results are given for runs on a CYBER 855, Convex C-1, Cray-2, Cray-XMP, and CYBER 205.

Dr. Eugene L. Poole Mail Stop 244 Nasa Langely Research Center Hampton, Va. 23665

Jeffery W. Frederick Dept. of Mechanical Engineering University of Virginia Charlottesville, Va. 22901

#187/11:00 AM

On the Solution of Block Tridiagonal Systems of Equations on Multi-Vector Machines - A Comparitive Study

The problem of solving scalar tridiagonal systems has been extensively studied by various researchers. This paper examines the solution of tridiagonal systems in which the non-zero elements of the system are mxm matrices, forming a block tridiagonal system of equations.

Two solution methods are examined: Ritchmeyer's algorithm, and a cyclic reduction based algorithm. It is shown that on multi-vector machines such as the Alliant FX/8, the reduction algorithm takes less time to execute compared to the Ritchmeyer's algorithm.

Ken Dowers, S. Lakshmivarahan, Sudarshan K. Dhall, and Julio Diaz Parallel Processing Institute School of Electrical Engineering and Computer Science University of Oklahoma Norman, OK 73019

#112/11:20 AM

High Performance Linear Algebra on the AMT DAP500

New techniques will be presented that greatly increase the linear algebra performance of the AMT DAP500, a highly parallel SIMD array machine rated at 9 MFLOPS. Many algorithms, such as solving equations or matrix multiplication, can be expressed as an accumulation of outer products of pairs of vectors. By mapping one vector across the array and multiplying it successively by numbers from the other vector, use can be made of new standard functions defined at a level above the individual array operation. Results will be presented: for example, solving 1000 x 1000 equations improves from 7 to over 30 MFLOPS at 32-bits.

Stewart F. Reddaway ICL Cavendish Road Stevenage, Herts SG1 2DY England

#134/11:40 AM

Performance of Blocked Gaussian Elimination on Multiprocessors

A performance analysis of blocked Gaussian elimination of a symmetric positive-definite system on a simple MIMD model is presented. The adopted MIMD model assumes a high bandwidth switch or network with bounded bandwidth ports to global and local memories. The blocking is shown to reduce global references to a non-interfering level, under certain mild restriction, so that synchronization and block-overhead are the important performance considerations. The analysis explains our experimental results on the BBN Butterfly multiprocessor: both the strongly limited speedup of the row-based algorithm and the far more extensive speedup of the block-based algorithm.

Elizabeth Jean O'Neil BBN Advanced Computers, Inc. 10 Fawcett Street Cambridge, MA 02238

Henno Allik
BBN Laboratories Inc.
Union Station
New London, CT 06320

FRIDAY, DECEMBER 4

10:40 AM - 12:00 Noon Santa Barbara Rooms A-C

Contributed Presentations 18

#98/10:40 AM,

Loop Blocking for Parallel Memory Optimization

Researchers have found that blocking of loops, judiciously applied, can improve the performance of computers by improving the utilization of virtual memory, cache memory, local memory or registers. We present loop blocking as a compiler optimization that is closely related to strip mining, a method used by compilers to generate code for vector registers. Loop blocking combined with index set splitting, loop interchanging and concurrentization is used to automatically optimize linear algebra algorithms for parallel machines with private memories.

Michael Wolfe Kuck and Associates, Inc. 1808 Woodfield Drive Savoy, IL 61874

#150/11:00 AM

Parallel FORTRAN: Why You Can't. How You Can.

With the advent of a variety of parallel processing systems, computational investigators are looking for convenient programming models and tools for executing their favorite codes. The first choice of investigators in the physical sciences and engineering appears to be access to parallel processing within FORTRAN. Unfortunately, the basic rules and requirements of an operational FORTRAN product strongly conflict with the most obvious primitives for parallel processing. I will discuss reasons why FORTRAN can not be formally extended for parallel processing without transforming the language beyond recognition (and beyond its operational requirements). I will then discuss an alternative for altering FORTRAN for parallel processing without any formal extensions. This is done with a set of a dozen compiler directives. The goal is to invent a set of directives that will be portable and allow efficient code to be generated for a large variety of parallel processing architectures.

Clifford Arnold ETA Systems, Inc. 1450 Energy Park Dr. St. Paul, MN 55108, USA #180/11:20 AM Performance Based Distributed Programming

Writing parallel and distributed programs will be very difficult. A number of approaches have been investigated to look at the viability of writing normal, sequential programs without consideration for the hardware architectures or the environmental characteristics. A separate decomposition configuration or specification is then written that will partition the software at runtime. By analyzing the software and hardware as separate steps with an automated specification generator, the program can optimally applied to a larger number of architectures. The automation techniques will be very important especially for the development of large systems for hard real-time systems.

Phillip Q. Hwang Director, ACSAD Laboratory Naval Surface Weapons Center, U33 White Oak Detachment Silver Spring, Maryland 20903-5000 (202)394-1351

#20/11:40 ÃM

A Programming Aide for Message-passing Systems

A programming aide, which will automatically partition and insert synchronization primitives where needed may increase the programming productivity and reduce the number of errors by an order of magnitude. This paper describes an interactive tool for message-passing systems. The tool performs the dataflow and dependency analyses, partitions programs into serial and parallel segments and simple loops into parallel processes, and inserts synchronization primitives automatically or interactively. It also provides analysis and performance estimates for complex loops. The time for developing programs on message-passing systems is expected to be drastically reduced.

Min-You Wu and Daniel D. Gajski Department of Computer Science University of Illinois at Urbana-Champaign Urbana, IL 61801

FRIDAY, DECEMBER 4

10:40 AM - 12:00 Noon San Diego Room

Contributed Presentations 19

#10/10:40 AM

A Gray-Code Scheme for Local Uniform Mesh Refinement on Hypercubes

Adaptive methods for PDEs can be viewed as a graph problem. Parallel methods must distribute this graph efficiently among the processors. In doing this, the cost of communication between processors and the structure of the graph must be considered. We divide this problem into two phases: labeling of graph nodes and subsequent mapping of these labels onto processors. We describe a new form of Gray-code which we call an interleaved Gray-code that allows easy labeling of graph nodes even when the maximal level of refinement is unknown, allows easy determination of nearby nodes in the graph, is completely deterministic, and often (in a well-defined sense) distributes the graph efficiently across a hypercube.

Computational experiments will be presented.

William D. Gropp Ilse C.F. Ipsen Dept. of Computer Science Yale University P.O. Box 2158 Yale Station New Haven, CT 06520

#43/11:00 AM

Iteration with Mesh Refinement for Systems of ODE's.

We consider block Picard - Lindelöf iteration (or "waveform relaxation") for systems of initial value problems. A rigorous environment for the following fact is given: one can refine the mesh during the iteration in a reliable way so that the total amount of work is a small multiple of that for the decoupled system.

Olavi Nevanlinna Helsinki University of Technology Institute of Mathematics SF-02150 Espoo, Finland

#108/11:20 AM

Parallelization of Adaptive Grid Domain Mappings

Adaptive grid domain mappings are used as a tool for grid adaption in the numerical solution of partial differential equations. They allow the power of adaptive moving-grid methods to be applied to difficult problems, without introducing irregularities into the grid which may make good vector or parallel numerical methods harder to achieve. The three main computational steps associated with these mappings are: 1) constructing the mapping by tensor product least squares with constraints, 2) evaluating the mapping by tensor product spline evaluation, 3) inverting the mapping by a two-dimensional secant method. The speaker will describe parallel implementations of these three steps on the Sequent Balance 21000 multiprocessor. Issues affecting performance and speedup will be discussed.

Calvin J. Ribbens
Department of Computer Science
Virginia Polytechnic Institute and State
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Blacksburg, VA 24061

#25/11:40 AM

Automated Mesh Decomposition and Concurrent Finite Element Analysis for Hypercube Multiprocessor Computers

A concurrent finite element formulation for linear and nonlinear transient analysis using an explicit time integration scheme has been developed for execution on hypercube multiprocessor computers. The formulation includes a new decomposition algorithm which automatically divides an arbitrary finite element mesh into regions and assigns each region to a processor on the hypercube. The decomposition algorithm is deterministic in nature and relies on a scheme which reduces the bandwidth of the matrix representation of the connectivities in the mesh. Numerical results obtained from a 32 processor Intel hypercube (the iPSC/d5 machine) are presented. Speedup factors of greater than 31 have been obtained thereby demonstrating 95% efficient use of the machine.

James G. Malone Engineering Mechanics Department General Motors Research Laboratories Warren, MI 48090-9057 TUESDAY, DECEMBER 1
5:00 - 7:00 PM
Pasadena Exhibit Hall
POSTER SESSION 1

#1 Scheduling of Precedence-Constrained Tasks on Multiprocessors

A major factor in the intelligent utilization of multiprocessor systems is the determination of how to assign computational tasks among processors and to schedule their execution. To this end, we consider a set of precedenceconstrained tasks, with arbitrary communication among them. We investigate several approaches to the problem of assigning and scheduling the tasks in order to achieve maximum parallelism and minimum communication overhead in both fullyconnected and hypercube multiprocessor ensembles. Three simple heuristics and an adaptation of simulated annealing are described. Results of computational experiments provide a comparative analysis of the performance of these heuristics, and give insight into their appropriate application for scheduling distributed task systems.

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M. A. Salama Jet Propulsion Laboratory California Institute of Technology Pasadena, California 91109

#2 ITERATIVE SUBOPTIMAL STRATEGIES FOR IMPROVING THE TRANSIENTS IN ADAPTIVE CONTROL SYSTEMS

Iterative solutions which involve the use of iterative optimization techniques are given to improve the adaptation transients in adaptive systems. Two suboptimal controllers are developed and applied to an equivalent discrete—time system valid to describe the behaviour of a recent adapti—

ve control scheme when at least one of the timevarying parameters entering the adaptation algorithm varies within a close domain being admissible for convergence of the estimates. The resulting suboptimal control strategies are translated into corrections of this parameter so as to recompute the adaptation algorithm over a finite horizon. Numerical simulations illustrate the efficiency of the proposed scheme.

M. DE LA SEN AND M.J. GONZALEZ-GOMEZ

DEPARTAMENTO DE FISICA FACULTAD DE CIENCIAS

UNIVERSIDAD DEL PAIS VASCO

LEIOA-BIZKAIA-SPAIN

#5 A Comparison of Hydrodynamics Multitasking Algorithms

Three one-dimensional, time-dependent hydrodynamics algorithms were multitasked on a Cray X-MP/416. An Eulerian method, a Lagrangian method with artificial viscosity, and a Lagrangian method using a Riemann solver were compared. Speedup results were obtained on a dedicated machine; that is with no one else running. The algorithm changes needed to multitask each of the methods are discussed. Differences in the speedups are discussed as well as the debugging problems encountered. The results were obtained for a sample problem involving a shock wave moving down a channel. The multitasking results are identical to the serial results and agree with the analytical solution.

David A. Mandell Computational Physics Group Los Alamos National Laboratory Los Alamos, NM 87545

ABSTRACTS: POSTER PRESENTATIONS

#6 EFFICIENT PARALLEL ALGORITHMS FOR CONTROLABILITY AND EIGENVALUE ASSIGNMENT PROBLEMS*

The design and analysis of time-invariant linear control systems give rise to a variety of interesting linear algebra problems. Numerically viable sequential algorithms now exist for most of these problems; however, efficient parallel algorithms are virtually non-existent. In this paper, we propose efficient parallel algorithms for multiinput controllability problem and the single-input pole assignment problem. A desirable feature of these algorithms is that they are composed only of basic linear algebraic operations such as vectormatrix multiplication, solution of a linear system and computationa of eigenvalues or singular values of a symmetric matrix, for which efficient parallel algorithms have already been developed. Thus, the proposed algorithms abve potentials for implementations on some existing and future parallel proces-

B. N. Datta and Karabi Datta Department of Mathematical Sciences Northern Illinois University DeKalb, IL 60115 USA

#7
Use of Matrix Determinant of Kronecker Product
In Parallel Matrix Computations

Using the idea of matrix determinant of Kronecker product of two matrices A and B, we show here how several computational problems in linear algebra like (a) finding if two matrices A and B have eigenvalues in common (b) finding a symmetrizer of a matrix and (c) computing the inertia of a nonhermitian matrix can be solved in parallel in o(log n) steps using a bounded number of processes. At this point, the results are of theoretical interests only. Some new properties of the matrix determinant are also developed in this paper.

Karabi Datta
Department of Mathematical Sciences
Northern Illinois University
DeKalb, IL 60115 USA

#8 A New Efficient Parallel Algorithm for Solving Tridiagonal Linear Systems

How to solve tridiagonal linear efficiently is of great signi new parallel algorithm is " ırst we transform the tridiar lions to a boundary value r second take [log N] ma' cANCELLED necessary in cancelled . then we S compute the coefficients simultanec ally we take one step to get t .on. If N processors are available only take 5[log N]+0(1) times. Oth cwise $22.5[N/P]+5[\log P]+O(1)$

times are needed if we have P processors (P<N). Comparing with the other methods, such as recursive doubling method, cyclic reduction method, etc., the new algorithm presented in this paper is more efficient.

Niaoshu Pan Department of Mathematics, Physics & Mechanics Nanjing Aeronautical Institute Nanjing, 210002 People's Republic of China

#11 Task Assignment in a Multiprocessor System

This paper presents a simple and efficient heuristic algorithm for assigning a set of partially ordered computational tasks with computation and communication constraints onto a three-dimensional bus-based multiprocessor system so that the total execution time is minimized. The algorithm uses two cost functions: (i) "Minimum Distance" and (ii) "Parallel Transfer" to minimize the completion time. Extensive simulation results show that the heuristic algorithm is highly effective. The results are compared to a "semirandom" assignment algorithm. The proposed algorithm has also been used to solve some practical numerical problems where optimal assignments have been obtained.

C. Siva Ram Murthy, V. Rajaraman Department of Computer Science and Automation Indian Institute of Science Bangalore 560012, INDIA

#12 Algorithm and Performance Notes Concerning Block LU Factorization

A common matrix structure occurring in static finite element analysis is a symmetric banded matrix having a significant number of zeros below the band, known as a skyline system. Since most finite element applications lead to symmetric, positive definite matrices, pivoting is not necessary, which preserves the original structure. This paper first reviews the standard implementation of Crout factorization for solving such a system which is limited to vector speeds. The remainder of the paper discusses blocking strategies for the matrix skyline for two dense algorithms. The first is the matrix-vector version of Crout factorization and the second is an approach based on products of multiple Gauss transformation matrices.

Jim Armstrong
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ABSTRACTS: POSTER PRESENTATIONS

#13

Use of Mathematical Switches to Solve Differential Equation Problems.

Mathematical switches have been used to solve differential equation problems in a multischeme approach. The use of the switches is equivalent to select automatically a suitable scheme among several schemes at every computing step. The use of three switches has been examined in a multiprocessor computer. Numerical results of one switch, the switch of exact error, confirm that a combined solution of several schemes is indeed more accurate than one of a single scheme. The use of switch of estimated accumulated error demonstrates that a suitable scheme selection automatically at every computing step is feasible. A multischeme approach with properly designed mathematical switches can certainly resolve many numerical difficulties which are not solvable today.

Yi-ling F. Chiang Computer and Information Science New Jersey Institute of Technology Newark, New Jersey, 07102

#16 Block Truncated-Newton Methods for Unconstrained Minimization

Truncated-Newton methods are a class of optimization methods suitable for largescale problems. At each major iteration, a search direction is obtained by approximately solving the Newton equations using an iterative method. In this way, matrix costs and second-derivative calculations are avoided, hence removing the major drawbacks of Newton's method. In this form, the algorithms are well-suited for vectorization. Further improvements in performance are sought by using block iterative methods for computing the search direction. In particular, conjugategradient-type methods are considered. These still permit vectorization of inner loops, while introducing parallelism at a coarser level.

Stephen G. Nash Dept. of Operations Research and Appl. Statistics George Mason University Fairfax, VA 22030 #21 Solving non-structural problems via a structural analysis package

MSC/Nastran is widely known and accepted as the world's leading finite element package for structural analysis problems. To accomplish thi the package has very advanced mathematical ٦t, especially in the numerical linear a ₹. It is not well known however, that an has a high level languague, which a ser to utilize these mathematical or non-struc tural problems. The name of this lange AP, which stands ving nonsquare children Programming. roblems arise when solor linear equations, whose .. differ in number. The pape sent the method of solving large _are problems with DMAP . linear Prac periences will be demonstrated on .cor and parallel computers. scalar

Louis Komzsik Senior Numerical Analyst Technical Services Department The MacNeal-Schwendler Corporation 815 Colorado Blvd. Los Angeles,CA 90041

#22 Calculating Good Initial Guesses for Iterative Solutions of Parabolic Partial Differential Equations in Parallel

An iterative procedure for solving parabolic partial differential equations with multiple space dimensions is presented. Such problems arise in fluid dynamics. Areas of rapid change are located by taking an explicit time step that exceeds the Courant condition on a coarse mesh and applying gradient based domain decomposition. The resulting subareas are solved in parallel. Schwarz outer iterations are used for overlapping subareas. The advanced subareas are superimposed on the explicit solution. The result is used as the initial guess for the coarse mesh iteration. Application of this method improved convergence rates on the coarse mesh tenfold.

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ABSTRACTS: POSTER PRESENTATIONS

#23

Systematic Method of Condition/Event Nets Design

Many different equivalence notions for comparing concurrent systems have been introduced in different contexts. Aim of this paper is to present minimization procedure and systematic method of condition/ event nets design. The notions considered are based on the assumption that the behaviour of a system can be observed or experienced by an observer only communicating with it. The observer can be considered as a possible environment in which the system is embedded and with whom the system interacts. If two different systems cannot be distinguished by any such observer, then they can be substituted in a con sistent way with respect to the interactions with the environment, whatever the environment is. The equivalence notions should be therefore congruences at least with respect to the parallel composition, i.e. in terms of nets with respect to the superposition of events.

Boleslaw Mikolajczak Computer and Information Science Department Southeastern Massachusetts University North Dartmouth, MA 02747

#26

Neural Nets and Sparse Matrices

Recently there has been increased interest in neural nets as novel approach to the handling of noisy and inexact data. Neural networks are modeled after powerful parallel biological computing systems and are of superior performance for a variety of cognitive tasks. Clearly the structure of a neural net can be modeled by a large, rectangular, sparse matrix. In this talk we show that some of the recent work by erson and his collaborators can be formula near algebra terms as sparse matrix ur ، ۱۱۱۵ We will express some of Anders .chms in sparse matrix terminology, .e convergence of a learning r the combination of recent advances in ix computations with neural nets or sibility of large scale simulation nets on todays supercomputers.

Horst D. Simon

Engineering Technology Applications Division Boeing Computer Services, M/S 7L-21 P.O. Box 24346 Seattle, WA 98124 #27 Large Memory Applications on the SCS-40

The SCS-40 is a minisupercomputer with a peak performance of about 44 MFLOPS, 4 Mwords of memory, and a Cray X-MP compatible architecture. In this paper we report about implementations of linear algebra algorithms on the SCS-40, which are designed to solve large dense problems using out-of-core techniques. Based on block algorithms we have implemented linear equation solvers and solvers for the generalized eigenvalue problem, which reduce I/O traffic and perform well on the vector architecture of the SCS-40. We will report numerical results and performance figures for problems with several thousand degrees of freedom.

R. E. Anderson, R. G. Grimes, A. C. Mong, and H. D. Simon

Engineering Technology Applications Division Boeing Computer Services, M/S 7L-21 P.O. Box 24346 Seattle, Washington 98124-0346

 $\frac{\#29}{A\ Parallel}$ Matrix Inversion Algorithm for Triangular Toeplitz Systems

In this paper, a parallel matrix inversion algorithm for n-order triangular Toeplitz systems is given. The number of processors needed by this algorithm is P=n, the number of parallel steps is $T = 0(\log^2 n)$. Thus, the speed up is $S = O(n/\log n)$. On the other hand, the fast division of polynomial has also corresponding parallel processing, and an obvious expression of the triangular Toeplitz matrix inversion is presented. Theorem 1. Suppose A is a n-order triangular Toeplitz matrix, then A can be calculated in $T=0(\log^2 n)$ time steps and needs at most P=n/2 processors. Corollary 1. The triangular system of equations Ax=b, where A is a triangular Toeplitz matrix of order n, can be solved in T=0(\log^2 n) time steps and needs P=n processors. Corollary 2. The division of two polynomials f(x)/g(x), where deg(f(x))=2n-1, deg(g(x))=n, can be calculated in $T=0(\log^2 n)$ time steps and at most P=n processors.

You Zhaoyong Li Lei Department of Mathematics, Xian Jiaotong University Shan Xi, China

#30
Two Supernodal Implementations of General
Sparse Factorization for Vector Computers

The introduction of hardware gather/scatter on recent vector supercomputers, such as the Cray X-MP, has made it possible to vectorize the computationally intensive kernels at the heart of general sparse Cholesky factorization (an indexed SAXPY, or so-called SAXPYI). By exploiting the supernode structure of the factor matrix, we are able to organize two new implementations around double-loop and even triple-loop computational kernels, whose speed may be greatly increased with standard loop unrolling techniques. The new implementations are generally twice as fast as an optimized sparse factorization module from SPARSPAK, with overall Computational rates ranging from 60 to 85 MFLOPS on several structural analysis matrices.

C. Cleve Ashcraft John G. Lewis Barry W. Peyton

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#33

Recursive Doubling Algorithm for Solution of Tridiagonal Systems on Hypercube Multiprocessors

The recursive doubling algorithm as developed by Stone solves a tridiagonal linear system of size n on a parallel computer with n processors using $O(\log n)$ parallel arithmetic steps. We give a limited processor version of the recursive doubling algorithm which solves the tridiagonal linear system using 0(n/p + log p) parallel arithmetic steps on a parallel computer with p processor where p<n. The main techniques rely on fast parallel prefix algorithms, for which we describe an efficient mapping on the hypercube using binary reflected Gray code. We show that for $p \not \ll n$ the recursive doubling algorithm solves the tridiagonal linear system with linear speed-up and constant efficiency over the sequential LU decomposition algorithm. The experimental results obtained on the Intel iPSC/d5 hypercube multiprocessor are presented. These results indicate greater gain in speed-up and efficiency over the existing sequential algorithms than our theoretical limits predict.

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#37
Matrix Processing Computers--The Logical Step into the Future

Matrix processing methods of computation generalize the fundamental concepts of scalar and vector computers. Naturally occurring multi-dimensional arrays of data form the basic operands for the matrix processing kernel operations. Control information, data transfer, processor organization, and memory management deal with matrices rather than with scalars. The multi-dimensional parallelism inherent in matrix processing can be effectively exploited in developing novel computer architectures for high-speed scientific computations.

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#39
Parallel Implementation of a Neumann-Dirichlet
Preconditioner in Domain Decomposition.

Alternating Neumann-Dirichlet subregions with intersecting separators are considered. Comparison of a domain decomposition method in parallel and sequential implementation is discussed. Investigation of computationally optimal number of processors on hypercube architecture is presented.

Wlodek Proskurowski Majid Haghoo Department of Mathematics, DRB 306 University of Southern California Los Angeles, CA 90089-1113

#40

The Network Emulation Problem

Loosely speaking, the network emulation problem asks how best to represent the processor adjacencies of one network with those of another. In general, we are willing to incur additional processor redundancy as long as we can maintain constant communication costs. Herein we address several questions about emulation, including: Which networks can be emulated in the plane? How much redundancy is necessary? How difficult is it to determine whether one network can emulate another?

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Donna J. Brown Coordinated Science Lab University of Illinois Urbana, IL 61801

Michael R. Fellows Department of Computer Science University of New Mexico Albuquerque, NM 87131

Michael A. Langston Department of Computer Science Washington State University Pullman, WA 99164-1210

#44

Performance Evaluation for Stochastic Dynamic Programming on Vector Multiprocessor

Performance has been measured for predictor corrector methods applied to the PDE of stochastic dynamic programming on the Argonne National Laboratory ACRF Alliant FX/8 vector multiprocessor. The dependence of the performance on the number of processors and the number of mesh points are given, as well as the relation to the vector length and memory size. The code has a large Amdahl fraction for both vectorization and parallelization.

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#45

PARALLEL RUNGE-KUTTA METHODS FOR ORDINARY DIFFERENTIAL EQUATIONS.

We investigate Runge-Kutta methods for ordinary differential equations, suitable for parallel processing across the method.

Classical explicit methods are not suitable for parallel processing. For non-stiff equations the best candidates seem to be fully implicit methods of high order constructed by collocation. We use fixpoint iteration and as starting values for the internal vectors we use a high order interpolant. A result is found giving the order of an iterate as a function of the order of the predictor and the number of iterations.

Numerical experiments on Cray X-MP/2 and Alliant show that we achieve the predicted speed up within reasonable limits. We also show that an ad-hoc constructed parallel variable step method performs better than the well-known code RKF45 on examples where the function evaluations are compute - intensive.

For stiff equations we construct efficient diagonally implicit RK methods and compare them with a well-known singly diagonally implicit method.

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Dr. Ivar Lie Institute for numerical mathematics University of Trondheim 7034 Trondheim Norway

#47 Scientific Computation in FORTRAN BX

Earlier this year the ANSI FORTRAN Committee completed its work on the Fortran 8X Standard and is in the process of submitting that standard to public review before its adoption by ANSI and ISO as the new standard.

Examples of the use of Fortran 8X in scientific computation will be given which illustrate the four main areas in which Fortran 8X extends beyond (its subset) Fortran 77: simplified array processing, generalised numerical precision, and convenient definition by the programmer of new types and new operations on those types.

Special emphasis will be given in the examples to the use of Fortran 8X in parallel algorithms.

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#50 Simulation of Larger Dimensions on the Intel iPSC Hypercube

A library of subroutines has been written and implemented on the Intel iPSC Hypercube which allows a programmer to simulate Hypercubes of higher dimensions. This subroutine library runs on a 5-dimensional MX, Hypercube and it is capable of simulating up to 9-dimensions. A special Gray code was developed to minimize communication overhead for 2-D and 3-D finite difference calculations. Tests which were run on a 2-D finite difference code are presented. The tests show that under certain conditions programs run faster when higher dimensions are simulated.

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#51 Parallel Implementation of a New Linear Matrix Solver

An Intel iPSC Hypercube implementation of a new stationary iterative method developed by Dr. John de Pillis is presented. This algorithm finds the solution vector \mathbf{x} for the invertable \mathbf{n} \mathbf{x} \mathbf{n} linear system $\mathbf{A}\mathbf{x} = (\mathbf{I} - \mathbf{B})\mathbf{x} = \mathbf{f}$ where A is a normalized matrix with real spectrum. The solution method coverges quickly because the Jacobi iteration matrix B is replaced by an equivalent iteration matrix with a smaller spectral radius. The parallel algorithm partitions A row-wise among all the processors in order to keep memory load to a minimum and to avoid duplicate computations. This method reduces only the outer loop in the iteration, which allows for exploitation of vector hardware. Example problems and timings will be presented.

Lisette de Pillis and Johnny Petersen Department of Computer Science Chr. Michelsen Institute Fantoftvegen 38 N-5036 Fantoft NORWAY

John de Pillis Bergen Scientific Centre Allégaten 36 N-5000 Bergen NORWAY Parallel Implementation of Fourier
Pseudospectral Methods on a Loosely
Coupled Array of Processors.

A domain decomposition algorithm was applied for the parallel performance of two dimensional Fast Fourier Transform (2-D FFT) in order to be used in the parallel solution of time dependent partial differential equations using pseudospectral methods. The algorithm was based on a hardware configuration consisting of shared bulk memories connected to attached processors. pseudospectral method has been implemented in parallel using data sorting for the 2-D FFT, and results have been obtained by solving a twodimensional time dependent parabolic differential equation. It was found that the parallel efficiency of the method is satisfactory for reasonable problem sizes.

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#54 Optimal Algorithms for Parallel Matrix Multiplication on Vector Computer

In this paper, two parallel algorithms for n-order matrix multiplication on vector computer are given. (1) Optimal inner product algorithm: the number of processors $P = n^3/(\log n)$; the number of parallel steps $T = 3\log n + (\log n)^2/3 - \log\log n + 0(1)$; efficiency E = O(1); (2) Optimal parallel strassen algorithm: the number of processors $p = n^2 \cdot 8^1/(\log n)$; the number of parallel steps $T = O(\log n)$; efficiency E = O(1). This two algorithms all achieve lower bound of complexity of parallel matrix multiplication in order of magnitude, and bound of the parallel processor attain optimal with the efficiency is O(1). The inner product algorithm whose efficiency is $O(1/\log n)$ and parallel strassen algorithm are improved.

You Zhaoyong Li Lei Xian Jiaotong University China

ABSTRACTS: CONTRIBUTED PRESENTATIONS

#55 Exploiting Multilevel Parallelism in PDE Software on the CRAY X-MP

This presentation provides an overview of how different types of parallelism in PDE software can be exploited on the CRAY X-MP with vectorization and the two available modes of multitasking: microtasking and macrotasking. A description is given of characteristics of these multitasking mechanisms and of experience in using them to enhance performance of two large codes applied to compressible full potential flow problems. The codes are (1) the general purpose package PLTMG of R. Bank et. al. which incorporates adaptive refinement of unstructured triangle meshes and multigrid solution techniques, and (2) a code utilizing logically uniform structured grids and preconditioned GMRES iterations. Various factors affecting multiprocessing performance are discussed.

Michael Bieterman Applied Mathematics Engineering Technology Applications Boeing Computer Services P.O. Box 24346, MS 7L-21 Seattle, WA 98124-0346

#57 An Extension of Ryckaer's Algorithm

A numerical algorithm integrating the Cartesian equations of motion for a system of interacting point particles subject to holonomic constraints is formulated. The algorithm is based on the original Ryckaert's method. However, it can be conveniently applied to the molecules in which some centers of force are coincident with very small masses, or are not coincident with any

Sheng-Bai Zhu Picosecond and Quantum Radiation Laboratory P.O. Box 4260 Texas Tech University Lubbock, TX 79409

#58 Solving Very Large Dense Systems of Linear Equations

Several experiments involving a block-based variant of the LU decomposition will be described. We will attempt to find out how much of a CRAY-XMP/48 with a 256 million work SSD (a fast external storage device) can be used to solve a very large system of linear equations (orders exceeding 10,000.) We hope to equal or exceed the computation rate of 718 Mflops reported by

Dongarra and Hewitt (January 1986 issue of SIAM J. Sci. Stat. Comput..) As time and machine availability permit, we will also report on similar experiments attempted on a CRAY-2.

Michael Pernice National Center for Atmospheric Research Scientific Computing Division P.O. Box 3000 Boulder, Colorado 80307-3000

#59 Low-Accuracy, High Speed Methods in the Numerical Solution of Ordinary Differential Equations

This paper discusses some new approaches to the numerical solution of ordinary differential equations utilizing parallel architecture. Historically, many problems have required small stepsizes and unacceptably slow integration times. The central idea of this paper is to predetermine h, the stepsize, and based on the speed of the computer, optimize the method for that given step size. Real time integration can be achieved, but the cost is a decrease in accuracy. The approach is to maximize the speed while sacrificing the least possible accuracy. While accuracy may be reduced, other properties may be preserved, for example, stability.

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Jean Michel Favre, Tektronix Inc., PO Box 500, 14150 Karl Braun Drive, Beaverton, OR, 97077

#60 Application of Supercomputers in MD study of Liquid Phase Isomerizations

With the aid of supercomputers and vectorization technique, we are able to perform full MD simulations for isomerizations in liquid. Some new features have been discovered. The entropy effect becomes weaker by tightening the freedom of motion for bond angle bending and bond stretching, and becomes negligible for the rigid isomers. The memory kernal and the rotational kinetic energy of the isomer may depend on the reaction coordinate and the barrier height. This dependence is related to the solvent viscosity and to the mass ratio of the isomer. Consideration of these effects may improve the Kramers-Grote-Hynes theoretical rates.

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#61

Fortran 9y: Extensions to FORTRAN for supporting network connected computers

The FORTRAN 8x Standard seems well-suited for a single vector processor. However, the Standard does not address issues relating to architectures like hypercubes.

There are two major, problem areas to address. One is the problem of communications between and control of (C²) independent processes. The second is the programmers' "virtual" program - the program seen as a whole.

We describe our fundamental model of network connected computer systems. This model is used to define the C^2 issues. We propose specific primitives based on our FORTRAN implementation on the FPS T-series. We indicate some ideas for defining virtual programs.

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#62

Parallel Algorithms for Market Equilibrium
Problems

In this paper we introduce both nonlinear and linearized parallel algorithms for the solution of a variety of market equilibrium problems formulated as variational inequality problems. These algorithms can be used to solve equilibrium problems such as the general spatial price equilibrium problem and the related constrained matrix problem. We provide computational results for the parallel algorithms and compare their performance to nonlinear and linearized serial algorithms. The computer utilized for this purpose is the IBM 3090 at the Cornell National Supercomputing Facility.

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#63

Concurrent I/O System for the Hypercube Multiprocessor

One of the deficiencies of existing parallel machines is the lack of adequate, high-bandwidth, parallel I/O facilities. Because of this, many interesting and important problems (image processing, data base) cannot be efficiently solved on parallel architectures. In this paper we will present the Concurrent I/O system (CIO) for the Hypercube multiprocessor. Parallel I/O systems differ from conventional I/O systems. Among challenging problems are efficient coordination of processors operating on an I/O resource (e.g. file), design of a distributed naming convention, design of the distributed cache, and a general facility for implementing I/O operations. We will discuss fundamental issues of I/O systems for parallel machines, describe the architecture of CIO, describe the CIO remote procedure call facility, and present some performance issues related to the cache design and routing strategy.

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#193

A Pseudo-Spectral Method Implemented on a SIMD Machine

In this paper we compare how effectively a vector computer and a SIMD machine solve a PDE using spectral methods. We solve Burger's equation using an implicit Chebyshev pseudo-spectral technique and optimize the method on both machines. We utilize parallel algorithms for a tridiagonal solver, FFT, and the Chebyshev recursion and compare run times on a Cray-l and a Connection Machine I. It is found that the method runs in NlnN time on the Cray-l and lnN time on the Connection Machine I, but the Cray-l still obtains the solution in a shorter time.

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#L-1
Parallel Multigrid Algorithms in Computational Fluid Dynamics

We will present progress on an ongoing project at RIACS/NASA Ames on parallel multigrid algorithms with applications to computational fluid dynamics (CFD). The long term goal of the project is to determine whether parallel computers and algorithms can make a significant contribution to the field of CFD in the near future. Our approach has been both theoretical and experimental — varying from studying model problems to actual parallel implementations of CFD codes. In particular, we wil try to speak on our experience in implementing a transonic flow code (FLO52) on the Intel iPSC hypercube.

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Ray S. Tuminaro
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#L-2

Matrix Decomposition on Message passing Architectures

We study various algorithms for solving linear systems of equations on different models of parallel architectures: Gauss, Jordan and Givens decompositions on message passing multiprocessors. We first recall some results concerning the shared data model which can be considered as a complete network of processors. In case of message passing model, the processors communicate through processor to processor channels. Y. Saad introduced a parallel Gaussian elimination algorithm on a linear array of processors based on a pipelined broadcasting of the pivot row. We derive a new algorithm by transmitting various pivot

rows between neighbouring processors. We show the asymptotic optimality of this algorithm. The preceding technics are applied to derive new parallel algorithms for the Gauss - Jordan method and the orthogonal decomposition using Givens rotations. Lastly we discuss various ways to incorporate pivoting strategies for numerical stability reasons and present results of execution on a 16 nodes Hypercube FPS T-20.

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#L-3

A Linear Processor Network for the Knapsack Problem

A linear processor network for solving the Knapsack Problem is presented. The parallel algorithm is based on a sequential algorithm from Horowitz & Sahni (E. Horowitz and S. Sahni, "Computing partitions with applications to the knapsack problem", J. of the ACM, vol 21, nº 2, april 74, pp 277-292) which is a depth-first branch and search one. Given n positive integers $W = (w_1, w_2, ..., w_n)$ and a positive integer M, the Knapsack Problem consists in finding an ensemble I \subseteq [1,n], such that $\sum w_i = M$, for $i \in I$. Each processor (PE) receives as input a combination c of $\{0,1\}^n$, its sum S(c) and a position i. Each PE stores n, the vector W and the searched value M. The basic arithmetic operation is a binary addition performed on the input for generating a new combination, the evaluation of the new sum which requires only one integer addition and the test on the equality between M and the generated sum. Strategies for pruning the search tree and for stopping the network are discussed, as well as how the parallelism is introduced into the starting algorithm. Some results of simulation on the FPS T-20 4-Cube Parallel Computer are presented.

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WEDNESDAY, DECEMBER 2

5:00 - 7:00 PM Pasadena Exhibit Hall

POSTER SESSION 2

#65 Fast Fourier Transforms on the BBN Butterfly Parallel Processor

We analyze the computation of one-dimensional FFT's on the Butterfly Parallel Processor. Our analysis deemphasizes the mathematical transformations which could be used to improve performance, and focusses instead on programming techniques. In particular, we describe the use of the Uniform System library to perform storage and task allocation for the scheduling-on-pairs FFT algorithm.

William Celmaster

Bolt Beranek and Newman Advanced Computers Inc. 10 Fawcett Street, Cambridge, MA 02238

#66 SCIENTIFIC COMPUTER PERFORMANCE EVALUATION:

A GUIDE FOR THE PERPLEXED

It is easy to compare the speed of <u>scalar</u> computers, but no recognized standards or methods exist for comparing the speed of <u>advanced</u> <u>scientific</u> computers. Typically, a user would like to know:

- On average, how fast is a given computer on the class of problems it is expected to run?
- 2. How reproducible is this average? The speed of vector processors is notoriously sensitive to how much of the workload has vector structure in both program and data.

We will present some answers to these questions, based on real application benchmarks, spreadsheet databases and nonparametric statistics. The methods used (jackknife and bootstrap) give means and variances for computer speed parameters, and we will explore speculations about a "modern portfolio theory" of computer performance based on these means and variances.

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#67

Fast Banded Matrix Multiplication on Parallel Computers

A fast algorithm for multiplying banded matrices is presented. The algorithm's implementation can be several orders of magnitude faster than some commercially available subroutines for banded matrix multiplication on vector and parallel computers. It can also be one order of magnitude faster on scalar machines. The performance of the method on various vector and parallel computers (including the ETA¹°, IBM 3090, Alliant FX/8, and Cyber 205) is given and compared. (This algorithm was developed for use in matrix inversion. See the accompanying paper by Magnan for a discussion of this application.)

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#68 Parallel and Vector Processing of the Incremental Dynamic Programming Algorithm

This paper discusses modifications to the incremental dynamic programming (DP) algorithm which allow the efficient solution of ten-dimensional problems on a supercomputer. These modifications include (1) parallel processing of the outer-loop over the state variables, (2) vector calculation of the inner-loop minimization over the control variables, and (3) the use of control vectors to accomplish the feasibility tests.

Timing comparisons are made for the algorithm running on a CDC CYBER 205, ETA-10, CRAY X-MP, ALLIANT FX/8, and a VAX 780. The CPU times show the ETA-10 to be over 4 times as fast than the CRAY X-MP for 64-bit precision. Timings on the ALLIANT are in progress with multitasking to split the work to be done in the stage loop among the eight processors.

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Dennis Morrow, Research Associate and Mohan Ramamurthy, Research Associate Supercomputer Computations Research Institute, 523 Keen Bldg., Tallahassee, FL 32306-4052

#69
Simulation of Transonic Flow on a
Concurrent Computer

With the increasing use of computational fluid dynamics for aerodynamic design, the demand for cost-effective computers has also increased. The transonic flow regime is one area of study that is important for all supersonic and high subsonic aircraft. This work involves the implementation of the widely-used FL057 code, an Euler equation-based transonic flow solver, on the AMETEK local-memory concurrent computer. As a result of various factors in the design of the original program, the conversion process is straightforward, and the concurrent version is shown to run very efficiently. The paper discusses the implementation philosophy, as well as the memory requirements, verification of results and performance, including speedups.

Mark E. Bassett Christopher J. Catherasoo AMETEK Computer Research Division 610 North Santa Anita Avenue Arcadia, California 91006 (818)445-6811

#71 <u>Distributed Orthogonal Factorization</u>: Givens and Householder Algorithms

We design and implement several algorithms for orthogonal factorization on distributed memory multiprocessors (e.g. hypercubes). Each algorithm is analyzed for its arithmetic and communication complexity. Two of the algorithms, the standard Givens sequence, and the greedy Givens sequence, are based on Givens rotations. Two algorithms employ Householder transformations but with different communication schemes (broadcast and pipelined ring). A fifth algorithm is a hybrid: it uses Householder transformations in an internal phase in which there is no communication, and Givens rotations in a recursive elimination phase. This algorithm possesses the low arithmetic cost of the Householder algorithms and the low communication cost of the Givens algorithms. Computation times from implementations on an intel IPSC hypercube agree quite well with our theoretical analyses.

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#72 A Constraint Algorithm for Maintaining Rigid Bonds in Molecular Dynamics Simulations of Large Molecules

An algorithm is described for efficiently maintaining rigid bonds in molecular dynamics simulations of polyatomic molecules. In this algorithm, a leapfrog integration scheme with constraint conditions is applied to all the bound particles. The procedure for computing the constraint forces requires a small number of iterations. The bond lengths do not have to be kept exactly fixed, but can deviate from their assigned values by small amounts. The algorithm is similar to the SHAKE algorithm and the algorithm of Memon et al. in that it iteratively corrects the constrained equations of motion. However, instead of applying successive corrections to particle positions or inverting a matrix, our method iteratively corrects the constraint forces required for maintaining the rigid bonds.

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#75 Pattern Recognition by Neural Networks on Hypercube

The objective of this work is to compare the performance characteristics of the parallel implementations of the Hopfield's model and the "back-propagation" model (Rumelhart et al.) in pattern recognition, and to investigate the transition from one model to the other. Specifically, the test case of recognition of Chinese characters is studied on hypercube topology computers. Quantitative analyses of the relationship between the amount of overlap in the training samples and the reliability of both models is presented. Also, the smooth transition linking the two models is discussed. The interpolating network has higher capacity than in the Hopfield's model and learns faster than in the "back-propagation" model.

Alex W. Ho, Wojtek Furmanski Concurrent Computation Project Mail Stop 206-49 California Intstitute of Technology Pasadena, CA 91125

#77

A Parallel Tridiagonal Equations Solver for the Butterfly Multiprocessor

A parallel block partitioning method to solve a tridiagonal system of linear equations is adapted to the Butterfly multiprocessor. A performance analysis of the programming experiments on the 32-node Butterfly is presented. The computational results verify the theoretical speedup and efficiency results of the parallel algorithm over its serial counterpart. Also, included is a comparative study of the performance runs of the same code on the Butterfly processor with hardware floating point versus the one with software floating point facility. The achieved results are shown to be within 82-90% of the predicted performance.

Swarn P. Kumar

Boeing Computer Services Engineering Technology Applications P.O. Box 24346, M/S 7L-21 Seattle, Washington 98124-0346

#78
Electromagnetic Scattering Analysis on the Hypercube

To investigate the applicability of a parallel computing architecture to the solution of large scale electromagnetic scattering problems, two existing sequential electromagnetic codes utilizing different algorithms have been converted to run on the Jet Propulsion Laboratory/California Institute of Technology Mark III Hypercube. The first was a frequency domain method of moments approach developed at Lawrence Livermore National Laboratory, "Numerical Electromagnetics Code" (NEC-2). It requires the solution of a large system of linear equations. The second was a time domain finite difference solution of Maxwell's equations for scattered fields. This presentation will describe the algorithms used in the parallel decompositions and discuss the performance results.

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#79

A Performance Comparison of Three Mini Supercomputers: the Alliant FX/8, the SCS-40, and the Convex C-1

A benchmark of three leading mini supercomputers has been carried out. Using the Los Alamos Standard Benchmark Set, a series of tests were run on each of the three machines. The results are compared with each other and with the same set of tests run on the Cray Research, Inc. K-MP series machine. In addition, the asynchronous parallel processing capability of one of the machines is explored in some detail. The results suggest a strong showing for this class of computers.

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#80

Multitasking a sparse matrix code on the Alliant FX/8

We have recently multitasked a code for the direct solution of sparse linear equations on the Alliant FX/8. We discuss several issues which are involved, all of which are of relevance to any shared memory multiprocessor. Among these issues are the dynamic allocation of data, the management of task queues, task spawning, and the effect of We also show controlling the granularity. runs of our code under the Schedule package from Argonne which presents a portable interface to users of parallel machines, allows the user to define his computational graph, and has very useful graphic output to a SUN workstation. Our tailored code attains a speed-up of about 6 on the eight processors of the Alliant. We suggest ways of improving it further.

Iain S. Duff Computer Science and Systems Division Harwell Laboratory, Oxfordshire, England OX11 ORA

#81

Parallel Triangular Matrix Inversion

The speaker will show that the information flow in parallel triangular matrix inversion is virtually identical to that in parallel Cholesky factorization. So, parallel triangular matrix inversion algorithms can be found by modifying parallel Cholesky factorization algorithms. Drawing from the extensive literature on parallel Cholesky factorization, the speaker will give parallel algorithms for triangular matrix inversion on systolic arrays, array processors, shared-memory multiprocessors and data flow machines.

David C. Fisher, Harvey Mudd College, Claremont, CA 91711 (Mathematics Dept.)

Linda Sattler, The Claremont Graduate School, Claremont, CA 91711 (graduate student in the Mathematics Dent.)

#82 Parallel Processing Bibliography (in Machine-Readable Form)

Rather than print yet another bibliography, the author is doing something different. Hardcopy bibliographies are just that: HARD! This means 1) they don't change, they are not 2) correctable, 3) updateable, or 4) reformattable for different citation styles, and 5) searching becomes difficult when bibliographies get too large. Soft bibliographies (in machine-readable form) are updateable, correctable, are reformattable, and search very quickly. It offers control librarians cannot offer.

The author started with several existing parallel processing bibliographies (including annotations) and has maintained this list for three years. The bibliography is available free of charge to ARPAnet sites (with updates) or for a nominal fee from COSMIC, NASA's software distribution facility. It contains over 7,000 references regarding parallel processing hardware, software, algorithms, and other useful citations. SIAM members are encouraged to contribute. Key word fields note surveys, required, and recommended readings for beginners.

E. N. Miya

Computational Research Branch NASA Ames Research Center Moffett Field, CA 94035 eugene@ames-nas.arpa UUCP: ames!eugene #85
Analysis of vectorisable preconditioners
for 3D problems

Incomplete factorisations of matrices deriving from 3D problems can be based on point, line or plane pivots. Numerical tests will be presented showing that, while recursive methods seem to display largely the same behaviour in all cases, strategies to obtain vectorisability (mainly by explicitly using approximations to inverses of pivots) will result in widely varying convergence speeds (including non-convergence!). Line methods appear the most interesting in terms of numbers of iterations/operations (and with vectorlength h-2/2 on average). A limit matrix analysis on sever strongly anisotropic model problems will analyse the difference between convergence of recursive and vectorisable methods.

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#86
Decomposition and Solution of Transient
Finite Element Problems on Hypercubes

Current research demonstrates that concurrent computers with the distributedmemory, hypercube MIMD architecture can be highly effective machines for finite element applications. Our research applies domain decomposition methods and concurrent equation solvers to examine performance issues for realistic structural and transient solid mechanics problems. The relative merits of iterative (conjugate gradient and Lanczos) and direct or hybrid solvers for a range of problems will be discussed, as well as the effective solution of a sequence of righthand sides. High concurrent efficiency and transparency at the user level are realized in application on the JPL/Caltech Hypercubes.

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Bahram Nour-Omid Mechanics and Materials Engineering Lab. Lockheed Palo Alto Research Laboratory Palo Alto, CA 94304

#87 Serial-to-Parallel Conversion of Real Time Continuous System Simulations

Many existing computer simulation programs of real time systems are amenable to parallel execution on local-memory, message-passing architecture parallel processors. We have developed a set of tools which partition an existing simulation program along module boundaries and map the resulting system to a hypercube parallel processor. Our conversion approach seeks to balance the computational load, minimize communication costs and maximize the number of processors employed.

Alan M. Baum Donald J. McMillan Computer Science Department General Motors Research Laboratories Warren, Michigan 48090-9055

#88

Design of a Real Time, Interactive Simulation
Computer with a Parallel Processing Organization

A special purpose computer organization is proposed for fast, real time, interactive simulation. The organization comprises a central processor, tens to thausants of identical peripheral processors (PP's), specially organized memory modules and a set of registers, achieving MIMD (Multiple Instruction stream, Multiple Data stream) parallel processing. Data are transfered by and among the memory modules, rather than by and among PP's, at the speed of memory cycle. A simple example of such organization will be discussed for a simulation of a typical dynamic phenomenon.

Yukoh Kobayashi, Ph.D. California Research and Development 1118 S. Lake Street Los Angeles, CA 90006

#92

Parallel Subspace Iteration for Solving the Generalized Eigenvalue Problem

Subspace iteration is used to compute m eigenpairs of a generalized eigenvalue problem, KX = λ MX, of size n, n >> m. A parallel execution time model of subspace iteration broken into basic phases of the computation, e.g., inverse iteration and Rayleigh-Ritz approximations, is defined. Model parameters account for computation and communication costs, synchronization and parallel processing overheads, both within and between phases.

Results of a parallel implementation of subspace iteration on a shared memory MIMD machine are presented. These results are used to test the validity of the execution time model.

Umesh Mahajan Merrell Patrick Department of Computer Science Duke University Durham, NC 27706

#93 Performance of the NCUBE Hypercube

The objective of this work is to study the performance characteristics of the NCUBE hypercube, and to investigate the use of the NCUBE parallel graphics interface. This talk also addresses the performance of the high-speed CrOS III crystalline operating system on the NCUBE. As a test problem we chose to study two-dimensional, convectively-dominated, fluid flows using a finite-difference technique which incorporates the flux-corrected transport algorithm. For this algorithm a speed of ~90 Kflops per node is attained. In designing algorithms for the NCUBE it is important to take into account the high message latency of the native VERTEX operating system.

Geoffrey C. Fox, David W. Walker Concurrent Computation Project Mail Stop 206-49 California Intstitute of Technology Pasadena, CA 91125 Gary R. Montry Sandia National Laboratories Albuquerque, NM 87185

#95 CONVERGENCE OF CERTAIN PARALLEL SOURCE ITERATION SCHEMES IN PARTICLE TRANSPORT

Combined inner/outer source iteration is a well-developed and widely implemented scheme for the serial solution of (linear) particle transport problems. Brickner, Hiromoto and Wienke (of LANL) have presented several parallel variants of inner/outer source iteration, and results from implementations of these upon a hypercube (as well as a shared-memory machine). These parallel source iteration schemes all correspond to decomposition by energy group.

We present a general model of parallel source iteration that encompasses all schemes in the work cited above. Matters related to convergence of this general class of parallel source iterations also are discussed.

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Beny Neta Department of Mathematics Naval Postgraduate School Monterey CA 93943

#96
Study of Program Termination Using Z-Transform
Theory

This paper deals with the termination of programs. A method of generating a set of recurrence equations for a given program is explained. Z-transform techniques are applied on these equations and the resulting equations in z are analysed to study the termination of the programs. This is a new approach.

The left shifting property is used for Conversion from discrete form to Z-transform. Then using the well known theorems and inverse transform the response of the equations is estimated. From the response termination is decided.

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- S. Arumugam Computer Centre, GCT Coimbatore-641013 INDIA

#97
Parallel Distributed Processing Models for
Natural Language Inference

Parallel procedures are defined on distributed network representations of natural language sentences that identify their logical and probabilistic consequences. These procedures are efficient on large natural language data bases and resolve computationally to the parallel identification of requisite patterns of interconnections among distributed occurrences of certain basic elements of which all network representations are wholly composed.

The proposed network representations can be systematically modified to yield a spectrum of parallel distributed processing models for natural language inference, ranging from logically complete models that yield only correct inferences to logically incomplete models that also yield certain incorrect inferences.

Peter G. Tripodes Department of Mathematics Los Angeles City College Los Angeles, CA 90029

#100 A Methodology for Dynamic Scheduling of Computation Tasks in Parallel Processing Systems

The use of parallel processor computer systems for high-performance scientific computing today requires systems designers to partition an appli-

cation into functions for each processor allocating spare processors for fault tolerance. A methodology is presented that avoids this static partitioning by dynamic assignment and reassignment of application tasks to processors. The computational intractability problem of optimal multiprocessor task assignment is circumvented through the use of heuristic and sub-optimal scheduling algorithms. The development and simulation results of two major categories of algorithms are presented including a centralized, computational-level scheduling approach, and a distributed, general multiprocessor ready-que oriented approach.

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#101

Analysis of Fortran and C Performance on the FPS T Series Parallel Supercomputer.

Benchmark Fortran and C programs have been written for the FPS T Series to evaluate the performance of these languages on this highly parallel scientific computer. Benchmarks were chosen which span a wide cross-section of computational parameters and were run on configurations of up to 128 nodes. Algorithm specification and detailed performance information are presented for each of the benchmarks, and the implications for general performance expectations are discussed.

Thomas E. Bauer and Ian A. Taylor Floating Point Systems, Inc. 3601 S.W. Murray Blvd. Beaverton, Oregon 97005

#102 Application of Book Thickness of Graphs to Parallel Computing

This paper applies the notion of book thickness to VLSI design and to the parallelization of sequential code. The book thickness of a graph is the smallest number of pages in a "book" in which the graph can be drawn with vertices along the binding. If the I/O ports of a chip, or smaller processing element, are regarded as vertices with computational paths as edges, then book thickness corresponds directly to time-division-multiplexing of the calculation. The pages of a program flow graph can be implemented in Mealy-type processing elements; an additional set of Moore-type PE's supervise the calculation.

Paul C. Kainen The Analytic Sciences Corporation 8301 Greensboro Dr. McLean, VA 22102

#106

Parallel Optimization of Large-Scale Linear and Nonlinear Networks

As a result of the availability of multiprocessors and multicomputers, there has been a renewed interest in decomposition techniques for largescale network optimization problems. In network optimization the granularity of the decompositions that are possible ranges from very coarse (corresponding to collections of network optimization subproblems) to very fine (in the sense of vector arithmetic). These topics will be addressed in the context of research on large-scale generalized networks and nonlinear networks. We will describe computational experience with two parallel computing systems at the Computer Sciences Department of the University of Wisconsin and possibly with other advanced architectures as well.

Robert R. Meyer University of Wisconsin-Madison Computer Sciences Department 1210 West Dayton Street Madison, Wisconsin 53706

#110 Pipeline SOR: A Parallel Iterative Method for Shared Memory Machines

We present a new parallel version of SOR called Pipeline SOR (PiSOR). This algorithm is designed to make efficient use of processors on a shared memory machine while maintaining the traditional serial SOR convergence properties. Iterations in the PiSOR algorithm are performed by different processors and are pipelined one after the other. Synchronization flags are used to prevent any iteration from overtaking the ones ahead of it. Results will be given from an implementation of the PiSOR algorithm on the Balance 21000.

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#111 Optimization Strategies for Pattern Selection Problems and Their Implementation on the BBN Butterfly Parallel Processor

We have used optimization techniques for finding steady solutions to dynamical systems for which Lyapunov functionals can be constructed. These algorithms have been implemented on a parallel processing computer, and we have compared the efficacy of such calculations with what can be accomplished by direct integration on a serial processor machine. The particular problem addressed is pattern

selection in Rayleigh-Bénard convection, for which the corresponding dynamical model equation is the Swift-Hohenberg equation. We conclude that since the dynamics induced by the various minimization algorithms leads to identical steady-state solutions, and as the solution time is a very strong function of the dynamical process used to move along the Lyapunov "surface", the construction of efficient dynamical processes on the Lyapunov "surface" is an effective method for finding steady solutions to the underlying dynamical system.

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William Celmaster
Eric Tenenbaum
Bolt, Beranek and Newman ACI
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#113 Efficient Implementation of Parallel 3-D

Seismic Algorithms on the CRAY X-MP

The recent need in 3-D seismic processing both in forward modeling and in migration has fostered the growth in algorithm development on the current generation of vector and/or parallel computers. Computation of meaningful sized problems requires extensive calculations and large 3-D data which must be retrieved and restored several thousand times during the computation.

With the proper choice of algorithms, problem sizes previously thought intractable can now be solved within reasonable time. This concept is demonstrated on the CRAY X-MP/48 with a 128MW SSD with examples of 3-D acoustic and elastic forward modeling using the Fourier method and a 3-D post stack depth migration.

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Chippewa Falls, WI 54729 f Moshe Reshef

Dan D. Kosloff Dept. of Geophysics Tel Aviv University Tel Aviv, Israel

Moshe Resner Cray Research, Inc. 2000 Bering Drive Suite 830 Houston, TX 77057

Christopher C. Hsiung

Cray Research, Inc.

900 Lowater Road

#114 Sparse Matrix Solution for Circuit Simulation on Multiprocessors

Motivated by the need to speed up the computationally intensive process of circuit simulation, this paper presents two approaches to parallel sparse LU decomposition: one which exploits parallelism at the level of row operations and the other at the level of elementary arithmetic operations. The various factors that determine the actual achieved performance of a parallel implementation of an algorithm are identified and a general framework is developed for evaluating their relative effects. This frame-work is then used along with measured results on various circuit matrices obtained from implementations on the ALLIANT FX/8 to contrast the performance of the two algorithms.

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V. Visvanathan AT&T Bell Laboratories, Room 3A-529 Murray Hill, NJ 07974

#115 Reliable Parallel Eliptic PDE Solution

Message-based multiprocessor systems are ideal vehicles for the implementation of relaxation algorithms. However, due to hardware and/or software errors, the end result of the computation may not be reliable. This work embeds reliability, through the use of a constraint predicate, into the class of eliptic PDE relaxation solutions to both flag faulty and ensure correct behavior of the problem solution. The presentation will show both graphically and analytically the effects of faulty behavior and the action of the constraint predicate on the faults.

Bruce M. McMillin and Lionel M. Ni Department of Computer Science Michigan State University East Lansing, Michigan 48824-1027

#116 Efficient Matrix Multiplication on Hypercube Multiprocessors

Matrix multiplication on hypercube multiprocessors is studied to determine the effectiveness of partitioning, pipelining, and architectural change on the performance of algorithms. The study shows that a linear speedup can be achieved by combining the above three factors together, even the input matrices must be downloaded from a host and the resultant matrix must be transferred back. This linear speedup follows from a minimization in the communication overhead through the overlapping between communication and computation. The approach presented here can be applied to other problems to develop more efficient parallel algorithms on hypercube multiprocessors.

Chung-Ta King and Lionel M. Ni Department of Computer Science Michigan State University East Lansing, Michigan 48824-1027

#117 Parallel Simulated Annealing Optimization of Distributed Database Networks

Design of a minimum cost network using standard linear programming codes may require hours of IBM 4341 time for a network as small as five nodes. Simulated annealing can find near optimal solutions in minutes on a PC. We have implemented a parallel version of simulated annealing on Transputers by subdividing the network, optimizing the sub-networks, then optimizing the entire network. The parallel version yields more tightly grouped results and a lower average cost than the sequential version for the same computing time. The parallel version performs better if the subnetworks are not compact but overlap each other entirely.

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Philippe Julien-Laferriere Xerox Corporation Paris, France

#118

Parallel Branch-and-Bound Algorithms on Hypercube Multiprocessors

Branch-and-bound (BB) algorithms are solution techniques for many optimization problems in Operations Research. They require considerable computational effort, making them candidates for parallel processing. This paper presents preliminary work on implementing parallel BB algorithms on hypercube multiprocessors. Integer programming is taken as a case study. Initially considering the special case of 0-1 integer programming, two generic parallel implementations of a BB algorithm for solving that problem on hypercube multiprocessors are presented. The two implementations serve as a feasibility study to determine factors that affect the performance of the general case. Experimental results from the NCUBE/ten with 64 processors demonstrate the performance of the two implementations.

Tarek S. Abdel-Rahman and Trevor N. Mudge Dept. of Electrical Engineering and Computer Science University of Michigan 1101 Beal Avenue Ann Arbor, Michigan 48109-2110

#119

Parallel Programming Without Tears

The design of the Myrias parallel computing system has emphasized ease of use. Selected do statements in Fortran 77 programs are replaced with pardo statements, which create sets of tasks that run in independent address spaces that are merged at task completion. pardos can be nested and combined with recursion. The Myrias implementation assigns tasks to processing elements, balances load between PEs, and moves data between tasks, all automatically. Programs are repeatable, and can be debugged using a symbolic debugger. This paper explains pardo, and develops a particle-in-cell code as an example application.

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Multiplying arbitrarily shaped matrices on a Data Parallel Computer

The Connection Machine is a Data Parallel Computer for which we have implemented routines for the multiplication of arbitrarily shaped matrices. Hence, inner-product, rank-1, AXPY, and matrix-vector computations are covered. All these operations are asymptotically logarithmic in time if there are sufficently many processors. However, practically they may be considered constant time operations on the Connection Machine by the availability of fast global operations such as copy and summation. We present data allocation schemes, matrix partitioning for maximally concurrent computation, complete algorithms and kernel functions.

S. Lennart Johnsson, Alan Ruttenberg, Alex Vasilevsky, and Jill Mesirov

Thinking Machines Corp. 245 First Street Cambridge, MA 02142

#L-5

Computing Fast Fourier Transforms on a Data Parallel Computer

The computation of large Fourier Transforms is of great interest for several signal processing applications, the solution of partial differential equations in three dimensional space, and in crystallography. In a data parallel computer an FFT on N elements can be performed in a time proportional to log_2N with a suitable communication system, and a number of processors proportional to N. The Connection Machine is a highly concurrent computer with a large primary storage. We discribe the implementation of a radix-2, complex FFT on the Connection Machine. Two techniques for the efficient use of the communication system that are described and implemented are pipelining and cube rotation. We also address the issue of allocation and reuse of twiddle factors.

S. Lennart Johnsson, Alan Ruttenberg, Alex Vasilevsky, and Guy Blelloch

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<u>L-7</u>
A Product Integration Method for Vortex Dynamics

A new high-order method for evaluating the multi-dimensional integrals appearing in the vorticity formulation of inviscid, incompressible fluid dynamics is presented. Unlike discrete vortex methods, the singularity in the integrand is analytically reduced, and does not require the introduction of a smoothing (finitecore) function, and its associated non-physical length scale.

The support of a general compact vorticity distribution is covered by a coordinate system periodic in (at least) one direction. By evaluating the integrals in iterated fashion, the integrals over the periodic coordinate are evaluated with spectral accuracy. The resulting scheme yields good accuracy at moderate resolution, and is naturally parallelizable on medium and fine grained levels.

Michael Shelley and Moshe Israeli Princeton University The Program in Applied and Computational Mathematics 217 Fine Hall Princeton, New Jersey 08544 #L-10 A Parallel Algorithm for Matrix Inversion

A variational algorithm for inverting matrices is presented. Since the algorithm involves mostly matrix multiplication and addition it is well suited for use on vector and parallel computers. It can be several orders of magnitude faster than methods which are in common use on these computers. An advantage of the algorithm is that it can obtain an approximate, sparse inverse of a sparse matrix. In addition, it is suitable for general classes of matrices and is robust for ill-conditioned matrices. The algorithm's performance on vector and parallel computers is demonstrated with example problems. (See the accompanying paper by Bertram and Magnan for an efficient banded matrix multiplication method developed for the implementation of the inversion algorithm.)

Jerry F. Magnan Department of Mathematics Florida State University Tallahassee, FL 32306

THURSDAY, DECEMBER 3 5:00 - 7:00 PM Pasadena Exhibit Hall POSTER SESSION 3

#121
OBTAINING NONUNIFORM SYSTOLIC DESIGNS USING MODIFIED MAPPING TECHNIQUES

Several systolic solutions for matrix triangularization using Givens' rotations have been derived manually. Many techniques for automatically mapping algorithms onto systolic arrays have been studied. Most mapping techniques assume uniformity of data dependence in the source algorithm or dataflow in the target design. The manually-derived designs for triangularization using Givens' rotations exhibit nonuniform dataflow, so it is difficult to apply proposed

mapping techniques to derive known systolic designs for this problem. This study uses a systolic array design tool called SAGE to obtain systolic designs equivalent to several manually-derived designs semi-automatically. Several novel solutions are also obtained.

Keith R. Allen & Roy P. Pargas Department of Computer Science Clemson University Clemson, SC 29634-1906 #130

A Parallel Processing Pre-Processor For the FORTRAN Language

Current implementations of parallel FORTRAN languages rely heavily on library routines accessed through the standard FORTRAN subroutine calling convention. This is neither convenient, reliable, or readable, and generally has a significant effect on programming production. This pre-processor introduces new FORTRAN "operators" or keywords that more clearly define the parallel operation being used and that also makes these parallel operations more readily apparent to the casual reader of the program (and even to the programmer!). The primary project goals were to 1) provide "correct" primitive operators such as barriers, critical regions, and others so that the programmer would have basic operators that are reliable; and 2) provide as much machineindependence as possible so that the programmer does not have to be concerned with what machine he will be using.

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#132

Inhomogeneous Parallel Supercomputers

Parallel computers have not offered a sufficient advance in performance or cost-effectiveness over conventional supercomputers to inspire widespread interest among scientists and engineers with computation-intensive problems. exploiting properties of the class of target applications, parallel computers can be designed which are at least 100 times more powerful and cost-effective than those built thus far. Such computers can be built using current off-the-shelf components. The speaker will present techniques for measuring the dependence of the ratio of communication to computation for computation-intensive problems on the decomposition grain size, and will show that this quantity is not scale invariant.

Anthony E. Terrano Department of Electrical and Computer Engineering Rutgers University Piscataway, NJ 08854 #133

FAULT-TOLERANT MAPPING OF LARGE SCALE RESOURCE ALLOCATION ALGORITHMS ONTO CONCURRENT PROCESSORS

This paper is concerned with the mapping of large scale resource allocation algorithms onto parallel computing architectures. The mapping problem is viewed as one of assigning the nodes of a directed, acyclic problem graph (representing the logical and data dependencies among the tasks constituting the algorithm) onto the nodes of an undirected system graph (denoting the parallel computing architecture). The objective is to minimize the completion time of the algorithm such that the redundancy, processor memory, and security constraints are satisfied. The dealys introduced by task queueing, message collision, and task precedence relationships are explicitly modeled. We present three algorithms to solve the mapping problem. The first is a two-stage heuristic that determines the order of task execution based on the critical path method (CPM), and then employs a one-step optimization algorithm to determine the task allocation. The second uses the idea of pair-wise exchange on task execution order to improve the performance of the heuristic algorithm. Finally, we develop an optimal mapping strategy using the A* algorithm, wherein a lower bound on the cost-to-go is obtained from the CPM, and an upper bound from the two-stage heuristic. The performance of the algorithms is illustrated on a variety of examples.

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#135

Research Issues in 1000-Processor Linear Algebra Algorithms in Finite Element Analysis*

Research in parallel finite element solution algorithms on a 10-processor ELXSI 6400 has identified several promising candidates for further investigation on 1000-processor systems. These algorithms include: (1) Jacobi preconditioned conjugate gradients (PCG), (2) incomplete LU factorization PCG, (3) unaccelerated and CG-accelerated multigrid based on Jacobi overrelaxation, and (4) hybrid direct-iterative solution methods. Limitations of these 10-processor algorithms are outlined, and progress towards achieving massive parallelism on a 1024-node NCUBE hypercube is discussed.

*This work performed at Sandia National Laboratories supported by the U.S. Department of Energy under contract number DE-ACO4-76DP00789.

Robert E. Benner and Gary R. Montry Sandia National Laboratories Albuquerque, NM 87185-5800

#137

Implementation of Iterative and Direct Solvers on the UNIX CRAY X-MP24

A comparison is made of a set of linear algebra routines on a vector - parallel supercomputer, CRAY X-MP24. The set contains routines for matrix multiplication and for the iterative and direct solutions of general systems of equations. These routines are designed to utilize a secondary solid stae storage device as well as the parallel processing features of the CRAY X-MP. Performance results are presented that indicate that the speed of these routines approach that of routines with all data in main memory and is close to the maximum speed of the processor.

Michael A. Driscoll and Elissa C. Pariser AT&T Bell Laboratories 600 Mountain Avenue 07974 Murray Hill, NJ

A Multiprocessor Model for a System **Evaluation Tool**

An interactive multiprocessor evaluation project is underway and a first version of the tool, PARET (Parallel Architecture Research and Evaluation Tool), has been implemented. A system model for evaluation must accurately represent significant behavior, but be simple and flexible enough for rapid construction and to maintain an interactive execution speed.

Directed flow graphs have been used extensively to model parallel algorithms with nodes representing software modules, and arcs showing the flow of data and control. This paradigm has been extended to include the entire multiprocessor, with nodes representing all primitive units of work which must be scheduled on a physical resource for execution. The author will discuss the model and its application to PARET.

Kathleen M. Nichols Room No. 4G-622 Computer & Robotics Research Lab. AT&T Bell Laboratories Crawfords Corner Road Holmdel, New Jersey 07733

A Case Study in Scientific Parallel Processing: Prefix-Based Computation

In this presentation the speaker will compare and contrast several implementations of computational problems whose underlying parallelism can be modeled as a parallel prefix computation, and use these results to investigate the various factors that dictate the performance of parallel programs. The problem of prefix computation which is used here as a case study finds widespread applications in computing linear recurrences, matrix multiplication, generalized Horner's expressions, various signal processing and graph theoretic algorithms. Here the speaker will focus on actual implementations of prefixbased applications using a variety of parallel processors (S/NET, NCUBE), software environment (MEGLOS, LINDA) and different programming paradigms (asynchronous, synchronous).

Binay Sugla Room No. 4G604 Computer and Robotics Research Laboratory AT&T Bell Laboratories Crawfords Corner Road Holmdel, New Jersey 07733

The Effects of Parallel Processing on #141 Matrix Computations

Increased computational speed and the power to manipulate numbers simultaneously are obvious effects of parallel processing. Yet in order to use the technology to its fullest, one must gain insight into its different implementations. The purpose of this presentation is to describe the more versatile implementations of parallel processing on commonly used matrix computations: implementations on algorithms designed to compute the solution of linear systems, the solution of the linear and non-linear least squares problems, and the solution of eigenvalue problems. Also implementation of parallel processing on corrective techniques to improve the accuracy of ill-conditioned problems will be included.

Laurie Melany Darmody 9521 Baltimore Avenue Laurel, Maryland 20707 #144
The Behavior of Conjugate Gradient Methods on a Multi-Vector Processor with a Hierarchical Memory

Conjugate gradient algorithms are efficient solvers for elliptic partial difference equations. In this paper, an analysis of some of the tradeoffs involved in the design of an efficient implementation of conjugate gradient based algorithms for a multi-vector processor with a two-level memory hierarchy is presented and supplemented by experimentals results obtained on an Alliant FX/8. The considered algorithms comprise the classical conjugate gradient algorithm, preconditioning techniques that are well suited for parallel computers such as polynomial preconditioners and a vectorized version of the incomplete Cholesky preconditioner as well as the reduced system approach.

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on leave from INRIA, domaine de Voluceau – Rocquencourt, B.P.105 78150 Le Chesnay, France

#147
The Impact of Hierarchical Memory Systems on Linear Algebra Algorithm Design

In 1986, work was done at CSRD which demonstrated that linear algebra algorithms based on the BLAS or extended BLAS were not able to achieve high performance on multivector processors with a hierarchical memory system due to a lack of data locality. For such machines, it is necessary to implement 'block' linear algebra algorithms in terms of matrix-matrix primitives. The design of efficient linear algebra algorithms for these architectures requires an analysis of the behavior of the matrix-matrix primitives and the resulting block algorithms as a function of certain system parameters (cache size, number of processors ...). This paper discusses the methodology and results of this analysis.

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* on leave from INRIA, domaine de Voluceau – Rocquencourt, B.P.105 78150 Le Chesnay, France An Analysis of a Multivector Processor Implementation of Cyclic Reduction

In this paper, the architectural issues involved in implementing algorithms to solve regular sparse linear systems, such as cyclic reduction, on a multivector processor with a hierarchical shared memory system are investigated. These issues include: data locality, cache bank conflict resolution strategies and the influence of MIMD execution on the address streams. Techniques are discussed which avoid some of the well-known difficulties associated with cyclic reduction on parallel and vector processors.

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#149 An Out-Of-Core Algorithm for Solving Dense Sets of Linear Equations on the FPS-164

Solution of large and dense linear systems on commercial parallel processors has not been extensively studied in the literature. Such systems of equations arise in the moment method solution to many problems in electromagnetic theory. A LU decomposition algorithm is presented to solve large and dense sets of linear equations with complex coefficients on the FPS 164. The algorithm takes into account the necessary I/O and CPU overlapping for computing the solution. We analyze the interrelationship between the memory size, the CPU speed and the I/O transfer rate and give projections for improving the performance.

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#152

Linear Permutations on a Multistage

Interconnection Network

Multistage interconnection networks with N inputs and outputs and log N stages pass only a limited set of permutations of data. In this paper we study the routing of the class of linear permutations on a multistage network. This class contains several permutations useful in parallel processing. A simple condition for passability of permutations in this class is formulated. For permutations which are not admissible in one pass, an algorithm for conflict-free routing using multiple passes is presented.

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#156

SOLUTION OF ELLIPTIC SYSTEMS ON HIGHLY PARALLEL COMPUTERS

The fastest known algorithms for the solution of an elliptic system all have one thing in common: they require $O(\log n)$ operations if the number of processors P = O(n). They differ greatly in their performance, however, when one examines the constant in front of the $\log (n)$ in this bound. They also differ in their accuracy when roundoff noise is present and in their ability to handle singular elliptic problems. Some, including the FFT based algorithms, will not handle variable coefficient problems. The fact that they differ in their communication requirements means that their performance is architecture dependent. One of the fastest of the highly parallel algorithms is the new algorithm PSMG, for Parallel Superconvergent Multigrid, which Oliver McBryan and the author are developing. PSMG reduces the error by a factor of 0.0014 on every iteration when solving a two-dimensional Poisson problem, for example, and converges almost as fast on nearby variable coefficient problems.

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#158

Minimum Spanning Tree on the Hmesh Architecture

A fast algorithm to compute minimum spanning tree of a given undirected graph on Hierarchical MESH connected computer (HMESH) is presented. The time complexity of the algorithm is $O(\log^2 n)$, where n is the number of nodes in the graph. HMESH is a broadcast bus VLSI architecture which consists of $n \times n$ processing elements (PE's) in a mesh connected structure and a hierarchy of broadcast buses in each row and column of the mesh structure such that each broadcast bus is connected to exactly k PE's, where k is a small constant. It is also shown how to compute connected components and transitive closure of a given undirected

graph in $O(\log^2 n)$ with a few modifications to the algorithm presented for computing minimum spanning tree.

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#159
Comparison of Methods of Solving Symmetric
Positive Definite Linear Systems on a Hypercube
Concurrent Computer

The solution of partitioned symmetric positive definite linear systems on a Hypercube concurrent computer using Cholesky decomposition, conjugate gradient iteration, and preconditioned conjugate gradient iteration are compared. The partitionings we use, in their most complete form, are similar to nested dissection carried out to a depth of k, the dimension of the hypercube. We give timings from the Intel iPSC Hypercube at Oak Ridge, as well as counts of significant computation units such as matrix-vector multiplications and forward and backward solves, along with counts of communication requirements such as fan in fan outs and nearest neighbor exchanges.

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#161 A Simulation Tool for Optimizing Production Programs

Sim40 is a high-level tool which simulates pipelined vector processor instruction timing for the SCS-40 and Cray X-MP. Sim40 produces a graphic clock-by-clock display of instruction resource utilization which automatically highlights resource contentions preventing instruction issue. Critical code may thus be highly optimized without requiring the programmer to have a detailed knowledge of machine timing characteristics. The novel integration of Sim40 with a symbolic debugger provides a familiar user interface for simulation while adding a powerful "watch address" capability to the debugger. Sim40 is also used to evaluate the performance of new CPU designs.

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#164

Early Experience With the PASM System Prototype

The PASM (PArtitionable SIMD/MIMD) parallel processing system is dynamically reconfigurable in either MIMD or SIMD modes. This research involved performance evaluation of this multi-mode capability of the 16-processor PASM prototype applied to matrix multiplication. Variations of an $O(n^2)$ square matrix multiplication algorithm were examined for SIMD, MIMD and hybrid SIMD/MIMD modes of parallelism. Results indicated that the SIMD version of the algorithm achieved approximately 65% better efficiency than the MIMD version and was 25% better than the hybrid SIMD/MIMD version. This was due to the ability of the micro-controllers in PASM to execute flow control instructions in parallel with SIMD arithmetic and network transfer operations.

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#165 Iterative Algorithms in a Data Driven Environment

A large class of numerical algorithms is iterative. These iterative algorithms have been traditionally implemented, in conventional Von Neuman environments, with REPEAT-UNTIL and WHILE loops. These REPEAT-UNTIL constructs severely limit the performance of parallel processors because they cannot be vectorized and/or multitasked. By inserting a FORALL loop (do i=1,n) inside these REPEAT-UNTIL constructs performance of these iterative algorithms on parallel processors can be improved. The performance of the proposed scheme on the MIT tagged token data-flow architecture has been simulated. Jacobi and Gauss Sidel iterations were used as benchmark algorithms. The proposed scheme nearly doubles the speedup compared to the conventional REPEAT-UNTIL loops.

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#166

Computational and Reliability Measures of Homogeneous Multiprocessor Systems

The advent of parallel processing has opened new possibilities in scientific and numerical computing. Computations that would be prohibitively long on uniprocessors can now be carried on parallel machines in a more reasonable time. However, the likelihood of one or more elements failing increases with the complexity and the number of elements in the system. Therefore, the overall reliability and fault-tolerance becomes a key issue in the design and implementation of such systems. The measures proposed here aim at evaluating the achievable tradeoffs between faster execution on a larger number of processors and a more reliable computation. Computational oriented measures are proposed that evaluate the probability of a sage completion of a computation on a given syste, By expressing the reliability as function of the number of processor-hours instead of time one can determine the amount of reliable computational work a system can deliver.

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#168 Adaptive Grids of Multivariate, Nth-Order Calculi

This paper outlines some of the analytictopological problems of decompositions of n-dimensional spaces with relevance for statistical problems; (it constitutes the preparatory stage for 'Nth-Order Canonical Calculi' - in preparation). The extension of wave function approaches to basic statistical problems (with a view to creation of compartible differential geometry of statistics) generates new problems of decompositions. embeddings and imbeddings of (hyper) spaces. These problems, taken together analytischtopologically, constitute a form of 'intersection set' for unifying, extended 'Statistical Mechanics' for pattern recognition (e.g. neural pattern computing), (bio) chemical spatio-temporal trasformation of evolution structures, and monitored reliability of complex systems et al. This contribuion is on the precisions, open problems and hypotheses.

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#171 Compiler-Driven Cache Management

Cache performance is critical in cache-based supercomputers, where the cache-miss/cache-hit memory reference delay ratio is typically large. Using compile-time analysis, program behavior can be predicted, and cache control directives can be embedded in the generated code. Thus, cache performance can be improved in a way not possible using conventional techniques. Given hardware able to selectively bypass the cache, cache performance can be increased: pollution can be minimized. Cache line replacement can also be controlled by the compiler (rather than by LRU, etc.), further enhancing performance. This new methodology provides optimal, or near optimal, cache performance without complex hardware.

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#172

CRegs: A Hardware Solution to the Alias Problem

Often, pointer and subscripted array references touch memory locations for which there are several possible aliases, hence compilers are unable to optimize code implementing them. The CReg (Cache-Register) mechanism combines the hardware structures of cache and registers to create an entirely new associative memory structure, which can be used either as processor registers or as a replacement for conventional cache memory. By permitting aliased names to be grouped together, CRegs resolve ambiguous alias problems in hardware, thereby enabling compilers to generate more efficient code. Hardware implementing CRegs is of comparable complexity to that implementing a similar-sized conventional cache.

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#173 Barrier MIMD: Beyond VLIW

In terms of compiler technology, there is a wide gap between the static SIMD and dynamic MIMD control flow models. Static control brings efficient synchronization and greater opportunity for compile-time optimization, but the SIMD model incorporates many additional constraints. VLIW (Very Long Instruction Word) computers relax these, while preserving the benefits of static control. We propose the Static-Barrier machine model as an efficient extension of static control flow beyond the expressive power of VLIW. Syn-

chronization barriers permit efficient VLIW-like instruction scheduling, but also insure reasonable performance where ambiguous inter-processor timing occurs either through differences in control-flow or through unpredictable instruction execution delays.

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#174

Efficient Iterative Solution of Elliptic PDE's

Most iterative techniques can be very efficient provided one knows the value of one or several special parameters. For SOR, one needs to determine the optimal over-relaxation factor W. Recent theoretical results give new insight into this problem, and the local relaxation approach allows efficient practical implementation of SOR.

We will review the recent developments mentioned above and extend them to the general 9 point stencil case. In addition we observe that SOR does not seem to be well-suited for high-order stencils. Thus, another general iterative method is proposed and analyzed. This discussion will be illustrated with applications, including some Finite Element problems which were implemented on the Connection Machine Computer.

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#176 A MIMD Eigenvalue Algorithm for a Parallel Transputer - Microcomputer System

A parallel MIMD algorithm is developed for determining the eigenvalues of a large matrix. It is implemented in the INMOS OCCAM language for a microcomputer-based Transputer system with the capability of utilizing a large number of Transputers operating in parallel. The algorithm is based on the QR algorithm and is tested on an IBM AT microcomputer combined with up to eight Transputers. Comparisons are made to a QR algorithm on a large computer pertaining to the speed of computation, utilization and efficiency. Improvement and extension of these results will be discussed.

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#179
Parallel Conjugate Gradients on a 1024Node Hypercube*

The conjugate gradient method is a preferred approach for the solution of sparse, structured linear systems generated by finite difference and finite element analyses of many two and three dimensional problems. On an ensemble architecture such as the hypercube, much of the conjugate gradient method is easily made massively parallel. An exception is the inner-product calculation, which, for example, can be accomplished quickly using a logarithmic time-complexity exchange algorithm. We will present both theoretical and measured performance for these calculations on a 1024-processor hypercube ensemble.

*This work performed at Sandia National Laboratories supported by the U.S. Department of Energy under contract number DE-AC04-76DP00789.

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#181
Performance of Some CFD Loops on
Several Parallel Processors

Three loops representative of most algorithms for computational fluid dynamics (CFD) have been run on three different parallel processors in ord n to understand the suitability of diff parallel architectures to CFD ar ns. The three loops have quite dif memory access requirements, trouble with inherant depe . which would inhibt parallelism .chines .LLIANT FX/8, which have been used a the Intel iPSC, and Butterfly. The es will be loops and their pe presented.

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and
E. Ducot . V. Klema
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Massachusetts Institute of Technology Cambridge, MA 02139 #183
A Parallel Algorithm for Computing the Generalized
Singular Value Decomposition (GSVD)

One of the recent interesting topics in numerical linear algebra is the GSVD of any two matrices having the same number of columns. The GSVD provides a new tool for the study of matrix pencil. The advent of the wide applications of the GSVD has aroused much interesting in developing a parallel algorithm for computing the GSVD. We will present a parallel implementation of C.C.Paige's sequential Jacobi-like GSVD algorithm. The algorithm is designed to implement efficiently on distributed-memory parallel computer architectures, such as Hypercube and Finite Element Machine.

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#163 Task Migration in Partitionable Parallel Processing Systems*

In parallel processing systems that can be partitioned into independent submachines a problem that needs to be addressed is the migration of tasks across partition boundaries. Motivations for task migration include freeing larger partitions, fault tolerance, and load balancing. In this research, the overhead and penalties associated with task migration in parallel systems with multistage cube networks are examined. The overhead of moving tasks is divided into direct and indirect overhead. The sum of direct and indirect overhead represents all delays experienced by the task that is migrated. Migration penalty is introduced as the delay incurred by tasks in other partitions as a consequence of the migration. This information will be used to make task migration decisions.

*This research was supported by the Air Force Office of Scientific Research under grant number F49620-86-K-0006, and by the Supercomputing Research Center under contract number MDA904-85-C-5027.

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#185

Finding Public Domain, Mathematical Software

A service exists by which users can request and receive, via electronic mail, free public domain mathematical software. The service is called netlib. The service runs at Bell Laboratories in Murray Hill, NJ for UUCPnet and Argonne National Laboratories in Argonne, IL for ARPAnet.

Netlib operates without human intervention. Users must therefore be relatively specific in their request(s) to netlib. The database of software is built around a collection of independent libraries. An index for each libary is available to briefly describe its contents.

If a requestor is interested in a specific algorithm or topic, which is often the case, software for it may be present in several libraries. For instance, 95 routines to compute eigenvalues exist in 10 different libraries.

We detail here the *find* facility of *netlib* that allows users to locate routines, across all libraries, by using keyword(s). The requestor submits a *netlib* request to find a given keyword(s). Information is returned to briefly describe the capabilities of the routine(s), and the information necessary to submit another *netlib* request for the routine(s), relevant to the keyword.

Several examples are given to illustrate its usage.

† Jack Dongarra and Eric Grosse, "Distribution of Mathematical Software via Electronic Mail", Comm. ACM, 30(403-407), 1987.

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#186

The Lawrence Livermore Loops with a Vectorizing C Compiler

The Lawrence Livermore Loops, hereafter LLL, was originally authored by F.H. McMahon[†] in Fortran. It is intended as a measure of the floating-point performance that is relevant to 'scientific computing.' It consists of 24 loops, or kernels, each representing a commonly used, or generally important, scientific algorithm. One of the more common utilizations of the LLL benchmark is for an approximate characterization of todays 'supercomputers.'

Clearly the language of scientific computation and supercomputing is, and will be for some time to come, Fortran. The C language offers many features that are utilitarian and aesthetic to some scientists involved in computation. While C is not dominant in scientific computation it is becoming more popular.

With a vectorizing C compiler the scientist can retain the features of C and also achieve the vector speed-ups that were previously confined to Fortran on the advanced architecture machines.

In this paper we present the results of the LLL coded in C, compiled with an automatic vectorizing C compiler and executed on a vector minisupercomputer, the Convex C1.

† F.H. McMahon, "The Livermore Fortran Kernels: A Computer Test of the Numerical Performance Range," Lawrence Livermore Laboratory, UCRL-53745, December 1986.

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#194

A GRAPHICAL PROGRAMMING LANGUAGE FOR THE NAVIER-STOKES COMPUTER

Traditional programming techniques are sometimes inadequate for dealing with the complexities of parallel architectures. Graphical programming tools are emerging as an alternative for dealing with certain aspects of parallelism, especially the definition of processor interactions. We describe a prototype graphical editor/assembler which has been implemented for a reconfigurable pipeline machine, the Princeton/NASA Navier-Stokes Computer. With this tool, users write programs by interactively "drawing" the pipeline configuration and annotating it with additional information about control flow and memory access. These schematic representations of chained vector operations are then automatically translated to object code. The concepts employed are extendable to machines with similar architectures.

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#195

A Comparison of Different Parallel Programming Languages Using Choleski's Method on a Multi-Processor Computer

Choleski's method for solving banded symmetric, positive definite systems is implemented on a multiprocessor computer using three FORTRAN based parallel programming languages, the Force, PISCES and Concurrent FORTRAN. The capabilities of the languages for expressing parallelism and their user friendliness are discussed, including readability of the code, debugging assistance offered, and expressiveness of the languages. The performance of the different implementations is compared. It is argued that PISCES, using the Force for fine-grained parallelism, is the appropriate choice for programming Choleski's method on the multiprocessor computer, Flex/32.

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#1-8

Performance of Some CFD Codes on the Alliant FX/8

Three CFD codes are ported to the Alliant FX/8. The first solves the 3-D unsteady Euler equations using an explicit finite-volume Runge-Kutta time stepping method. The second solves the same problem using the Beam and Worming implicit finite difference method. The third is ARC2D from NASA Ames which solves the unsteady 2-D Navier-Stockes equations using an implicit method. We present extensive observations and results on the performance of these codes on the FX/8. Careful interaction with parrallelizing compiler improves the performance some. Better results are obtained by simple recoding of different segments in the programs.

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Experiences with the Use of the Cray X-MP/48 Performance Monitoring Tools

We report our experiences with the performance monitoring tools of the Cray X-MP/48. The goal is workload characterization of vectorized/multitasked/ and microtasked computational fluid dynamics codes. Our results revealed that the Cray tools are very informative when executing code on a single CPU. However, when codes are multitasked and/or microtasked, little multiprocessing relevant information could be obtained using the current tools. Some ideas on hardware/software instrumentation of the machine for workload characterization of multiprocessed codes are discussed.

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#L-11 Combined Vector and Parallel Processing on the IBM 3090

The IBM 3090 Computer System is capable of taking advantage of both vector and parallel features in a single FORTRAN program. Portions of code with large vector content can be run with reduced CPU time by taking advantage of the built-in 3090 Vector Facility. Portions of the code that are independent can be split off and run on separate processors of the multiprocessor 3090 system, thus reducing the elapsed time for running the entire program. For codes containing portions of independent code with vector content, the independent sections can be executed on separate processors each with their own vector facility. Vectorization of the code is done automatically by the VS FORTRAN Version 2 compiler. The use of the parallel processing capabilities can be exploited by the programmer by use of the VS FORTRAN multitasking facility, or automatically by the compiler by the recently announced Parallel Fortran Product available by special bid from IBM. The capabilities of the hardware and software for utilizing both the vector and parallel capabilities will be discussed. Examples of speedups in both CPU and elapsed runtime will be shown with applications currently taking advantage of these features.

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#L-12 Applications of Level 2 BLAS in the NAG Library

The Level 2 BLAS as specified by Dongarra et al (1986), "An extended set of Fortran basic linear algebra subprograms", Argonne National Laboratory Report, are becoming established as a standard interface to commonly occurring matrix-vector operations, for which efficient implementations are becoming available on various machines. The strategy for the NAG Library is to make calls to Level 2 BLAS wherever possible and maintain the portability of the higher level algorithms. We illustrate this by examples taken from the solution of non-linear equations and non-linear least squares problems. Performance measurements are quoted from various vector-processing machines in support of the strategy.

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#L-13 (POSTER SESS. 1, TUESDAY, DEC. 1, 5:00-7:00 PM)

Using Almost Block Diagonal Codes when Solving ODE's on Supercomputers

We describe briefly the two algorithms which are in common use for solving the almost block diagonal systems which arise in the collocation solution of boundary value solvers in ordinary differential equations and we analyse the cost of using each of them on a vector machine. We present our analysis in terms of the size and order of the differential system and the distribution of the boundary conditions. By considering some extreme cases we are able to demonstrate conclusively the degree of superiority of the favoured algorithm for this application. Finally we make some observations about solving boundary value problems using collocation methods on shared memory parallel machines, highlighting those areas where more work needs to be done.

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R. I. Hay Department of Statistics and Computational Maths. University of Liverpool Liverpool L69 3BX England #191 Generalized SOR - New Spectra from Old

Linear system $(I - \tilde{B})\tilde{x} = \tilde{f}$ is equivalent to $Ax = (I_n - B)x = f$ if, for arbitrary initial vectors y_0 and \tilde{y}_0 , the sequences $\{y_j\}, \{\tilde{y}_j\}$ by $y_j = By_{j-1} + f$, $\tilde{y}_j = B\tilde{y}_{j-1} + \tilde{f}$ $j = 0, 1, 2, 3, \cdots$ where $\{\tilde{y}_j\} \to \tilde{x}$. Solution vector $x = \pi(\tilde{x})$ must be easy to obtain from \tilde{x} . If, moreover, $\tilde{y}_j \to \tilde{x}$ faster than $\{y_j\} \to x$, then $\{\tilde{y}_j\}$ is said to be an acceleration of $\{y_j\}$. This is the case if and only if the spectral radii have the properties:

(1) $\rho(\tilde{B}) < \rho(B)$ and (2) $\rho(\tilde{B}) < 1$. The SOR theory, for example, supplies a parameterized family of alternate matrices

 $\{\tilde{B}=B_{\omega},\tilde{f}=f_{\omega}\}$ where $\tilde{x}=x$. We have $\tilde{y}_{j}\to\tilde{x}$ at least twice as fast as modified SOR when $\sigma(B)$ is pure imaginary, and $\tilde{y}_{j}\to\tilde{x}$ always converges when $\sigma(A)$ (hence, $\sigma(B)$) is real. In fact, our method converges for all systems where $\sigma(B)$ is either real or pure imaginary, even if $\rho(B)>1$. No symmetry, consistent ordering or positivity is ever required.

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NAME .	DAY	TIME	END-TIME	SESSION	ABST.PAGE	коом
Abdal Dahman T C	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh. Hall
Abdel-Rahman, T.S. Abello, James	Tue. PM	5:00	7:00	P/S 1	A27	Pasadena Exh. Hall
Abu-Sufah, Walid	Thurs. PM	5:00	7:00	P/S 3	A50	Pasadena Exh. Hall
Agerwala, Tilak	Fri. AM	8:30	9:15	I/P 13	4	San Diego
Aitchison, Peter	Thurs. PM	5:00	7:00	P/S 3	A47	Pasadena Exh. Hall
Allen, Keith	Thurs. PM	5:00	7:00	P/S 3	A41	Pasadena Exh. Hall
Allen, Richard	Thurs. AM	11:20	11:40	C/P 13	A15	San Gabriel A-C
Allen, Tom	Wed. AM	11:40	12:00	C/P 9	A10	Sacramento
Allik, Henno	Fri. AM	11:40	12:00	C/P 17	A19	Sacramento Sacramento
Anderson, Edward	Thurs. AM	10:40	11:00	C/P 14	A15 A25	Pasadena Exh. Hall
Anderson, R. E.	Tue. PM Wed. PM	5:00 5:00	7:00 7:00	P/S 1 P/S 2	A23 A37	Pasadena Exh. Hall
Andert, Ed Arbib, Michael	Tue. PM	1:30	2:15	I/P 3	2	San Diego
Armstrong, Jim	Tue. PM	5:00	7:00	P/S 1	A23	Pasadena Exh. Hall
Arnold, Clifford	Fri. AM	11:00	11:20	C/P 18	A20	Santa Barbara A-C
Arumugam, S.	Wed. PM	5:00	7:00	P/S 2	A37	Pasadena Exh. Hall
Ashcraft, C. C.	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
✓Astfalk, Greg	Thurs. PM	5:00	7:00	P/S 3	A49	Pasadena Exh. Hall
Atapattu, Daya	Wed. PM	4:40	5:00	C/P 10	A11	San Diego
Axelrod, Timothy	Wed. AM	11:40	12:00	C/P 7	A8	San Gabriel A-C
Axelsson, 0.	Thurs. PM	4:00	4:20	C/P 15	A17	San Diego
Babb, Robert G.	Tue. AM	11:20	11:40	C/P 1	A1	San Diego San Gabriel A-C
Baca, Lorraine	Thurs. AM	11:20 11:00	11:40 11:20	C/P 13 C/P 6	A15 A7	Santa Barbara A-C
Baden, Scott	Wed. AM Thurs. PM	5:00	7:00	P/S 3	A48	Pasadena Exh. Hall
Bai, Zhaojun	Thurs. PM	4:40	5:00	C/P 15	A17	San Diego
Bailey, David Barszcz, Eric	Tue. PM	5:00	7:00	P/S 1	A31	Pasadena Exh. Hall
Bassett, Mark	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall
Bauer, Thomas	Wed. PM	5:00	7:00	P/S 2	A37	Pasadena Exh. Hall
Baum, Alan	Wed. PM	5:00	7:00	P/S 2	A36	Pasadena Exh. Hall
Benner, Robert	Thurs. PM	5:00	7:00	P/S 3	A42	Pasadena Exh. Hall
✓Benner, Robert	Thurs. PM	5:00	7:00	P/S 3	A48	Pasadena Exh. Hall
Benten, Muhammad	Wed. AM	11:00	11:20	C/P 7	A8	San Gabriel A-C
Berger, Marsha	Tue. AM	9:15	10:00	I/P 2	2	San Diego
Herry, Michael	Wed. AM	11:40	12:00 7:00	C/P 8	A9 A32	San Diego Pasadena Exh. Hall
Bertram, Richard	Wed. PM	5:00 5:00	7:00	P/S 2 P/S 1	A29	Pasadena Exh. Hall
Bieterman, Michael	Tue. PM Tue. PM	4:20	4:40	C/P 5	A6	San Diego
Bischof, C. H. Bixler, J. P.	Wed. AM	11:00	11:20	C/P 9	A10	Sacramento
Blelloch, Guy	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh.Hall
Blue, James	Fri. AM	11:20	11:40	C/P 16	A18	San Gabriel A-C
Bonomo, John	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall
Boppana, R. V.	Thurs. PM	5:00	7:00	P/S 3	A45	Pasadena Exh. Hall
Borasky, M. Edward	Wed. PM	5:00	7:00	P/S 2	A32	Pasadena Exh. Hall
Boris, J. P.	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall Sacramento
Bosworth, Ken	Wed. AM	11:20	11:40	C/P 9	A10 A51	Pasadena Exh. Hall
Brankin, Richard	Thurs. PM	5:00 11:40	7:00 12:00	P/S 3 C/P 7	A31	San Gabriel A-C
Brooks, Eugene	Wed. AM Thurs. AM	11:40	11:40	C/P 12	Al4	San Diego
Brooks, Eugene		5:00	7:00	P/S 1	A27	Pasadena Exh. Hall
Brown, Donna	Tue. PM	2.00	7.00	1/01	, ALI	

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NAME	DAY	TIME	END-TIME	SESSION	ABST.PAGE	ROOM
Brussino, G.	Thurs. AM	11:00	11:20	C/P 14	A15	Sacramento
Calalo, Ruel	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Casavant, Thomas	Thurs. PM	5:00	7:00	P/S 3	A48	Pasadena Exh. Hall
Casavant, Thomas	Thurs. PM	5:00	7:00	P/S 3	A46	Pasadena Exh. Hall
Catherasoo, C. J.	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall
Cathey, W. Thomas	Tue. PM	2:15	3:00	I/P 4	2	San Diego
Celmaster, William	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall Santa Barbara A-C
Celmaster, William	Wed. AM	11:20 5:00	11:40 7:00	C/P 6 P/S 2	A7 A32	Pasadena Exh. Hall
Celmaster, William	Wed. PM Wed. AM	11:00	11:20	C/P 8	A9	San Diego
Chan, David	Tue. PM	5:00	7:00	P/S 1	A31	Pasadena Exh. Hall
Chan, Tony Chandran, Sharat	Fri. AM	10:40	11:00	C/P 16	A17	San Gabriel
Chi, Chi-Hung	Thurs. PM	5:00	7:00	P/S 3	A47	Pasadena Exh. Hall
Chiang, Yi-Ling	Tue. PM	5:00	7:00	P/S 1	A24	Pasadena Exh. Hall
Chin, Ray	Tue. AM	11:00	11:20	C/P 4	A5	Santa Barbara A-C
Chin, Ray	Wed. PM	2:15	3:00	I/P 8	3 .	San Diego
Chow, Shui-Nee	Wed. AM	10:40	11:00	C/P 9	A9 A28	Sacramento Pasadena Exh. Hall
Christidis, Z.	Tue. PM	5:00 11:40	7:00 12:00	P/S 1 C/P 3	A4	San Gabriel A-C
Chu, Eleanor	Tue. AM Tue. PM	5:00	7:00	P/S 1	A29	Pasadena Exh. Hall
Coray, Chris	Tue. PM	5:00	7:00	P/S 1	A31	Pasadena Exh. Hall
Crivelli, Luis	Thurs. AM	10:40	11:00	C/P 11	A12	Santa Barbara A-C
Crockett, Thomas	Thurs. PM	5:00	7:00	P/S 3	A49	Pasadena Exh. Hall
Eybenko, George	Wed. AM	11:40	12:00	C/P 9	A10	Sacramento
Daghi, Afshin	Tue. AM	11:20	11:40	C/P 2	А3	Sacramento
Daoud, Amjad	Thurs. PM	5:00	7:00	P/S 3	A50	Pasadena Exh. Hall
Darmody, L. M.	Thurs. PM	5:00	7:00	P/S 3	A43	Pasadena Exh. Hall San Gabriel, A-C
Darmohray, Gregory	Wed. AM	11:40	12:00	C/P 7	A8	
Darmohray, Gregory	Thurs. AM	11:20	11:40 7:00	C/P 12 P/S 1	A14 A23	San Diego Pasadena Exh. Hall
Datta, Biswa	Tue. PM Tue. PM	5:00 5:00	7:00	P/S 1	A23	Pasadena Exh. Hall
Datta, Karabi Davis, Larry	Fri. AM	10:40	11:00	C/P 16	A17	San Gabriel A-C
De la Sen, M.	Tue. PM	5:00	7:00	P/S 1	A22	Pasadena Exh. Hall
de Pillis, John	Thurs. PM	5:00	7:00	P/S 3	A51	Pasadena Exh. Hall
de Pillis, John	Tue. PM	5:00	7:00	P/S 1	A28	Pasadena Exh. Hall
de Pillis, Lisette	Tue. PM	5:00	7:00	P/S 1	A28	Pasadena Exh. Hall
deDoncker-Kapenga,E.		10:40	11:00	C/P 6	A6	Santa Barbara A-C San Gabriel A-C
deDoncker-Kapenga,E.		11:40	12:00 11:20	C/P 13 C/P 17	A15 A19	Sacramento
Dhall, Sudarshan	Fri. AM Fri. AM	11:00 11:00	11:20	C/P 17	A19	Sacramento
✓Diaz, Julio Dietz, Henry	Thurs. PM	5:00	7:00	P/S 3	A47	Pasadena Exh. Hall
DiNucci, David	Tue. AM	11:20	11:40	C/P 1	Al	San Diego
Dongarra, Jack	Thurs. AM	11:00	11:20	C/P 12	A13	San Diego
Dongarra, Jack	Thurs. PM	5:00	7:00	P/S 3	A49	Pasadena Exh. Hall
Dongarra, Jack	Thurs. PM		4:00	C/P 15	A16	San Diego
Doolen, Gary	Thurs. AM		10:00	I/P 10	4	San Diego Sacramento
Dowers, Ken	Fri. AM	11:00	11:20	C/P 17 P/S 3	A19 A43	Pasadena Exh. Hall
riscoll, Michael	Thurs. PM	5:00 11:00	7:00 11:20	C/P 12	A13	San Diego
Du Croz, Jeremy	Thurs. AM Thurs. PM	5:00	7:00	P/S 3	A51	Pasadena Exh. Hall
Du Croz, Jeremy	Thurs. AM		11:20	C/P 12	A13	San Diego
Duff, Iain	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Dunki-Jacobs, R.	Tue. AM	11:40	12:00	C/P 2	A3	Sacramento
Dyksen, John	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall
Edwards, Mickey	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall
Ægecioglu, Omer	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
✓Eijkhout, Victor	Wed. PM	5:00	7:00	P/S 2	A35	Pasadena Exh. Hall San Gabriel A-C
Eirola, Timo	Tue. AM	10:40	11:00	C/P 3 C/P 7	A3 A8	San Gabriel A-C
Eskow, Elizabeth	Wed. AM	11:20 5:00	11:40 7:00	P/S 3	A46	Pasadena Exh. Hall
Evripidou, P.	Thurs. PM Thurs. PM		3:00	I/P 12	4	San Diego
Faber, Vance	Thurs. AM		11:00	C/P 11	A12	Santa Barbara A-C
Favre, Jean Michel	Tue. PM	5:00	7:00	P/S 1	A29	Pasadena Exh. Hall
Fellows, Michael	Tue. PM	5:00	7:00	P/S 1	A27	Pasadena Exh. Hall
Ferreira, A. G.	Tue. PM	5:00	7:00	P/S 1	A31	Pasadena Exh. Hall
Fineberg, Samuel	Thurs. PM	5:00	7:00	P/S 3	A46	Pasadena Exh. Hall

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NAME	DAY	TIME	END-TIME	SESSION	ABST.PAGE	ROOM
Fisher, David	Tue. AM	10:40	11:00	C/P 2	A2	Sacramento
Fisher, David	Wed. PM	5:00	7:00	P/S 2	A35	Pasadena Exh. Hall
Foster, Richard	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh. Hall
Foulser, David	Tue. PM	5:00	7:00	P/S 1	A26 A36	Pasadena Exh. Hall Pasadena Exh. Hall
Fox, Geoffrey	Wed. PM Fri. AM	5:00 10:40	7:00 11:00	P/S 2 C/P 17	A19	Sacramento
Frederick, Jeffrey Frederickson, Paul	Thurs. PM	5:00	7:00	P/S 3	A45	Pasadena Exh. Hall
Furmanski, Wojtek	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall
Gajski, Daniel	Fri. AM	11:40	12:00	C/P 18	A20	Santa Barbara A-C
Gallivan, Kyle	Thurs. PM	5:00	7:00	P/S 3	A44	Pasadena Exh. Hall
Callopoulos, E.	Tue. PM	3:40	4:00	C/P 5	A5	San Diego
Gallopoulos, E.	Thurs. AM	4:20	4:40	C/P 15 C/P 10	A17 A11	San Diego San Diego
Gannon, Dennis Gaudiot, J. L.	Wed. PM Thurs. PM	4:40 5:00	5:00 7:00	P/S 3	A46	Pasadena Exh. Hall
Geist, George	Thurs. AM	11:40	12:00	C/P 12	A14	San Diego
George, J. Alan	Tue. AM	11:40	12:00	C/P 3	A4	San Gabriel A-C
Gerbi, Stephane	Thurs. AM	11:40	12:00	C/P 14	A16	Sacramento
Gonzalez-Gomez, M.	Tue. PM	5:00	7:00	P/S 1	A22	Pasadena Exh. Hall
Gottlieb, David	Tue. AM	8:30	9:15	I/P 1	2	San Diego
Greengard, Leslie	Thurs. AM	8:30	9:15	I/P 9	3	San Diego Pasadena Exh. Hall
Grimes, R. G.	Tue. PM	5:00	7:00	P/S 1 C/P 19	A25 A21	San Diego
Gropp, William	Fri. AM Thurs. PM	10:40 5:00	11:00 7:00	P/S 3	A49	Pasadena Exh. Room
Grosse, Eric Gustafson, John	Thurs. PM	5:00	7:00	P/S 3	A48	Pasadena Exh. Hall
Haghoo, Majid	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
Haidvogel, Dale	Fri. AM	11:00	11:20	C/P 16	A18	San Gabriel A-C
Hammarling, Sven	Thurs. AM	11:00	11:20	C/P 12	A13	San Diego
Hammond, Steven	Tue. AM	11:40	12:00	C/P 2	A3 A27	Sacramento Pasadena Exh. Hall
Hanson, Floyd	Tue. PM	5:00	7:00 12:00	P/S 1 C/P 2	A27 A3	Sacramento
Hardy, Robert	Tue. AM Wed. AM	11:40 11:00	11:20	C/P 9	A10	Sacramento
Harimoto, J. Harrod, Bill	Tue. AM	11:20	11:40	C/P 3	A4	San Gabriel A-C
Hebert, D. J.	Tue. AM	11:40	12:00	C/P 4	A5	Santa Barbara A-C
Hedstrom, G. W.	Tue. AM	11:00	11:20	C/P 4	A5	Santa Barbara A-C
Merbin, R.	Thurs. AM	11:00	11:20	C/P 14	A15	Sacramento
Herbin, K.	Thurs. AM	11:40	12:00	C/P 14	A16	Sacramento
Ho, Alex W.	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall Pasadena Exh. Hall
Howard, Steve	Thurs. PM	5:00	7:00	P/S 3	A45 A38	Pasadena Exh. Hall
Heu Heigo-Ming	Wed. PM Fri. AM	5:00 11:00	7:00 11:20	P/S 2 C/P 16	A18	San Gabriel A-C
Hsu, Hsiao-Ming ✓Mwang, Kai	Wed. PM	3:40	4:00	C/P 10	Ali	San Diego
Hwang, Kai	Thurs. AM	11:20	11:40	C/P 14	A16	Sacramento
Hwang, Phillip	Fri. AM	11:20	11:40	C/P 18	A20	Santa Barbara A-C
Hyatt, Robert	Thurs. PM	5:00	7:00	P/S 2	A42	Pasadena Exh. Hall
Imbriale, William	Wed. PM	5:00	7:00	P/S 2 C/P 19	A34 A21	Pasadena Exh. Hall San Diego
✓Ipsen, Ilse Israeli, Moshe	Fri. AM Wed. PM	10:40 5:00	11:00 7:00	P/S 2	A41	Pasadena Exh. Hall
Jacobs, Ken	Thurs. PM	5:00	7:00	P/S 3	A50	Pasadena Exh. Hall
Jalby, William	Thurs. PM	5:00	7:00	P/S 3	A44	Pasadena Exh.Hall
Jeffrey, William	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall
Jespersen, Dennis	Tue. PM	5:00	7:00	P/S 1	A31	Pasadena Exh. Hall
Lessup, E. R.	Wed. AM	11:20	11:40 7:00	C/P 8 P/S 2	A9 A40	San Diego Pasadena Exh. Hall
Johnsson, S. L. Jones, Mark	Wed. PM Thurs. PM	5:00 -5:00	7:00	P/S 3	A50	Pasadena Exh. Hall
Jordan, Harry	Wed. AM	11:00	11:20	C/P 7	A8	San Gabriel A-C
Julien-Laferriere, P.		5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
Kainen, Paul	Wed. PM	5:00	7:00	P/S 2	A37	Pasadena Exh. Hall
Kapenga, John	Wed. AM	10:40	11:00	C/P 6	A6	Santa Barbara A-C
Kapenga, John	Thurs. AM	11:40	12:00	C/P 13	A15	San Gabriel A-C
Katti, C. P.	Wed. PM Wed. AM	5:00 8:30	7:00 9:15	P/S 2 I/P 5	A36 3	Pasadena Exh. Hall San Diego
Kennedy, Ken Keyes, David	Thurs. AM	11:40	12:00	C/P 11	A13	Santa Barbara A-C
King, Chung-Ta	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
King, chung ia	Wed. PM	5:00	7:00	P/S 2	A36	Pasadena Exh. Hall
Koc, Cetin	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
Konno, Chisato	Tue. AM	11:40	12:00	C/P 1	A2	San Diego
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Kosloff, Dan	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall
Krueger, Robert	Thurs. AM Wed. PM	11:20 5:00	11:40 7:00	C/P 11	A12	Santa Barbara A-C
Lakshmivarahan, S.	Fri. AM	11:00	11:20	P/S 2 C/P 17	A34 A19	Pasadena Exh. Hall Sacramento
Lambrakos, S. G.	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall
Lancaster, Dale	Thurs. PM	5:00	7:00	P/S 3	A49	Pasadena Exh. Hall
Langston, Michael	Tue. PM	5:00	7:00	P/S 1	A27	Pasadena Exh. Hall
, Laub, Alan J.	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
Lee, Allisen	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
Lee, Daeshik	Thurs. PM	4:20	4:40	C/P 15	A17	San Diego
Lee, Mann Ho	Wed. PM	4:40	5:00	C/P 10	All	San Diego
Lee, Rong Lei, Li	Thurs. PM Tue. PM	5:00 5:00	7:00 7:00	P/S 3 P/S 1	A42 A25	Pasadena Exh. Hall Pasadena Exh. Hall
Lei, Li	Tue. PM	5:00	7:00	P/S 1	A28	Pasadena Exh. Hall
Leiserson, C. E.	Tue. AM	11:00	11:20	C/P 2	A2	Sacramento
Leuze, Michael	Wed. AM	11:40	12:00	C/P 6	A7	Santa Barbara A-C
Lewis, John	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
Lewis, John G.	Tue. AM	11:00	11:20	C/P 2	A2	Sacramento
Liddell, Heather	Wed. AM	10:40	11:00	C/P 7	A7	San Gabriel A-C
Lie, Ivar	Tue. PM	5:00	7:00	P/S 1	A27	Pasadena Exh. Hall
Liewer, Paulett C.	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Lin, Avi Lindheim, Jan	Wed. AM	10:40	11:00	C/P 8	A8	San Diego
Louri, Ahmed	Tue. PM Wed. PM	5:00 3:40	7:00 4:00	P/S 1 C/P 10	A28 All	Pasadena Exh. Hall San Diego
Lubeck, Olaf	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Luh, Peter	Thurs. PM	5:00	7:00	P/S 3	A42	Pasadena Exh. Hall
Lyons, James	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
zenga, Gregory	Wed. PM	5:00	7:00	P/S 2	A35	Pasadena Exh. Hall
Ma, Shing	Tue. AM	11:00	11:20	C/P 3	A3	San Gabriel A-C
Magnan, Jerry	Wed. PM	5:00	7:00	P/S 2	A41	Pasadena Exh. Hall
Magnan, Jerry	Wed. PM	5:00	7:00	P/S 2	A41	Pasadena Exh. Hall
Magnan, Jerry F. Mahajan, Umesh	Wed. PM Wed. PM	5:00 5:00	7:00 7:00	P/S 2 P/S 2	A32 A36	Pasadena Exh. Hall Pasadena Exh. Hall
Malone, James	Fri. AM	11:40	12:00	C/P 19	A21	San Diego
Mandell, David	Tue. PM	5:00	7:00	P/S 1	A22	Pasadena Exh. Hall
Mansfield, Lois	Thurs. PM	5:00	7:00	P/S 3	A45	Pasadena Exh. Hall
Manshadi, Farzin	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Margaret L. Simmons	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Marinescu, Dan	Wed. PM	4:00	4:20	C/P 10	All	San Diego
Mayes, Peter Mbaeyi, Peter N. O.	Thurs. PM Thurs. PM	5:00 5:00	7:00 7:00	P/S 3 P/S 3	A51 A46	Pasadena Exh. Hall
McBryan, Oliver	Wed. PM	4:20	4:40	C/P 10	All	Pasadena Exh. Hall San Diego
McMillan, Donald	Wed. PM	5:00	7:00	P/S 2	A36	Pasadena Exh. Hall
McMillin, Bruce	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
Mehrotra, Piyush	Fri. AM	9:15	10:00	I/P 14	4	San Diego
Meier, Ulrike	Thurs. PM	5:00	7:00	P/S 3	A44	Pasadena Exh. Hall
Mesirov, Jill	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh. Hall
Meurant, Gerard Meyer, Gerard	Tue. PM Fri. AM	4:40 11:40	5:00 12:00	C/P 5 C/P 16	A6	San Diego
Meyer, Robert	Wed. PM	5:00	7:00	P/S 2	A18 A38	San Gabriel A-C Pasadena Exh. Hall
Middleton, David	Thurs. PM	5:00	7:00	P/S 3	A49	Pasadena Exh. Hall
Mikolajczak, B.	Tue. PM	5:00	7:00	P/S 1	A25	Pasadena Exh. Hall
Miya, E. N.	Wed. PM	5:00	7:00	P/S 2	A35	Pasadena Exh. Hall
Mong, A. C.	Tue. PM	5:00	7:00	P/S 1	A25	Pasadena Exh. Hall
Montry, Gary	Wed. PM	5:00	7:00	P/S 2	A36	Pasadena Exh. Hall
Montry, Gary	Thurs. PM	5:00	7:00	P/S 3	A42	Pasadena Exh. Hall
Morrow, Dennis Mudge, Trevor	Wed. PM Wed. PM	5:00 5:00	7:00 7:00	P/S 2 P/S 2	A32	Pasadena Exh. Hall Pasadena Exh. Hall
Nagumo, M.	Wed. PM	5:00	7:00	P/S 2	A40 A33	Pasadena Exh. Hall
Magurney, Anna	Tue. PM	5:00	7:00	P/S 1	A30	Pasadena Exh. Hall
Naik, Vijay	Thurs. AM	10:40	11:00	C/P 12	A13	San Diego
✓Najjar, Walid	Thurs. PM	5:00	7:00	P/S 3	A46	Pasadena Exh. Hall
Mash, Stephen	Tue. PM	5:00	7:00	P/S 1	A24	Pasadena Exh. Hall
Welson, Paul	Wed. PM	5:00	7:00	P/S 2	A36	Pasadena Exh. Hall
Neta, Beny	Wed. PM	5:00	7:00	P/S 2	A36	Pasadena Exh. Hall
Nevanlinna, O.	Tue. AM	10:40	11:00	C/P 3	A3	San Gabriel A-C

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	NAME	DAY	TIME	END-TIME	SESSION	ABST.PAGE	ROOM
	Nevanlinna, 0.	Fri. AM	11:00	11:20	C/P 19	A21	San Diego
	-₩i, Lionel	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
	Ni, Lionel	Wed. AM	10:40	11:00	C/P 9	A9	Sacramento
L	Michols, Kathleen	Thurs. PM		7:00	P/S 3	A43	Pasadena Exh. Hall
	Norsett, Syvert	Tue. PM Thurs. AM	5:00 11:00	7:00 11:20	P/S 1 C/P 11	A27 A12	Pasadena Exh. Hall Santa Barbara A-C
	Nour-Omid, Bahram Nour-Omid, Bahram	Wed. PM	5:00	7:00	P/S 2	A35	Pasadena Exh. Hall
	O'Neil, E. J.	Fri. AM	11:40	12:00	C/P 17	A19	Sacramento
	Oran, E. S.	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall
	Ortiz, M.	Thurs. AM	11:00	11:20 7:00	C/P 11 P/S 3	A12 A41	Santa Barbara A-C Pasadena Exh. Hall
	∽Pargas, Roy Pariser, Elissa	Thurs. PM Thurs. PM	5:00 5:00	7:00	P/S 3	A43	Pasadena Exh. Hall
	Parkinson, D.	Wed. AM	10:40	11:00	C/P 7	A7	San Gabriel A-C
	Patrick, Merrell	Thurs. AM	10:40	11:00	C/P 12	A13	San Diego
	Patrick, Merrell	Thurs. PM	5:00	7:00	P/S 3	A50 A3	Pasadena Exh. Hall San Gabriel A-C
*	Patrick, Merrell -Patrick, Merrell	Tue. AM Wed. PM	11:00 5:00	11:20 7:00	C/P 3 P/S 2	A36	Pasadena Exh. Hall
	Patterson, Jean	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
	Pattipati, Krishna	Thurs. PM	5:00	7:00	P/S 3	A42	Pasadena Exh. Hall
,	Peir, Jih-Kwon	Fri. AM	11:00	11:20	C/P 16	A18 A24	San Gabriel A-C Pasadena Exh. Hall
ب	Perkins, A. Louise	Tue. PM Tue. PM	5:00 5:00	7:00 7:00	P/S 1 P/S 1	A24 A29	Pasadena Exh. Hall
•	Pernice, Michael Detersen, Johnny	Tue. PM	5:00	7:00	P/S 1	A28	Pasadena Exh. Hall
	Peyton, Barry	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
	Podrazik, Louis	Fri. AM	11:40	12:00	C/P 16	A18 A19	San Gabriel A-C Sacramento
V	Poole, Eugene	Fri. AM Wed. PM	10:40 5:00	11:00 7:00	C/P 17 P/S 2	A33	Pasadena Exh. Hall
	Pothen, Alex Prasanna-Kumar, V.	Tue. AM	11:20	11:40	C/P 2	A3	Sacramento
	Price, C. C.	Tue. PM	5:00	7:00	P/S 1	A22	Pasadena Exh. Hall
L	Proskurowski, W.	Tue. PM	5:00	7:00	P/S 1	A26	Pasadena Exh. Hall
-	Puckett, J. A. Puckett, J. A.	Thurs. AM Thurs. AM	10:40 11:00	11:00 11:20	C/P 13 C/P 13	A14 A14	San Gabriel A-C San Gabriel A-C
	Raefsky, Arthur	Wed. PM	5:00	7:00	P/S 2	A35	Pasadena Exh. Hall
	Raghavan, Padma	Wed. PM	5:00	7:00	P/S 2	A33	Pasadena Exh. Hall
ا	Raghavendra, C. S.	Thurs. PM	5:00	7:00	P/S 3	A45	Pasadena Exh. Hall Pasadena Exh. Hall
_	Rajaraman, V.	Tue. PM Wed. PM	5:00 5:00	7:00 7:00	P/S 1 P/S 2	A23 A32	Pasadena Exh. Hall
	Ramamurthy, Mohan Reddaway, Stewart	Fri. AM	11:20	11:40	C/P 17	A19	Sacramento
سا	Reed, Daniel	Wed. PM	1:30	2:15	I/P 7	3	San Diego
	Rees, Michael	Tue. AM	11:20	11:40 7:00	C/P 4 P/S 2	A5 A38	Santa Barbara A-C Pasadena Exh. Hall
	Reshef, Moshe Reynders, J. V. W.	Wed. PM Tue. PM	5:00 5:00	7:00	P/S 1	A30	Pasadena Exh. Hall
11	Pribbene Calvin	Fri. AM	11:20	11:40	C/P 19	A21	San Diego
ب س 1	Rice, John	Wed. PM	4:00	4:20	C/P 10	All	San Diego
	Routile, Charles	Thurs. AM	11:40	12:00	C/P 12	A14	San Diego San Diego
	Rosing, M. Rosner, Robert	Tue. AM Wed. PM	11:00 5:00	11:20 7:00	C/P 1 P/S 2	A1 A38	Pasadena Exh. Hall
	Roy, Joydeep	Thurs. PM	5:00	7:00	P/S 3	A42	Pasadena Exh. Hall
	Ruttenberg, Alan	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh. Hall
	Saad, Youcef	Tue. PM	3:40	4:00	C/P 5	A5 A44	San Diego Pasadena Exh. Hall
	Saad, Youcef Saad, Youcef	Thurs. PM Thurs. AM	5:00 10:40	7:00 11:00	P/S 3 C/P 14	A15	Sacramento
	Sadayappan, P.	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
_	Safavi, Ali	Tue. AM	11:20	11:40	C/P 2	A3	Sacramento
,	Saji, Miyuki	Tue. AM	11:40	12:00	C/P 1	A2	San Diego
L	Salama, M. A.	Tue. PM	5:00	7:00	P/S 1	A22 4	Pasadena Exh. Hall San Diego
	Saltz, Joel Sameh, Ahmed	Thurs. PM Tue. AM	1:30 11:20	2:15 11:40	I/P 11 C/P 3	A4	San Gabriel A-C
	Sameh, Ahmed	Wed. AM	11:40	12:00	C/P 8	A9	San Diego
	Sameh, Ahmed	Thurs. PM		7:00	P/S 3	A44	Pasadena Exh. Hall
	Sandee, Daan	Wed. PM	5:00 5:00	7:00 7:00	P/S 2	A32 A35	Pasadena Exh. Hall Pasadena Exh. Hall
	Sattler, Linda Schach, Stephen	Wed. PM Wed. AM	5:00 11:40	7:00 12:00	P/S 2 C/P 6	A33 A7	Santa Barbara A-C
	Schaefer, Mark	Tue. AM	11:20	11:40	C/P 3	A4	San Gabriel A-C
	Schmidt, R. J.	Thurs. AM	10:40	11:00	C/P 13	A14	San Gabriel A-C
	Schmidt, R. J.	Thurs. AM	11:00	11:20	C/P 13	A14	San Gabriel A-C

 $[\]overline{C/P}$ = Contributed Presentation I/P = Invited Presentation P/S = Poster Session

NAME	DAY	TIME	END-TIME	SESSION	ABST.PAGE	ROOM
Schnabel, Robert	Tue. AM	11:00	11:20	C/P 1	Al	San Diego
Schnabel, Robert	Wed. AM	11:20	11:40	C/P 7	A8	San Gabriel A-C
Schwederski, Thomas	Thurs. PM	5:00	7:00	P/S 3	A48	Pasadena Exh. Hall
Schwederski, Thomas Schwederski, Thomas	Thurs. PM Thurs. PM	5:00 5:00	7:00 7:00	P/S 3 P/S 3	A46 ⁻ A47	Pasadena Exh. Hall Pasadena Exh. Hall
croggs, Jeffrey	Tue. AM	11:00	11:20	C/P 4	A5	Santa Barbara A-C
Shei, Bruce	Wed. PM	4:40	5:00	C/P 10	All	San Diego
Shelley, Michael	Wed. PM	5:00	7:00	P/S 2	A41	Pasadena Exh. Hall
Shen, Yun-Qiu	Wed. AM	10:40	11:00	C/P 9	A9	Sacramento
Shoaff, William	Wed. AM	11:00 5:00	11:20 7:00	C/P 8 P/S 3	A9 A48	San Diego Pasadena Exh. Hall
Siegel, Howard Jay Simmons, Margaret	Thurs. PM Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Simon, Horst	Tue. PM	5:00	7:00	P/S 1	A25	Pasadena Exh. Hall
Simon, Robert	Wed. PM	5:00	7:00	P/S 2	A38	Pasadena Exh. Hall
Sincovec, R. F.	Tue. AM	10:40	11:00	C/P 1	Al	San Diego
Siva Ram Murthy, C. Sjogreen, Bjorn	Tue. PM Tue. AM	5:00 10:40	7:00 11:00	P/S 1 C/P 4	A23 A4	Pasadena Exh. Hall Santa Barbara A-C
Smooke, Mitchell	Thurs. AM	11:40	12:00	C/P 11	A13	Santa Barbara A-C
Sonnad, V.	Thurs. AM	11:00	11:20	C/P 14	A15	Sacramento
Sonnad, V.	Thurs. AM	11:40	12:00	C/P 14	A16	Sacramento
Sonnad, V.	Tue. PM	5:00	7:00	P/S 1	A28	Pasadena Exh. Hall
Sorensen, D. C.	Tue. AM Wed. AM	11:00 11:20	11:20 11:40	C/P 4 C/P 8	A5 A9	Santa Barbara A-C San Diego
Sorensen, D. C. Sorensen, D. C.	Thurs. PM	3:40	4:00	C/P 15	Al6	San Diego
Frinidhi, Hassan	Thurs. PM	5:00	7:00	P/S 3	A44	Pasadena Exh. Hall
Stevenson, D. E.	Tue. PM	5:00	7:00	P/S 1	A30	Pasadena Exh. Hall
Stiles, G. S.	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
Stiles, G. S.	Wed. AM	11:20	11:40	C/P 9	A10 A43	Sacramento Pasadena Exh. Hall
Sugla, Binay Cillivan, Francis	Thurs. PM Fri. AM	5:00 11:20	7:00 11:40	P/S 3 C/P 16	A18	San Gabriel A-C
Suter, Bruce	Thurs. RM		7:00	P/S 3	A42	Pasadena Exh. Hall
Szyld, Daniel	Tue. AM	11:00	11:20	C/P 3	А3	San Gabriel A-C
Tang, Wei Pai	Wed. AM	9:15	10:00	I/P 6	3	San Diego
Taylor, Ian	Wed. PM	5:00	7:00	P/S 2	A37	Pasadena Exh. Hall
Tenenbaum, Eric	Wed. PM Thurs. PM	5:00 5:00	7:00 7:00	P/S 2 P/S 3	A38 A42	Pasadena Exh. Hall Pasadena Exh. Hall
Tesch, John	Thurs. PM	5:00	7:00	P/S 3	A50	Pasadena Exh. Hall
Thomson, Chris	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh. Hall
Tomboulian, Sherryl	Thurs. PM	5:00	7:00	P/S 3	A49	Pasadena Exh. Hall
Topka, Terry	Tue. AM	11:40	12:00	C/P 2	A3	Sacramento
Tourancheau, B. Tripodes, Peter	Tue. PM Wed. PM	5:00 5:00	7:00 7:00	P/S 1 P/S 2	A31 A37	Pasadena Exh. Hall Pasadena Exh. Hall
Tuminaro, Kay	Tue. PM	5:00	7:00	P/S 1	A31	Pasadena Exh. Hall
Umetani, Yukio	Tue, AM	11:40	12:00	C/P 1	A2	San Diego
Mandewalle, Stefan	Tué. PM	4:00	4:20	C/P 5	A6	San Diego
Varma, Anujan	Thurs. PM	5:00	7:00	P/S 3	A45	Pasadena Exh. Hall
Vasilevsky, Alex Venkatachalam, P.A.	Wed. PM Wed. PM	5:00 5:00	7:00 7:00	P/S 2 P/S 2	A 40 A37	Pasadena Exh. Hall Pasadena Exh. Hall
Vincent, Ina	Thurs. PM	5:00	7:00	P/S 3	A47	Pasadena Exh. Hall
Visvanathan, V.	Wed. PM	5:00	7:00	P/S 2	A39	Pasadena Exh. Hall
Voigt, Robert	Thurs. PM	5:00	7:00	P/S 3	A50	Pasadena Exh. Hall
Waleffe, Fabian	Thurs. PM Wed. PM	5:00	7:00	P/S 3	A47	Pasadena Exh. Hall
Walker, David	Thurs. AM	5:00 11:20	7:00 11:40	P/S 2 C/P 14	A36 A16	Pasadena Exh. Hall Sacramento
Wasserman, Harvey	Wed. PM	5:00	7:00	P/S 2	A34	Pasadena Exh. Hall
Watson, Layne	Wed. AM	11:00	11:20	C/P 9	A10	Sacramento
Wilson, Alan	Tue. PM	5:00	7:00	P/S 1	A27	Pasadena Exh. Hall
Wilson, Dan	Wed. PM	5:00	7:00	P/S 2	A40	Pasadena Exh. Hall
Witkowski, Andrew Wolfe, Michael	Tue. PM Fri. AM	5:00 10:40	7:00 11:00	P/S 1 C/P 18	A30 A20	Pasadena Exh. Hall Santa Barbara A-C
omble, David	Thurs. AM	11:20	11:40	C/P 13	A15	San Gabriel A-C
Woodyard, Mark	Thurs. PM	5:00	7:00	P/S 3	A44	Pasadena Exh. Hall
Wu, Min-You	Fri. AM	11:40	12:00	C/P 18	A20	Santa Barbara A-C
Yamabe, Michiru	Tue. AM	11:40	12:00	C/P 1	A2	San Diego
Zhang, He Zhaoyong, You	Wed. AM Tue. PM	10:40 5:00	11:00 7:00	C/P 8 P/S 1	A8 A25	San Diego Pasadena Exh. Hall
Zhaoyong, You	Tue. PM	5:00	7:00	P/S 1	A28	Pasadena Exh. Hall
Zhu, Sheng-Bai	Tue. PM	5:00	7:00	P/S 1	A29	Pasadena Exh. Hall

 $[\]overline{C/P}$ = Contributed Presentation I/P = Invited Presentation P/S = Poster Session

REGISTRATION INFORMATION

The registration desk will be located in the California Foyer on the second floor of the hotel and will be open as listed below.

Monday, November 30 Tuesday, December 1 Wednesday, December 2 Thursday, December 3

Friday, December 4

4:00 PM-8:00 PM 7:00 AM-7:30 PM 7:30 AM-6:00 PM

7:30 AM-6:00 PM 7:30 AM-1:00 PM

	•	SIAG/ SC*	SIAM Member	Non Member	Students
Conference	Advance	\$ 90	\$ 95	\$125	\$20
Fees	On-Site	\$120	\$125	\$155	\$20

^{*} Discounted fee for members of the SIAM Activity Group on Supercomputing

Beef and Beer Party

Tuesday, December 1 7:30 PM Pasadena Exhibit Hall Westin Bonaventure \$15.00

Non SIAM Members

Non-member registrants are encouraged to join SIAM in order to obtain the member rate for conference registration and enjoy all the other benefits of SIAM membership.

Special Note

There will be no prorated fees. No refunds will be issued once the conference has started.

Telephone Messages

The telephone number at the Westin Bonaventure is 1-213-624-1000. The Westin Bonaventure will either connect you to the SIAM desk or forward a message. Messages received at the SIAM registration desk are posted on a bulletin board in the registration area.

SIAM CORPORATE MEMBERS

Non-member attendees who are employed by the following institutions are entitled to the SIAM member rate.

Aerospace Corporation Amoco Production Company AT&T Bell Laboratories **Bell Communications Research Boeing Company** Cray Research, Inc. Culler Scientific Systems Corporation E.I. Du Pont de Nemours and Company Eastman Kodak Company Exxon Research and Engineering Company General Electric Company **General Motors Corporation** Giers Schlumberger GTE Laboratories, Inc. Hollandse Signaalapparaten B.V. IBM Corporation Institute for Computer Applications in Science and Engineering (ICASE) IMSL, Inc. MacNeal-Schwendler Corporation Marathon Oil Company Martin Marietta Energy Systems Mathematical Sciences Research Institute Standard Oil Compay of Ohio (SOHIO) Supercomputing Research Center, a division of Institute for Defense Analyses Texaco, Inc. United Technologies Corporation

Credit Cards

SIAM is now accepting credit cards for the payment of registration fees and special functions.

GENERAL INFORMATION

EXHIBITS

Software/Hardware Books/Journals

The exhibit area is in the Pasadena Exhibit Hall of the hotel. The exhibit times are 9:30 AM-10:00 PM, Tuesday, December 1; 9:30 AM-7:00 PM, Wednesday and Thursday, December 2 and 3; and 9:30 AM-1:00 PM Friday, December 4. The exhibits setup time will begin at 2:00 PM, Monday, November 30; breakdown will begin at 1:00 PM and will end at 4:00 PM, Friday, December 4.

Special Notice to: All Conference Participants

SIAM requests conferees to refrain from smoking in the session rooms during lectures. Thank you.

UPCOMING CONFERENCES

May 23 – 26, 1988

Third SIAM Conference on Linear Algebra
The Concourse Hotel
Madison, WI

June 13-16, 1988
Fourth SIAM Conference on Discrete
Mathematics
Cathedral Hill Hotel
San Francisco, CA

July 11 – 15, 1988 SIAM Annual Meeting Hyatt Regency Hotel Minneapolis, MN

Third SIAM Conference on Parallel Processing for Scientific Computing

December 1-4, 1987

Los Angeles, California

EXHIBITORS

Software/Hardware

Active Memory Technology Intel Scientific Computers

Alliant Computer Systems Kuck & Associates, Inc.

Amdahl Corporation Multiflow Computer, Inc.

Applied Dynamics NCUBE

Convex Computer Corp. Numerical Algorithms Group

Floating Point Systems Numerix Corporation

IBM Corporation Saxpy Computer Corporation

Sequent Computer Systems

Publishers

Academic Press, Inc. Cambridge University Press

American Institute of Physics North Holland—A Division of Elsevier Science Publishers B.V.

American Society of Mechanical

Engineers Penn State Press

Baywood Publishing Company, Inc. Plenum Publishing

Butterworth Publishers SIAM

Note: This list includes only those exhibitors who submitted their contracts for exhibit space prior to the printing of the final program.